

# Convex Optimization

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# Convex Optimization

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## ■ General Optimization Problem

$$\begin{aligned} & \text{minimize } f_0(x) \\ & \text{subject to } f_i(x) \leq 0, \quad i = 1, 2, \dots, m \\ & \quad \quad \quad h_i(x) = 0, \quad i = 1, 2, \dots, p \end{aligned}$$

$$x = (x_1, x_2, \dots, x_n)$$

$$f_0 : \mathbb{R}^n \rightarrow \mathbb{R}$$

$$f_i : \mathbb{R}^n \rightarrow \mathbb{R}$$

$$h_i : \mathbb{R}^n \rightarrow \mathbb{R}$$

## ■ Example

$$\begin{aligned} & \text{minimize } -\int_0^\infty B \log_2(1 + P_t(\gamma)\gamma/\bar{P})p(\gamma)d\gamma \\ & \text{subject to } \int_0^\infty (P_t(\gamma)/\bar{P})p(\gamma)d\gamma \leq 1 \end{aligned}$$

- A problem is a **convex optimization problem** if
  - Objective function & Inequality constraint functions: convex
  - Equality constraint functions: linear (or affine)
- **Convexity** guarantees some nice properties to solve it.

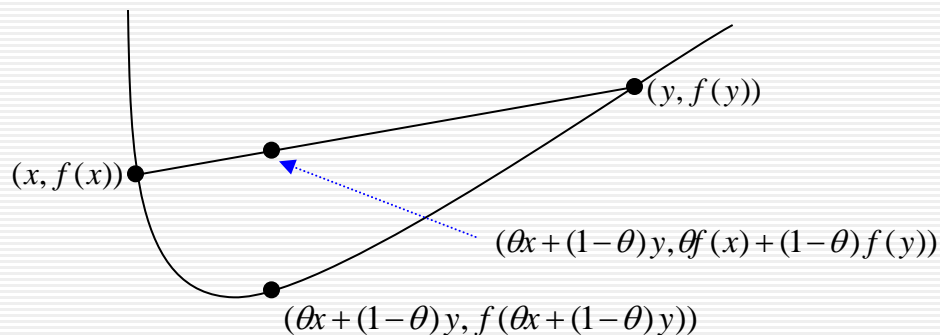
# Convex function (1)

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$f : \mathbb{R}^n \rightarrow \mathbb{R}$  is convex

if  $\text{dom } f$  is a convex set and for all  $x, y \in \text{dom } f$  and  $\theta$  with  $0 \leq \theta \leq 1$ ,

$$f(\theta x + (1-\theta)y) \leq \theta f(x) + (1-\theta)f(y)$$



$h : \mathbb{R}^n \rightarrow \mathbb{R}^m$  is affine

if it is a sum of a linear function and a constant, i.e.,

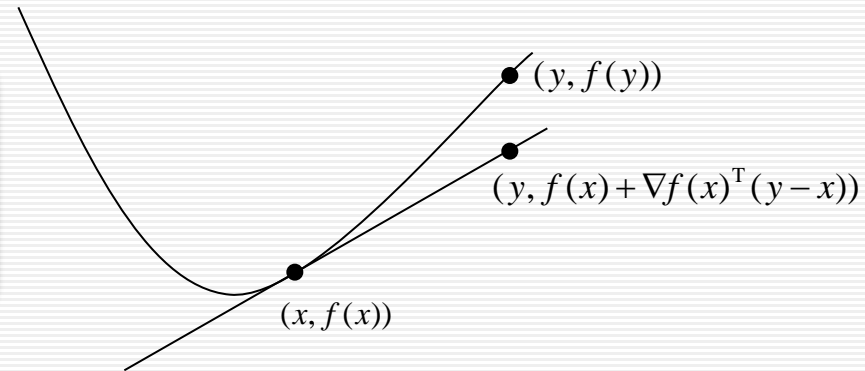
$$h(x) = Ax + b \quad \text{where } A \in \mathbb{R}^{m \times n} \text{ and } b \in \mathbb{R}^m$$

# Convex function (2)

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## First-order Condition

When  $f$  is differentiable,  
 $f$  is convex if  $\text{dom } f$  is a convex set,  
 $f(y) \geq f(x) + \nabla f(x)^T (y - x)$



## Second-order Condition

When  $f$  is twice differentiable,  
 $f$  is convex if  $\text{dom } f$  is a convex set and  
 $\nabla^2 f(x) \geq 0$  for all  $x \in \text{dom } f$

