

# Artificial Intelligence

## Chapter 8

# Uninformed Search

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Biointelligence Lab  
School of Computer Sci. & Eng.  
Seoul National University

# Outline

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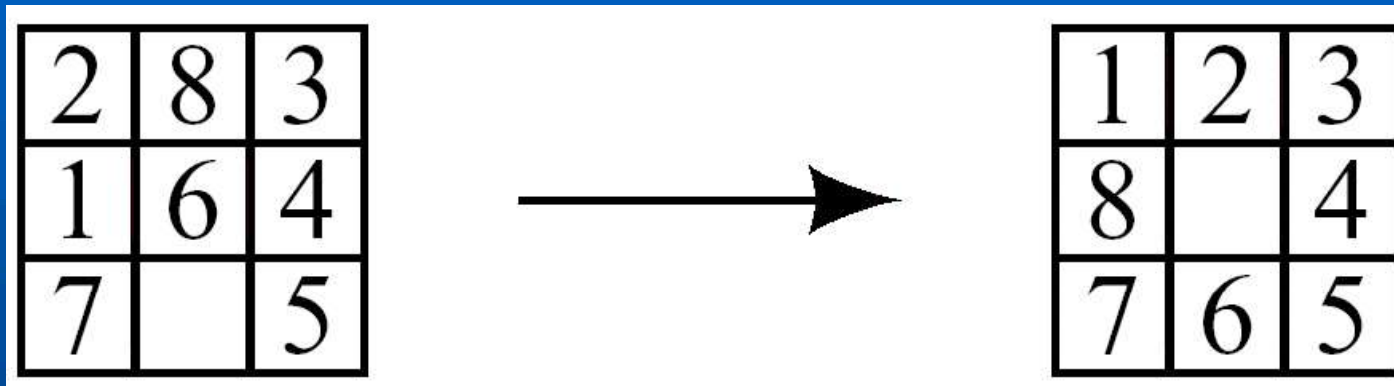
- Search Space Graphs
- Depth-First Search
- Breadth-First Search
- Iterative Deepening

# 1. Formulating the State Space

- For huge search space we need,
  - ◆ Careful formulation
  - ◆ Implicit representation of large search graphs
  - ◆ Efficient search method

# 1. Formulating the State Space (Cont'd)

- e.g.) 8-puzzle problem



- ◆ state description
  - 3-by-3 array: each cell contains one of 1-8 or blank symbol
- ◆ two state transition descriptions
  - 8×4 moves: one of 1-8 numbers moves up, down, right, or left
  - 4 moves: one blank symbol moves up, down, right, or left

# 1. Formulating the State Space (Cont'd)

- ◆ The number of nodes in the state-space graph:
  - $9!$  (= 362,880)
- ◆ State space for 8-puzzle is
  - Divided into two separate graphs : not reachable from each other

## 2. Components of Implicit State-Space Graphs

- 3 basic components to an implicit representation of a state-space graph
  1. Description of start node
  2. *Actions*: Functions of state transformation
  3. *Goal condition*: *true-false* valued function
- 2 classes of search process
  1. *Uninformed* search: no problem specific information
  2. *Heuristic* search: existence of problem-specific information

# 3. Breadth-First Search

- Procedure

1. Apply all possible operators (*successor function*) to the start node.
2. Apply all possible operators to all the direct successors of the start node.
3. Apply all possible operators to their successors till goal node found.
  - ♠ *Expanding* : applying successor function to a node

