



# Data Structure

## Lecture#14: Non-Binary Trees (Chapter 6)

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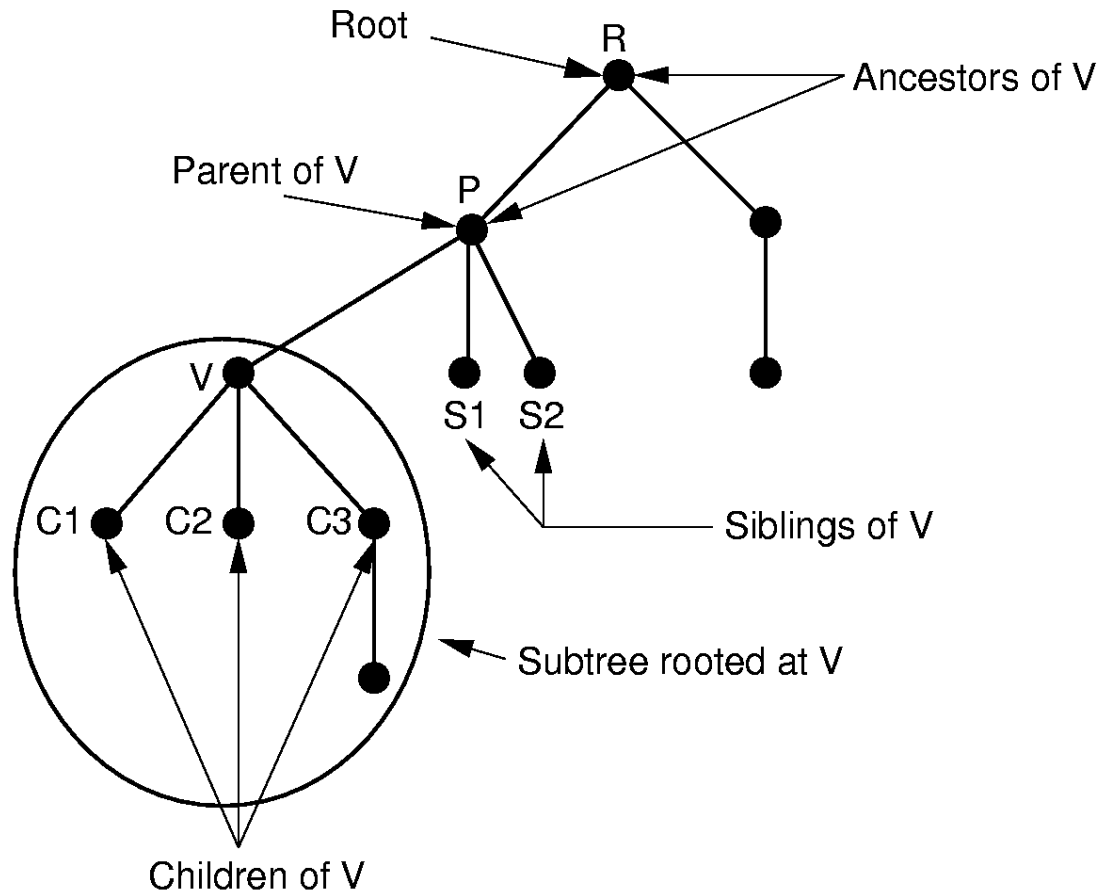


# In This Lecture

- The concept of the general tree, its ADT, and its operation
- Motivation and the main idea of the Union/Find operation
- Path compression for Union/Find



# General Trees





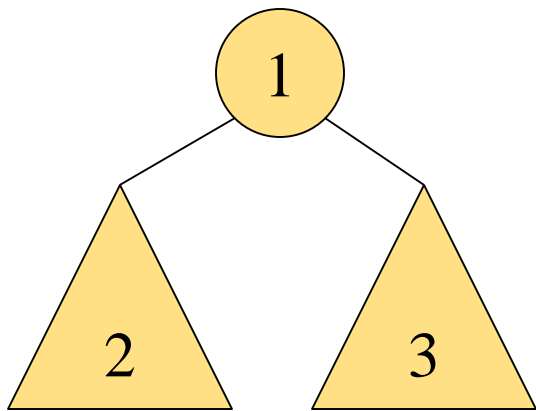
# General Tree Node

```
interface GTNode<E> {
    public E value();
    public boolean isLeaf();
    public GTNode<E> parent();
    public GTNode<E> leftmostChild();
    public GTNode<E> rightSibling();
    public void setValue(E value);
    public void setParent(GTNode<E> par);
    public void insertFirst(GTNode<E> n);
    public void insertNext(GTNode<E> n);
    public void removeFirst();
    public void removeNext();
}
```