# Primal Problem & Dual Problem

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## ■ Primal Problem

$$\begin{aligned} &\text{minimize } f_0(x) \\ &\text{subject to } f_i(x) \leq 0, \quad i = 1, 2, \dots, m \\ &\quad \quad \quad h_i(x) = 0, \quad i = 1, 2, \dots, p \end{aligned}$$

## ■ Lagrangian $L: R^n \times R^m \times R^p \rightarrow R$

$$L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x)$$

— Lagrangian multipliers:  $\lambda = (\lambda_1, \lambda_2, \dots, \lambda_m)$   $\nu = (\nu_1, \nu_2, \dots, \nu_p)$

## ■ Lagrange Dual Function

$$g(\lambda, \nu) = \inf_{x \in D} L(x, \lambda, \nu)$$

—  $D$ : a domain of a decision vector  $x$

# Lower Bound Property (1)

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- Suppose  $\tilde{x}$  is a feasible point for the primal problem:

$$f_i(\tilde{x}) \leq 0 \text{ and } h_i(\tilde{x}) = 0$$

- Since  $\lambda \geq 0$ ,  $\sum_{i=1}^m \lambda_i f_i(\tilde{x}) + \sum_{i=1}^p \nu_i h_i(\tilde{x}) \leq 0$

$$\text{Then, } L(\tilde{x}, \lambda, \nu) = f_0(\tilde{x}) + \sum_{i=1}^m \lambda_i f_i(\tilde{x}) + \sum_{i=1}^p \nu_i h_i(\tilde{x}) \leq f_0(\tilde{x})$$

$$\text{Thus, } g(\lambda, \nu) = \inf_{x \in D} L(x, \lambda, \nu) \leq L(\tilde{x}, \lambda, \nu) \leq f_0(\tilde{x})$$

Since  $g(\lambda, \nu) \leq f_0(\tilde{x})$  holds for every feasible point  $\tilde{x}$ ,  $g(\lambda, \nu) \leq f_0(x^*)$

$$g(\lambda, \nu) \leq p^* \text{ for any } \lambda \geq 0 \text{ and any } \nu,$$

