Chapter 1 Introduction

Why study logic design?

- it is the implementation basis for all modern computing devices
 - building large things from small components
 - provide a model of how a computer works
- More important reasons
 - the inherent parallelism in hardware is often our first exposure to parallel computation
 - it offers an interesting counterpoint to software design and is therefore useful in furthering our understanding of computation, in general

What will we learn in this class? (1/2)

- The basics of logic design
 - Boolean algebra, logic minimization, state, timing, CAD tools
- The concept of state in digital systems
 - analogous to variables and program counters in software systems

What will we learn in this class? (2/2)

- How to specify/simulate/compile/realize our designs
 - hardware description languages (HDLs)
 - tools to simulate the workings of our designs
 - Iogic compilers to synthesize the hardware blocks of our designs
 - mapping onto programmable hardware
- Contrast with software design
 - sequential and parallel implementations
 - □ specify algorithm as well as computing/storage resources it will use

Applications of logic design

- Conventional computer design
 - CPUs, busses, peripherals
- Networking and communications
 - □ phones, modems, routers
- Embedded products
 - in cars, toys, appliances, entertainment devices
- Scientific equipment
 - testing, sensing, reporting
- The world of computing is much bigger than just PCs!

Global processor/controller market



PC CPUs constitute less than 2% of the market

What is logic design? (1/2)

- Digital hardware consists of components
- Components or building blocks
 - Switches built from semiconductor transistors
 - Most basic element
 - Higher level circuits such as logic gates and memories
- A logic designer should choose the right component to solve logic design problems
- Constraints: size, cost, performance, power consumption
 - Cost vs. size

* A circuit is an interconnected collection of switches

What is logic design? (2/2)

- Each component has
 - a set of input wires
 - a set of output wires
 - Each wire is set to some analog voltage value
 - But will be interpreted as either 1 or 0 (digital abstraction)
- Transistors react to the voltage levels on the input wires
 - Switch their state and cause a change in output wires
 - At macro scale, a component that contains transistors reacts to input voltage values
- Depending on the way a circuit reacts to the input voltages
 - Combinational logic circuits
 - Sequential logic circuits

What is digital hardware?

- Collection of devices that sense and/or control wires, which carry a digital value (i.e., a physical quantity that can be interpreted as a "0" or "1")
 - example: digital logic where voltage < 0.8v is a "0" and > 2.0v is a "1"
- Primitive digital hardware devices
 - logic computation devices (sense and drive)
 - are two wires both "1" make another be "1" (AND)
 - is at least one of two wires "1" make another be "1" (OR)
 - is a wire "1" then make another be "0" (NOT)
 - memory devices (store)
 - store a value
 - recall a previously stored value



Computation: abstract vs. implementation

- Computation has been a mental exercise (paper, programs)
- This class is about physically implementing computation using physical devices that use voltages to represent logical values
- Basic units of computation are:

representation:	"0", "1" on a wire set of wires (e.g., binary integer)
assignment:	x = y
data operations:	x + y - 5
control:	
sequential statements:	A; B; C
conditionals:	if x == 1 then y
loops:	for (i = 1 ; i == 10, i++)
procedures:	A; proc(); B;

 We will study how each of these are implemented in hardware and composed into computational structures

Switches: basic element of physical implementations

Implementing a simple circuit



close switch (if A is "1" or asserted) and turn on light bulb (Z)



open switch (if A is "0" or unasserted) and turn off light bulb (Z)

 $Z \equiv A$

Z and A are equivalent boolean variables

Switches (cont'd)

Compose switches into more complex ones (Boolean functions):



Switching networks

- Switch settings
 - determine whether or not a conducting path exists to light the light bulb
- To build larger computations
 - use a light bulb (output of the network) to set other switches (inputs to another network).
- Connect together switching networks
 - to construct larger switching networks, i.e., there is a way to connect outputs of one network to the inputs of the next.



