



Loop Invariant Computation and Code Motion

- * Loop-invariant computation
- * Algorithms for code motion
- * Summary: 13.1.2 (global CSE), 13.2

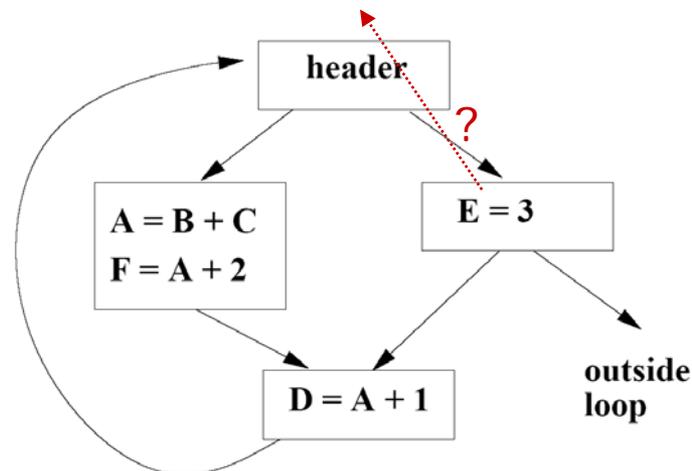
Loop-Invariant Computation and Code Motion

* Loop invariant computation

- * A computation whose value does not change as long as control stays within the loop

* Loop invariant code motion (LICM)

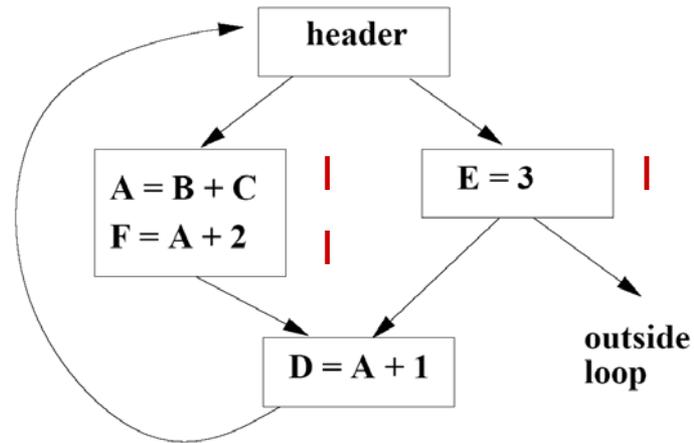
- * Move a loop invariant statement within a loop to the pre-header of the loop





Example

- * Which statement is a loop-invariant?
- * Which statement can we move to the preheader?





LICM Algorithm

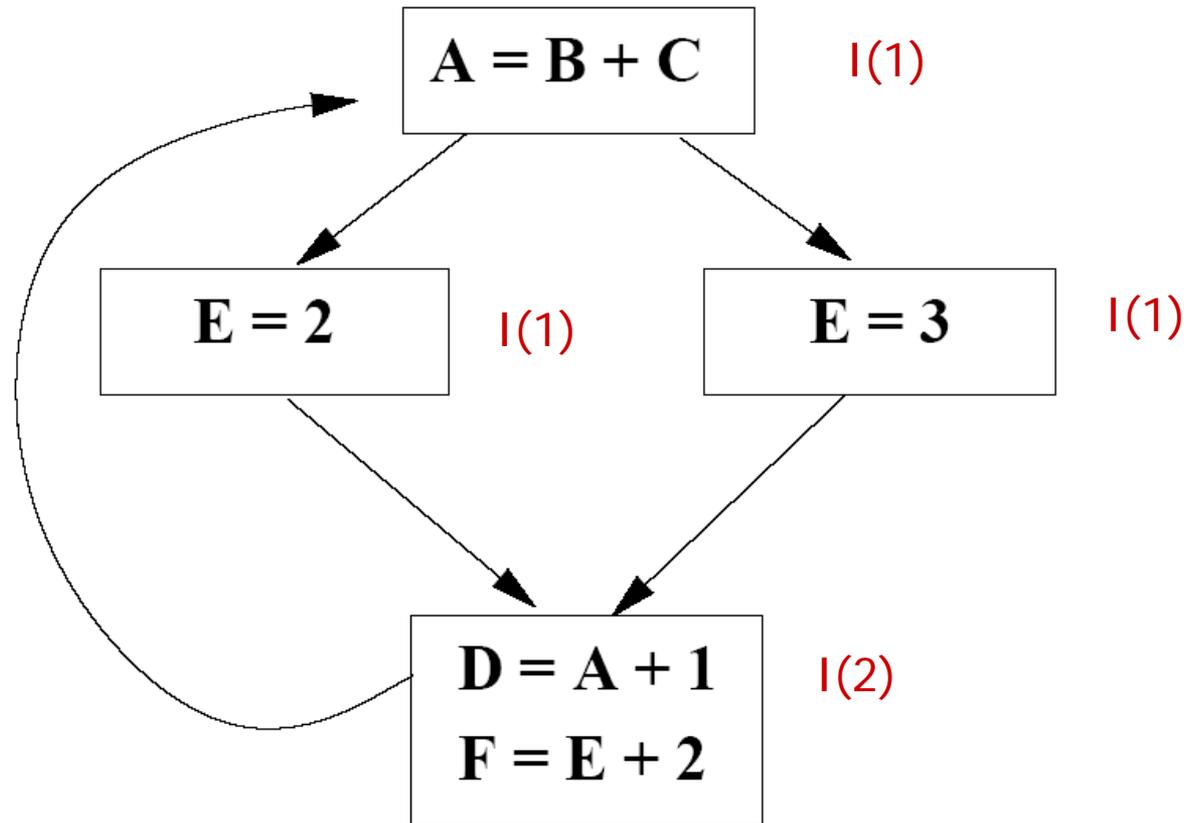
- * Observations
 - * Loop invariant: operands are defined outside loop or are defined by loop invariants
 - * Code motion: not all invariant statements can be moved to the pre-header
- * Algorithm
 - * Detect loop invariant computations
 - * Check conditions for code motion
 - * Code transformation