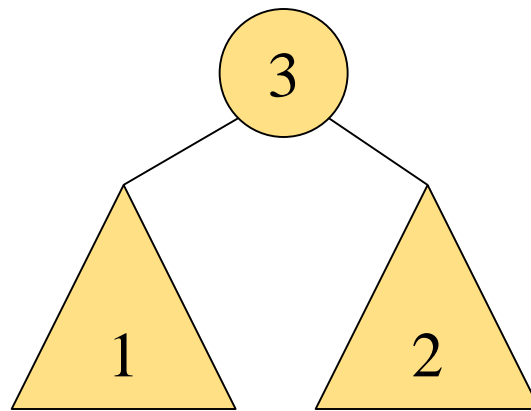


# General Tree Traversal

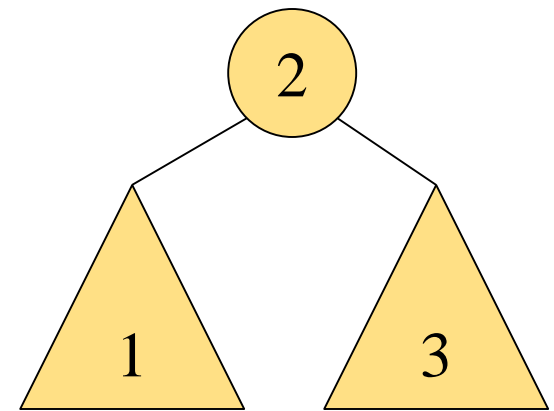
## ■ Binary Tree



Preorder Traversal

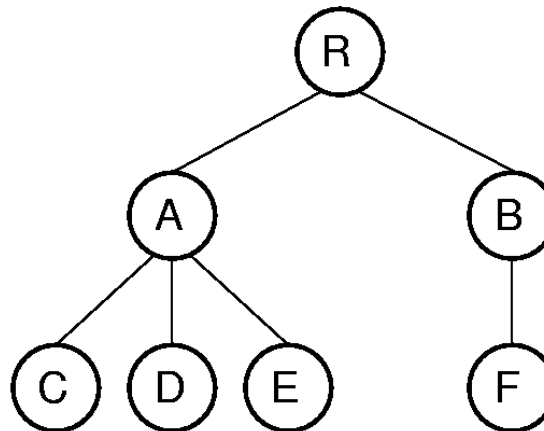


Postorder Traversal



Inorder Traversal

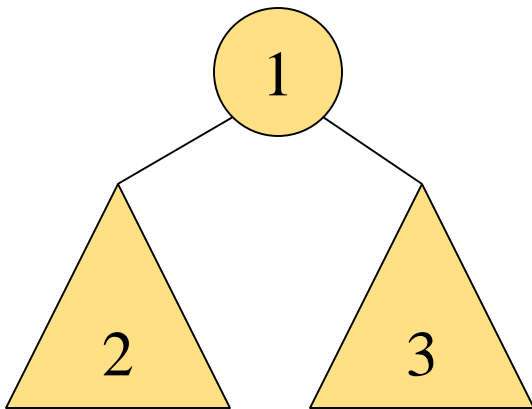
## ■ Non-binary Tree?