A mechanical switch



• A semiconductor switch or transistor

Transistor networks

- Modern digital systems are designed in CMOS technology
 - MOS stands for Metal-Oxide on Semiconductor
 - C is for complementary because there are both normally-open and normally-closed switches: nMOS and pMOS
- MOS transistors act as voltage-controlled switches

* CMOS: complementary metal-oxide semiconductor

n-channel (or n-type) MOS (nMOS) circuit

- Three terminals: source-gate-drain (or S-G-D for short)
- three layers: polysilicon (used to be metal) SiO2 substrate



- If G is at positive voltage, electrons in the substrate will move toward G terminal, which sets up a channel between S and D
 - And D is at high voltage, current will flow from drain to source
- Metal is replaced by polysilicon which is more adhesive

* n+: heavily doped n-type semiconductor

p-channel (or p-type) MOS (pMOS) circuit

- Three terminals: source-gate-drain (or S-G-D for short)
- Same principle, but reverse doping and voltage
 - Source (Vss) is positive with regard to drain (Vdd)
- Bubble indicates the inverted behavior



- If G is at positive voltage, the current does not flow
- If G is at ground level, the current flows

MOS transistors

- MOS transistors have three terminals: drain, gate, and source
 - they act as switches in the following way: if the voltage on the gate terminal is (some amount) higher/lower than the source terminal, then a conducting path will be established between the drain and source terminals



n-channel open when voltage at G is low closed when voltage at G is high



p-channel closed when voltage at G is low open when voltage at G is high



- A simple component is made up of two transistors
 - X: input
 - Y: output
 - What is this function?
- In CMOS circuits, pMos and nMOS are used in pair





what is the relationship between x, y and z1/z2?

Х	у	z1	z2
0 volts	0 volts		
0 volts	3 volts		
3 volts	0 volts		
3 volts	3 volts		

Three input NAND gate

- Y pulls low if ALL inputs are 1
- Y pulls high if ANY input is 0



- In general, the more inputs, the more transistors (TRs)
- In CMOS, a variable requires a pair of TRs

Digital vs. analog

- Convenient to think of digital systems as having only discrete, digital, input/output values
- In reality, real electronic components exhibit continuous, analog behavior
 - Why do we make the digital abstraction anyway?
 - switches operate this way
 - easier to think about a small number of discrete values
 - Quantization error, though
- Why does it work?
 - does not propagate small errors in values
 - always resets to 0 or 1

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Combinational logic symbols

- Common combinational logic systems have standard symbols called logic gates
 - Buffer, NOT



Combinational vs. sequential digital circuits

A simple model of a digital system is a unit with inputs and outputs:



- Combinational means "memory-less"
 - a digital circuit is combinational if its output values only depend on its (current) input values

Combinational vs. sequential digital circuits



(a) Combinational

(b) Sequential

Output = f(In)

Output = f(In, Previous In)

Sequential logic

Sequential systems

 exhibit behaviors (output values) that depend not only on the current input values, but also on previous input values

Example of combinational and sequential logic

- Combinational:
 - input A, B
 - wait for clock edge
 - observe C
 - wait for another clock edge
 - observe C again: will stay the same
- Sequential:
 - input A, B
 - wait for clock edge
 - observe C
 - wait for another clock edge
 - observe C again: may be different



Abstractions

- Some we've seen already
 - digital interpretation of analog values
 - transistors as switches
 - switches as logic gates
 - use of a clock to realize a synchronous sequential circuit
- Some others we will see
 - truth tables and Boolean algebra to represent combinational logic
 - encoding of signals with more than two logical values into binary form
 - state diagrams to represent sequential logic
 - hardware description languages to represent digital logic
 - waveforms to represent temporal behavior

An example

- Calendar subsystem: number of days in a month (to control watch display)
 - Combinational logic
 - used in controlling the display of a wrist-watch LCD screen
 - □ inputs: month, leap year flag
 - outputs: number of days

```
Implementation in software
integer number_of_days ( month, leap_year_flag)
{
  switch (month) {
    case 1: return (31);
    case 2: if (leap year flag == 1) then return (29)
                                     else return (28);
    case 3: return (31);
    . . .
                                       P.16
    case 12: return (31);
    default: return (0);
```

Implementation as a combinational digital system

- Encoding:
 - how many bits for each input/output?
 - binary number for month
 - four wires for 28, 29, 30, and 31
- Behavior:
 - combinational
 - truth table specification



					care
month	leap	d28	d29	d30	d31
0000		_	_	_	<u> </u>
0001	_	0	0	0	1
0010	0	1	0	0	0
0010	1	0	1	0	0
0011	—	0	0	0	1
0100	_	0	0	1	0
0101	—	0	0	0	1
0110	—	0	0	1	0
0111	_	0	0	0	1
1000	—	0	0	0	1
1001	—	0	0	1	0
1010	_	0	0	0	1
1011	_	0	0	1	0
1100	—	0	0	0	1
1101	—	-	—	_	—
111–	_	-	—	—	—