$p^*$  : optimal solution of the primal problem

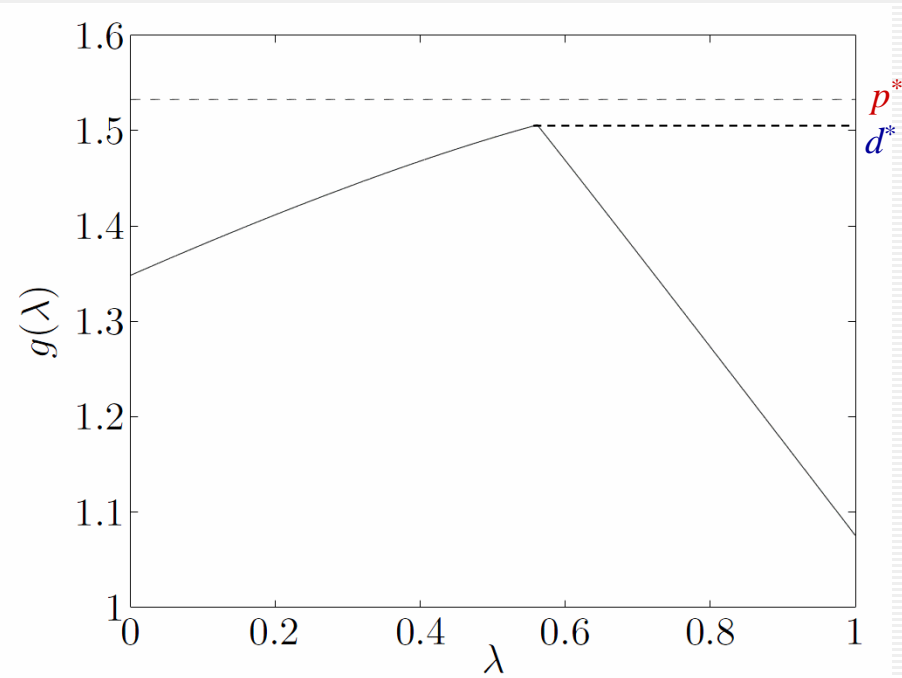
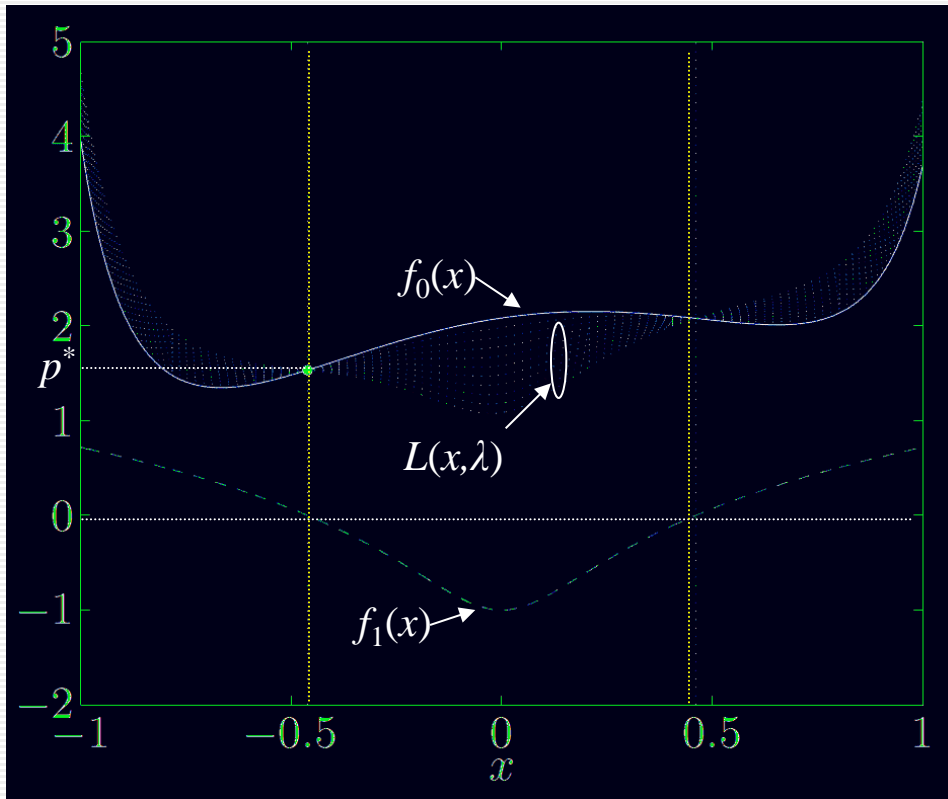
# Lower Bound Property (2)

Example :

$$\begin{aligned} & \text{minimize} && f_0(x) \\ & \text{subject to} && f_1(x) \leq 0 \end{aligned}$$



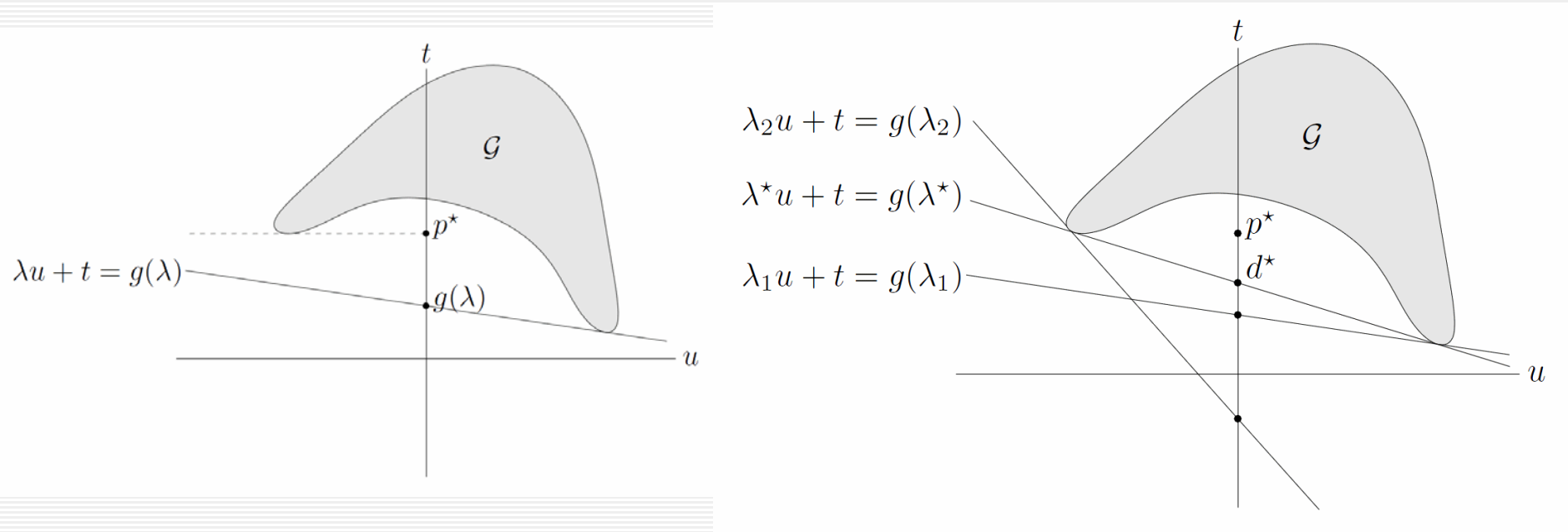
$$\begin{aligned} L(x, \lambda) &= f_0(x) + \lambda f_1(x) \\ g(\lambda) &= \inf_x L(x, \lambda) \quad \text{for } \lambda \geq 0 \end{aligned}$$