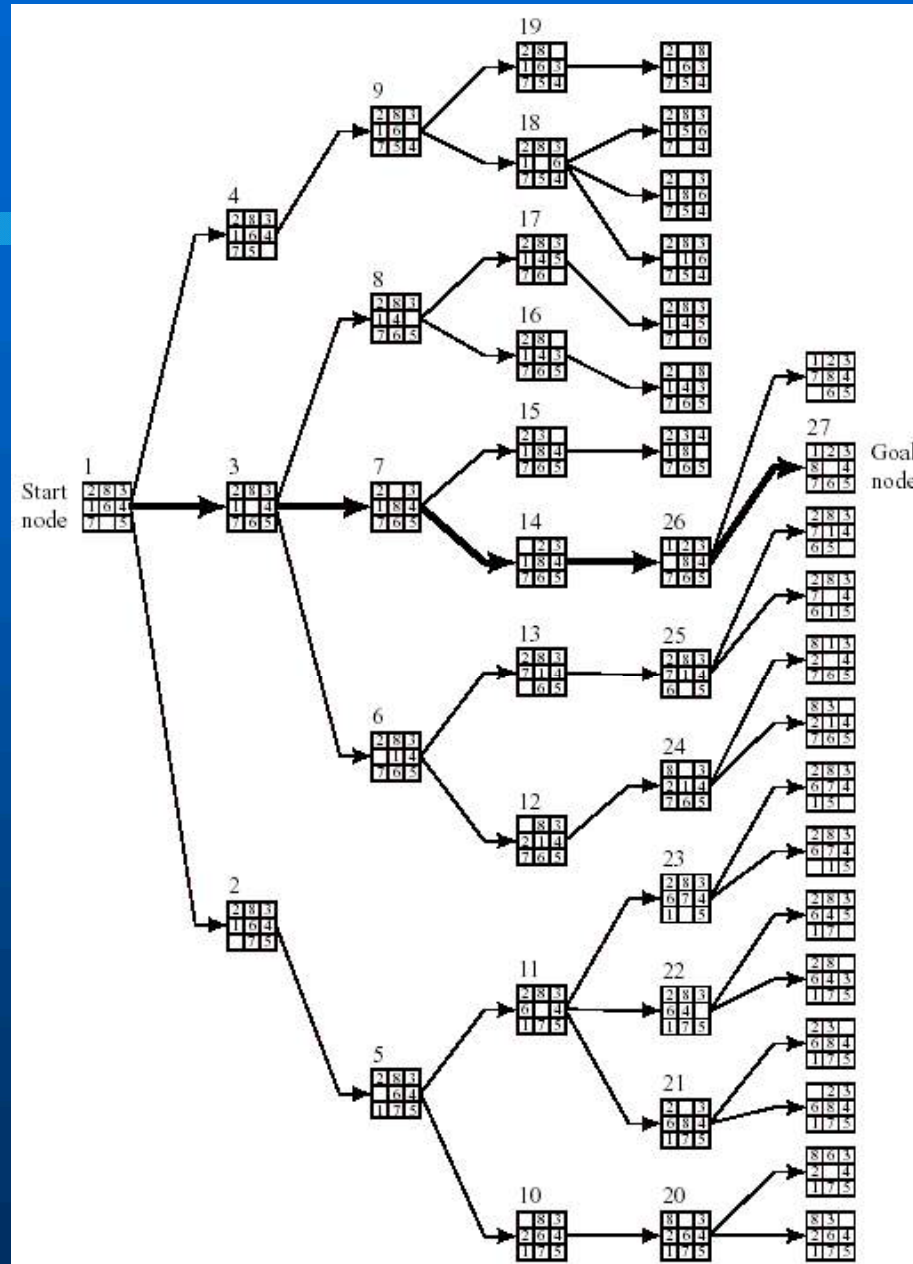


Figure 8.2 Breadth-First Search of the Eight-Puzzle



### 3. Breadth-First Search (Cont'd)

- Advantage
  - ◆ Finds the path of minimal length to the goal.
- Disadvantage
  - ◆ Requires the generation and storage of a tree whose size is exponential the the depth of the shallowest goal node
- *Uniform-cost* search [Dijkstra 1959]
  - ◆ Expansion by *equal cost* rather than equal depth

# 4. Depth-First or Backtracking Search

- Procedure

- ◆ Generates the successor of a node just one at a time.
- ◆ Trace is left at each node to indicate that additional operators can be applied there if needed.
- ◆ At each node a decision must be made about which operator to apply first, which next, and so on.
- ◆ Repeats this process until the *depth bound*.
- ◆ *chronological Backtrack* when search depth is depth bound.

# 4. Depth-First or Backtracking Search (Cont'd)

- 8-puzzle example
  - ◆ Depth bound: 5
  - ◆ Operator order: left → up → right → down