Don't



Combinational example (cont'd)

- d28 = m8'•m4'•m2•m1'•leap'
- d29 = m8'•m4'•m2•m1'•leap
- d30 = (m8'•m4•m2'•m1') + (m8'•m4•m2•m1') + (m8•m4'•m2'•m1) + (m8•m4'•m2•m1) = (m8'•m4•m1') + (m8•m4'•m1)
- d31 = (m8'•m4'•m2'•m1) + (m8'•m4'•m2•m1) + (m8'•m4•m2'•m1) + (m8'•m4•m2•m1) + (m8•m4'•m2'•m1') + (m8•m4'•m2•m1') + (m8•m4•m2'•m1')



Combinational example (cont'd)

- d28 = m8'•m4'•m2•m1'•leap'
- d29 = m8'•m4'•m2•m1'•leap
- d30 = (m8'•m4•m2'•m1') + (m8'•m4•m2•m1') + (m8•m4'•m2'•m1) + (m8•m4'•m2•m1) + (m8•m4'•m2•m1)
- d31 = (m8'•m4'•m2'•m1) + (m8'•m4'•m2•m1) + (m8'•m4•m2'•m1) + (m8'•m4•m2•m1) + (m8•m4'•m2'•m1') + (m8•m4'•m2•m1') + (m8•m4•m2'•m1')



Another example (Door combination lock)

- punch in 3 values in sequence and the door opens; if there is an error the lock must be reset; once the door opens the lock must be reset
 - Sequential logic
 - □ inputs: sequence of input values, reset
 - Numeric number: 4 wires
 - outputs: door open/close
 - memory: must remember combination

or always have it available as an input

```
Implementation in software
```

```
integer combination_lock ( ) {
    integer v1, v2, v3;
    integer error = 0;
    static integer c[3] = 3, 4, 2;
```

```
while (!new_value( ));
v1 = read_value( );
if (v1 != c[1]) then error = 1;
```

```
while (!new_value( ));
v2 = read_value( );
if (v2 != c[2]) then error = 1;
```

Array index starts from 1

```
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```

v3 = read_value();
if (v2 != c[3]) then error = 1;

while (!new_value());

}

```
if (error == 1) then return(0); else return (1);
```

Implementation as a sequential digital system

- Encoding:
 - how many bits per input value?
 - how many values in sequence?
 - how do we know a new input value is entered?
 - how do we represent the states of the system?

Behavior:

- clock wire tells us when it's ok to look at inputs (i.e., they have settled after change)
- sequential: sequence of values must be entered
- sequential: remember if an error occurred
- finite-state specification



Sequential example (cont'd): abstract control

- Finite-state diagram
 - states: 5 states
 - represent point in execution of machine
 - each state has inputs and outputs
 - transitions: 6 from state to state, 5 self transitions, 1 global
 - changes of state occur when clock says it's ok
 - based on value of inputs
 - inputs: reset, new, results of comparisons





Sequential example (cont'd): data-path vs. control

- Internal structure
 - data-path
 - storage for combination
 - comparators

control

- finite-state machine controller
- control for data-path
- state changes controlled by clock



* Multiplexer (MUX)

Sequential example (cont'd): FSM

- Finite-state machine (FSM)
 - refine state diagram to include internal structure





* state is not input, but internal variable

Sequential example (cont'd): encoding

- Encode state table
 - □ state can be: S1, S2, S3, OPEN, or ERR
 - needs at least 3 bits to encode: 000, 001, 010, 011, 100
 - and as many as 5: 00001, 00010, 00100, 01000, 10000
 - choose 4 bits: 0001, 0010, 0100, 1000, 0000
 - output mux can be: C1, C2, or C3
 - needs 2 to 3 bits to encode
 - choose 3 bits: 001, 010, 100
 - output open/closed can be: open or closed
 - needs 1 or 2 bits to encode
 - choose 1 bits: 1, 0

Sequential example (cont'd): encoding

Encode state table



Sequential example (cont'd): controller implementation

Implementation of the controller





terminologies



Summary

- That was what the entire course is about
 - converting solutions to problems into combinational and sequential networks effectively organizing the design hierarchically
 - doing so with a modern set of design tools that lets us handle large designs effectively
 - taking advantage of optimization opportunities