# Geometric Representation

$$\Omega = \{(f_1(x), f_2(x), \dots, f_m(x), h_1(x), h_2(x), \dots, h_p(x)), f_0(x) \in R^m \times R^p \times R \mid x \in D\}$$

$$p^* = \inf \{t \mid (u, v, t) \in \Omega, u \preceq 0, v = 0\}$$



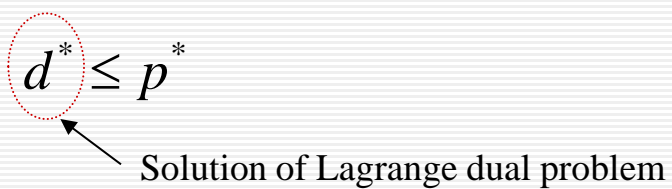


# Lagrange Dual Problem

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- What is the **best lower bound** from the Lagrange dual function?
- Lagrange Dual Problem

$$\begin{aligned} &\text{maximize } g(\lambda, \nu) \\ &\text{subject to } \lambda_i \geq 0, \quad i = 1, 2, \dots, m \end{aligned}$$

- This dual problem is a **convex optimization problem**, although the primal problem is not convex
- Thus, generally it is easier to solve than the primal problem.
- Duality
  - Weak Duality (always holds):  $d^* \leq p^*$   

  - **Strong Duality**:  $d^* = p^*$

# Strong Duality Condition

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- Convex Primal Problem

$$\begin{aligned} & \text{minimize } f_0(x) \\ & \text{subject to } f_i(x) \leq 0, \quad i = 1, 2, \dots, m \\ & \quad \quad \quad Ax = b \end{aligned}$$

- $f_0(x), f_1(x), \dots, f_m(x)$ : convex functions

- Slater's constraint qualification

$$\begin{aligned} & \text{There exists an } x \in \text{relint } D \text{ such that} \\ & \quad \quad \quad f_i(x) < 0, \quad i = 1, \dots, m, \\ & \quad \quad \quad Ax = b \end{aligned}$$

- $x$ : strictly feasible point

- When the primal problem is convex

- if Slater's condition holds, strong duality holds

# Strong Duality Condition

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- Primal Problem

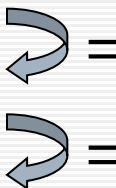
$$\begin{aligned} & \text{minimize } f_0(x) \\ & \text{subject to } f_i(x) \leq 0, \quad i = 1, 2, \dots, m \\ & \quad \quad \quad h_i(x) = 0, \quad i = 1, 2, \dots, p \end{aligned}$$

- Lagrange Dual Problem

$$\begin{aligned} & \text{maximize } g(\lambda, \nu) = \inf_x \left( f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x) \right) \\ & \text{subject to } \lambda_i \geq 0 \quad \text{for all } i \end{aligned}$$

- Strong Duality:  $f_0(x^*) = g(\lambda^*, \nu^*)$

- $x^*$ : primal optimal point,  $(\lambda^*, \nu^*)$ : dual optimal point

- $$\begin{aligned} g(\lambda^*, \nu^*) &= \inf_x \left( f_0(x) + \sum_{i=1}^m \lambda_i^* f_i(x) + \sum_{i=1}^p \nu_i^* h_i(x) \right) \\ &\leq f_0(x^*) + \sum_{i=1}^m \lambda_i^* f_i(x^*) + \sum_{i=1}^p \nu_i^* h_i(x^*) \\ &\leq f_0(x^*) \end{aligned}$$


# KKT Optimality Conditions

- $f_0, f_1, \dots, f_m, h_1, \dots, h_p$  should be differentiable.