Detecting loop invariant computation

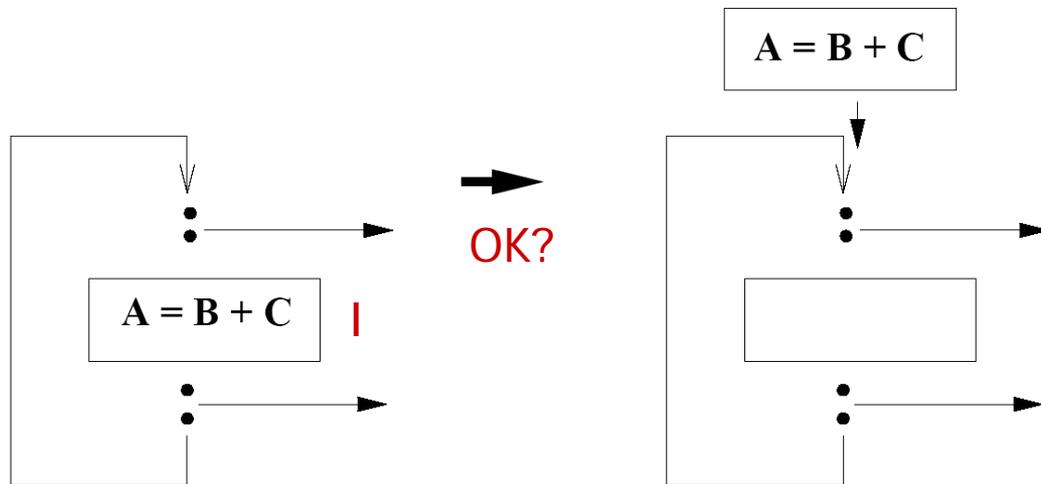
- * Compute **reaching definitions**
- * Repeat: **mark $A=B+C$ as invariant if**
 - * All reaching definitions of B are **outside** of the loop, or there is exactly **one reaching definition** for B and it is **from a loop-invariant** statement inside the loop
 - * Check similarly for C
- * **Until** no changes to the set of loop-invariant statements occur

Example



Conditions for Code Motion

- * Correctness: movement does not change the program semantics
- * Performance: code should not be slowed down



Need information on

- * Control flow of the loop: **dominate all the exits**
- * Other definitions: **no other definitions of A**
- * Other uses: **all uses of A are dominated by block b**