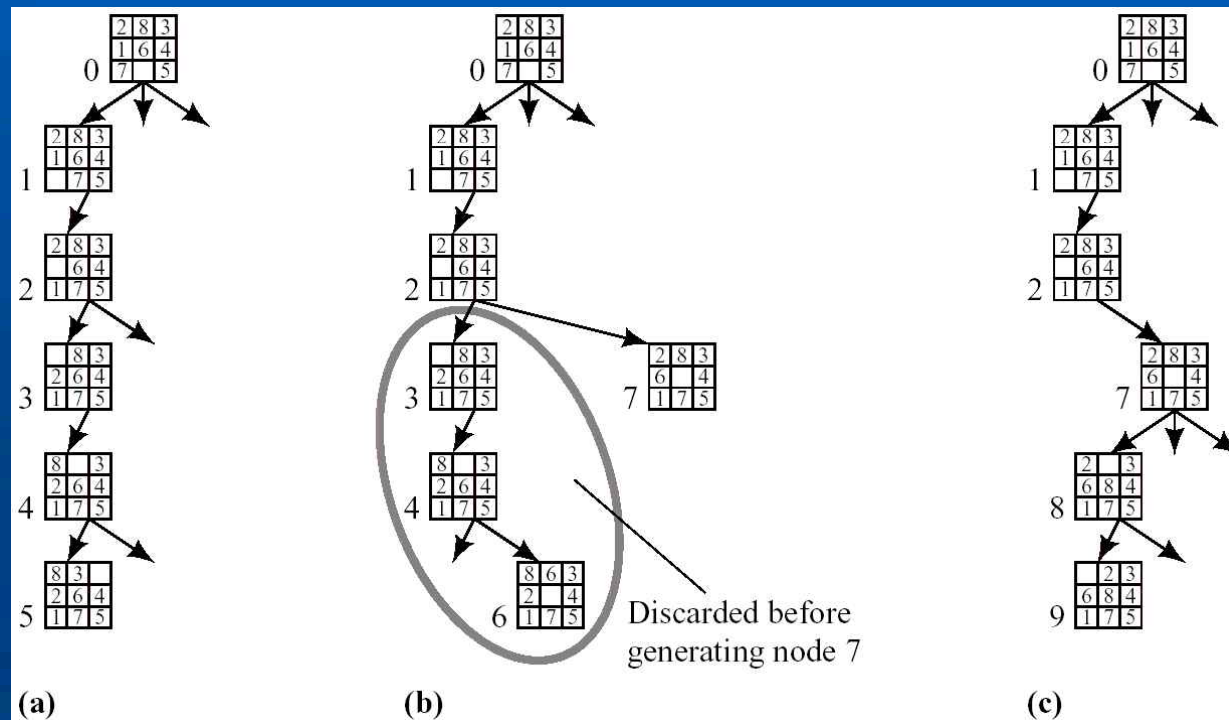
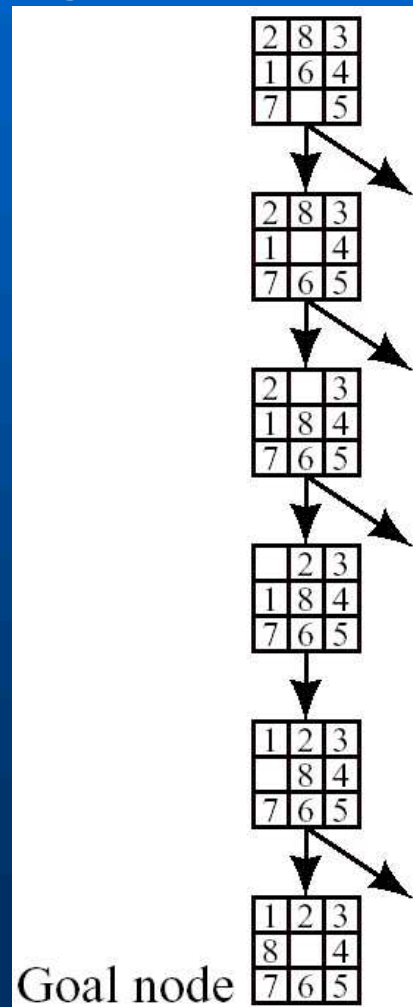


Figure 8.3 Generation of the First Few Nodes in a Depth-First Search

# 4. Depth-First or Backtracking Search (Cont'd)

- ◆ The graph when the goal is reached in depth-first search



# 4. Depth-First or Backtracking Search (Cont'd)

- Advantage

- ◆ Low memory size: linear in the depth bound
  - saves only that part of the search tree consisting of the path currently being explored plus traces

- Disadvantage

- ◆ No guarantee for the minimal state length to goal state
- ◆ The possibility of having to explore a large part of the search space

# 5. Iterative Deepening

- Advantage
  - ◆ Linear memory requirements of depth-first search
  - ◆ Guarantee for goal node of minimal depth
- Procedure
  - ◆ Successive depth-first searches are conducted – each with depth bounds increasing by 1