- $g(\lambda^*, \nu^*)$   
$$= \inf_x \left( f_0(x) + \sum_{i=1}^m \lambda_i^* f_i(x) + \sum_{i=1}^p \nu_i^* h_i(x) \right) \Rightarrow \nabla f_0(x^*) + \sum_{i=1}^m \lambda_i^* \nabla f_i(x^*) + \sum_{i=1}^p \nu_i^* \nabla h_i(x^*) = 0$$
$$= f_0(x^*) + \sum_{i=1}^m \lambda_i^* f_i(x^*) + \sum_{i=1}^p \nu_i^* h_i(x^*)$$
$$= f_0(x^*) \Rightarrow \sum_{i=1}^p \lambda_i^* f_i(x^*) = 0 \Rightarrow \lambda_i^* f_i(x^*) = 0$$

- **KKT conditions**

$$f_i(x^*) \leq 0, \quad i = 1, \dots, m$$

$$h_i(x^*) = 0, \quad i = 1, \dots, p$$

$$\lambda_i^* \geq 0, \quad i = 1, \dots, m$$

$$\lambda_i^* f_i(x^*) = 0, \quad i = 1, \dots, m,$$

$$\nabla f_0(x^*) + \sum_{i=1}^m \lambda_i^* \nabla f_i(x^*) + \sum_{i=1}^p \nu_i^* \nabla h_i(x^*) = 0,$$

# Example: Water Filling (1)

## ■ Problem Description

- The problem is to allocate power to a set of  $n$  subchannels so that the total data rate is maximized.
- The total power allocated to  $n$  subchannels is fixed to  $P_T$ .

## ■ Problem Formulation

- $P_i$ : the power of subchannel  $i$ ,
- $x_i$ : the normalized power of subchannel  $i$ :  $P_i/P_T$
- $\gamma_i$ : SNR of subchannel  $i$  when the transmission power is  $P_T$
- Total data rate:  $\sum_{i=1}^n B \log_2 \left( 1 + \frac{P_i}{P_T} \gamma_i \right) = \sum_{i=1}^n B \log_2 (1 + x_i \gamma_i)$

maximize  $\sum_{i=1}^n B \log_2 (1 + x_i \gamma_i)$

subject to  $x_i \geq 0, i = 1, \dots, n$

$\sum_{i=1}^n x_i = 1$

→  $B \sum_{i=1}^n \log_2 (1/\gamma_i + x_i) - B \sum_{i=1}^n \log_2 (\gamma_i)$

→ minimize  $-\sum_{i=1}^n \ln(1/\gamma_i + x_i)$

# Example: Water Filling (2)

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- Primal Problem (Convex Problem)

$$\begin{array}{ll} \text{minimize} & -\sum_{i=1}^n \ln(1/\gamma_i + x_i) \\ \text{subject to} & \mathbf{x} \geq 0, \quad \mathbf{1}^T \mathbf{x} = 1 \end{array}$$

- Lagrangian

$$-\sum_{i=1}^n \ln(1/\gamma_i + x_i) - \sum_{i=1}^n \lambda_i x_i + \nu(\sum_{i=1}^n x_i - 1)$$

# Example: Water Filling (3)

- KKT condition:

- $\mathbf{x}^*$ : primal optimal point,  $(\lambda^*, \nu^*)$ : dual optimal point

- $\mathbf{x}^* \geq 0, \mathbf{1}^T \mathbf{x}^* = 1, \lambda^* \geq 0,$

- $\lambda_i^* x_i^* = 0, \quad i = 1, \dots, n,$

- $\frac{-1}{1/\gamma_i + x_i^*} - \lambda_i^* + \nu^* = 0, \quad i = 1, \dots, n.$



- $\mathbf{x}^* \geq 0, \mathbf{1}^T \mathbf{x}^* = 1,$

- $\nu^* \geq \frac{1}{1/\gamma_i + x_i^*}, \quad i = 1, \dots, n,$

- $x_i^* (\nu^* - 1/(1/\gamma_i + x_i^*)) = 0, \quad i = 1, \dots, n,$

- Solution

- if  $\nu^* < \gamma_i, x_i^* > 0 \Rightarrow \nu^* = \frac{1}{1/\gamma_i + x_i^*} \Rightarrow x_i^* = \frac{1}{\nu^*} - \frac{1}{\gamma_i}$