Code Motion Algorithm

Given a set of nodes in a loop

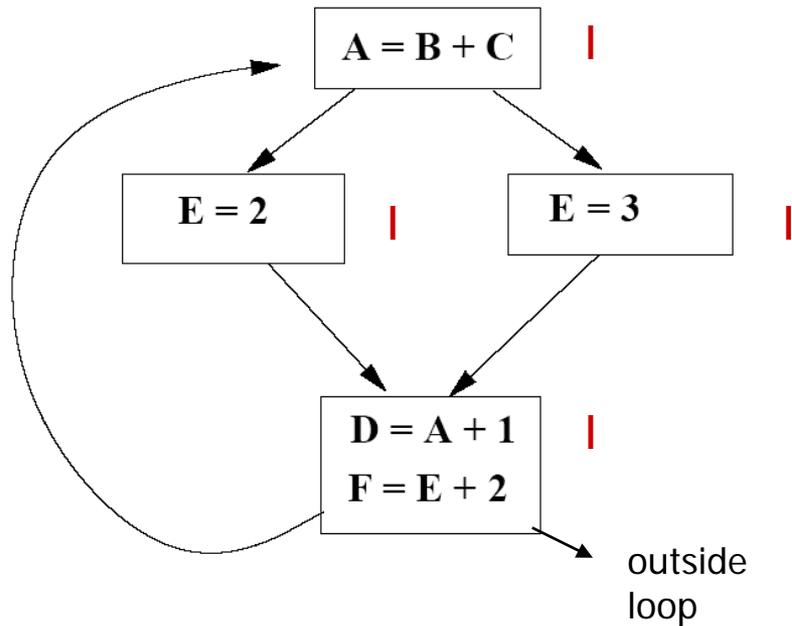
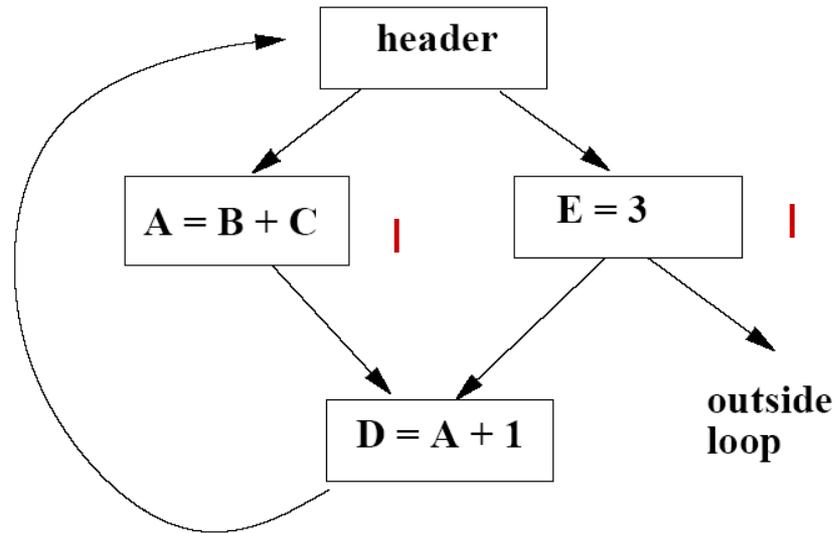
- * Compute reaching definitions
- * Compute loop invariant computation
- * Compute dominators
- * Find the exits of the loop, which are nodes whose successors are located outside loop



Code Motion Details

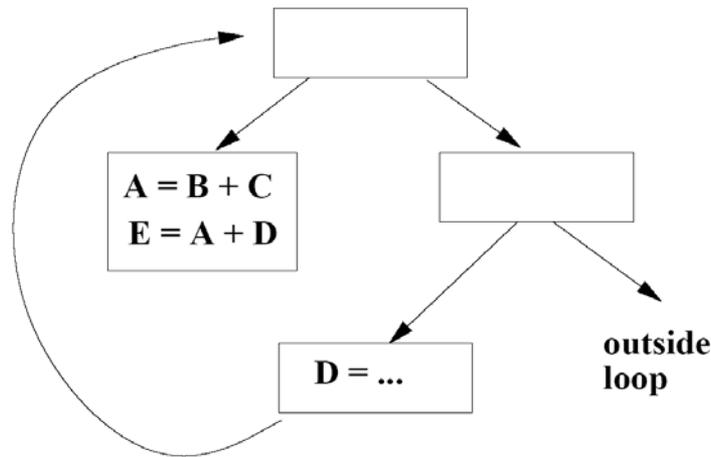
- * Candidate statement for code motion
 - * Loop invariant
 - * Located in a block that dominates all the exits of the loop
 - * Assign to a variable not assigned to elsewhere in the loop
 - * Located in a block that dominate all blocks in the loop that contain the use of the variable
- * Visit blocks in a reverse-postorder
 - * Move the candidate statement to pre-header if all the invariant operations it depends on have been moved

Examples



More Aggressive Optimizations

- * Gamble on: most loops get executed
 - * Can we relax constraint of dominating all exits?
 - * If liveness check and exception check is satisfied





Landing Pads

Code for most loops is generated in a **test-repeat-until** form to simplify dominance

Landing pads:

While p DO s



```
if (p) {  
    preheader  
    repeat  
        s  
    until not p  
}
```