# 5. Iterative Deepening (Cont'd)

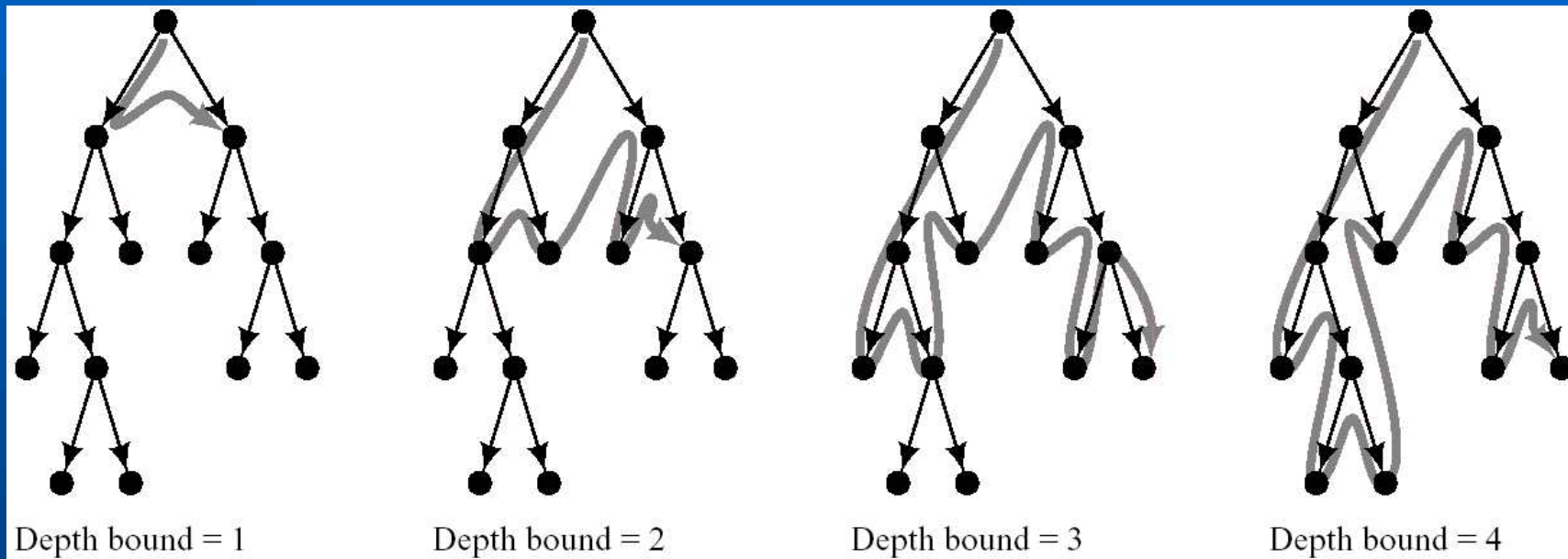


Figure 8.5 Stages in Iterative-Deepening Search

## 5. Iterative Deepening (Cont'd)

- The number of nodes
  - ◆ In case of breadth-first search

$$N_{\text{bf}} = 1 + b + b^2 + \dots + b^d = \frac{b^{d+1} - 1}{b - 1} \quad (b : \text{branching factor}, d : \text{depth})$$

- ◆ In case of iterative deepening search

$$\begin{aligned} N_{\text{df}_j} &= \frac{b^{j+1} - 1}{b - 1} : \text{number of nodes expanded down to level } j \\ N_{\text{id}} &= \sum_{j=0}^d \frac{b^{j+1} - 1}{b - 1} \\ &= \frac{1}{b - 1} \left[ b \left( \sum_{j=0}^d b^j \right) - \sum_{j=0}^d 1 \right] = \frac{1}{b - 1} \left[ b \left( \frac{b^{d+1} - 1}{b - 1} \right) - (d + 1) \right] \\ &= \frac{b^{d+2} - 2b - bd + d + 1}{(b - 1)^2} \end{aligned}$$



## 5. Iterative Deepening (Cont'd)

- ◆ For large  $d$  the ratio  $N_{id}/N_{df}$  is  $b/(b-1)$
- ◆ For a branching factor of 10 and deep goals, 11% more nodes expansion in iterative-deepening search than breadth-first search
- ◆ Related technique *iterative broadening* is useful when there are many goal nodes

## 6. Additional Readings and Discussion

- Various improvements in chronological backtracking
  - ◆ *Dependency-directed backtracking* [Stallman & Sussman 1977]
  - ◆ *Backjumping* [Gaschnig 1979]
  - ◆ *Dynamic backtracking* [Ginsberg 1993]