- if  $\nu^* \geq \gamma_i, x_i^* \leq 0 \Rightarrow x_i^* = 0$

$$x_i^* = \begin{cases} 1/\nu^* - 1/\gamma_i & \text{if } \gamma_i > \nu^* \\ 0 & \text{if } \gamma_i \leq \nu^* \end{cases}$$

# Convex Optimization

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1. Formulate a primal problem
2. Derive the Lagrangian
3. Construct the Lagrange dual function
  - Concave
  - Not necessarily differentiable
4. Solve the Lagrange dual problem (convex optimization)
  - There are various methods
  - Gradient (subgradient) method
    - differential: projected gradient method
    - not differential: projected subgradient method
5. Prove the strong duality



# Gradient/Subgradient Method

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- Consider the following general concave maximization:

$$\begin{array}{l} \text{maximize } g(\lambda) \\ \text{subject to } \lambda \in \chi \end{array}$$

(where  $\chi$  is a convex set)

- The methods generate a sequence of feasible points  $\{\lambda(k)\}$  as

$$\lambda(k+1) = [\lambda(k) + \alpha(k) \times s(k)]_{\chi}$$

- $s(k)$ : a gradient(subgradient) of  $g$  at the point of  $\lambda(k)$ 
  - It depends on whether  $g$  is differentiable or not.
- $[\cdot]_{\chi}$ : the projection onto the feasible set  $\chi$
- $\alpha(k)$ : a positive step size at the  $k$ th iteration

# Subgradient Method

- Subgradient property

$$g(y) \leq g(x) + s(x) \times (y - x)$$

- $s(x)$  is a subgradient of  $g$  at  $x$

- Simple example

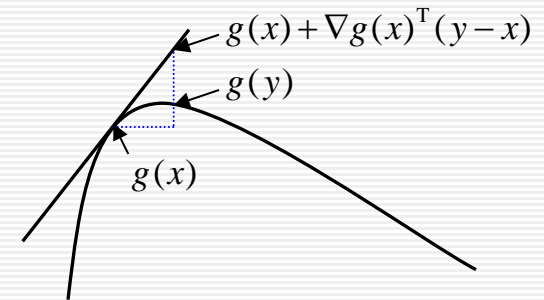
$$g(\lambda) = \min_{x \in X} \{f_0(x) + \lambda f_1(x)\}$$

$$x(\lambda) = \arg \min_{x \in X} \{f_0(x) + \lambda f_1(x)\}$$

$$\begin{aligned} g(\mu) &= \min_{x \in X} \{f_0(x) + \mu f_1(x)\} \\ &\leq f_0(x(\lambda)) + \mu f_1(x(\lambda)) \\ &= f_0(x(\lambda)) + \lambda f_1(x(\lambda)) + (\mu - \lambda) f_1(x(\lambda)) \\ &= g(\lambda) + (\mu - \lambda) f_1(x(\lambda)) \end{aligned}$$

$f_1(x(\lambda))$  is a subgradient of  $g$  at  $\lambda$

Gradient property



# Subgradient Method for Solving Lagrange Dual Problem

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- For a given  $\lambda, \nu$

$$\begin{aligned}x(\lambda, \nu) &= \arg \min_{x \in D} L(x, \lambda, \nu) \\ &= \arg \min_{x \in D} \left\{ f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x) \right\}\end{aligned}$$

- Subgradients of  $g$  at  $(\lambda, \nu)$ :  $f_i(x(\lambda, \nu)), h_i(x(\lambda, \nu))$

- For obtaining an optimal  $\lambda^*, \nu^*$

$$\left. \begin{aligned}\lambda_i^{(k+1)} &= \lambda_i^{(k)} + \delta^{(k)} \times f_i(x(\lambda^{(k)}, \nu^{(k)})) \\ \nu_i^{(k+1)} &= \nu_i^{(k)} + \delta^{(k)} \times h_i(x(\lambda^{(k)}, \nu^{(k)}))\end{aligned} \right\} \Rightarrow \text{converge to } \lambda^*, \nu^*$$

$\delta^{(k)}$  : positive step size at the  $k$ th iteration

$$\delta^{(k)} > 0, \sum_{k=1}^{\infty} \delta^{(k)} = \infty, \sum_{k=1}^{\infty} (\delta^{(k)})^2 < \infty \text{ (ex. } \delta^{(k)} = c/k)$$