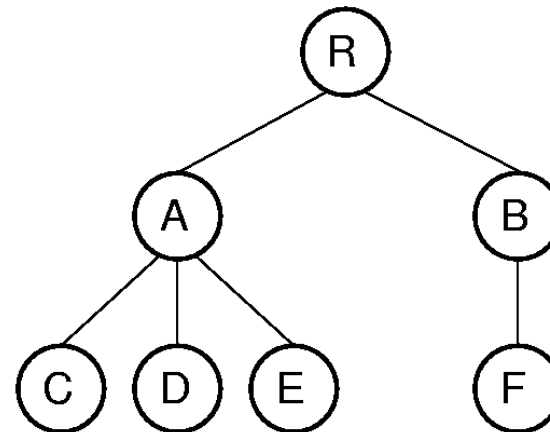


# Preorder Traversal

```
/** Preorder traversal for general trees */  
static <E> void preorder(GTNode<E> rt) {  
    PrintNode(rt);  
    if (!rt.isLeaf()) {  
        GTNode<E> temp = rt.leftmostChild();  
        while (temp != null) {  
            preorder(temp);  
            temp = temp.rightSibling();  
        }  
    }  
}
```



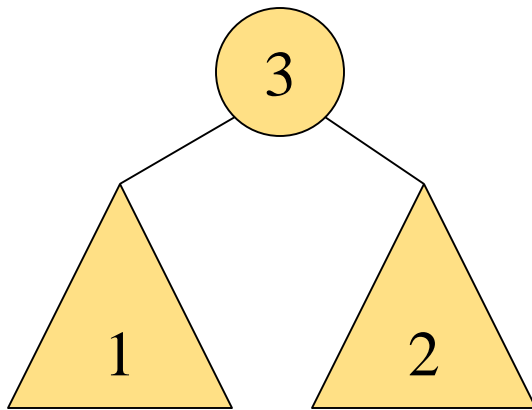
Preorder in binary tree



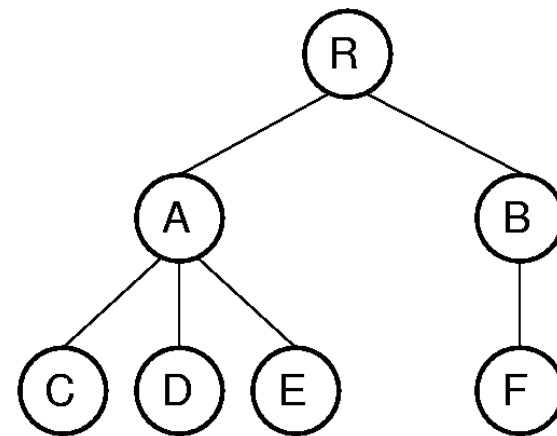
Well defined: R A C D E B F



# Postorder Traversal



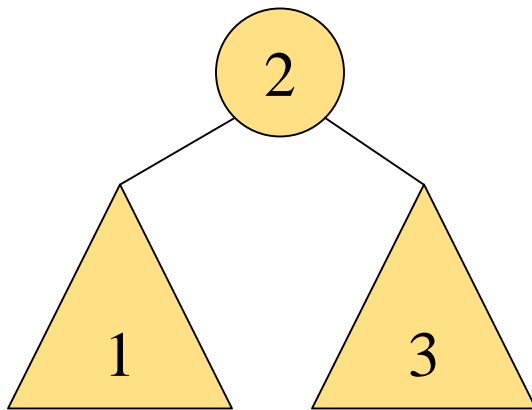
Postorder in binary tree



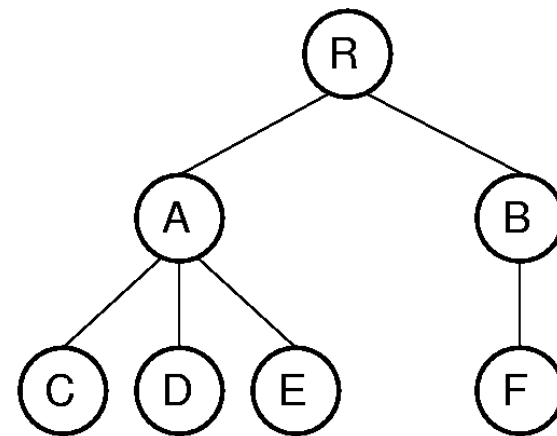
Well defined: C D E A F B R



# Inorder Traversal



Inorder in binary tree

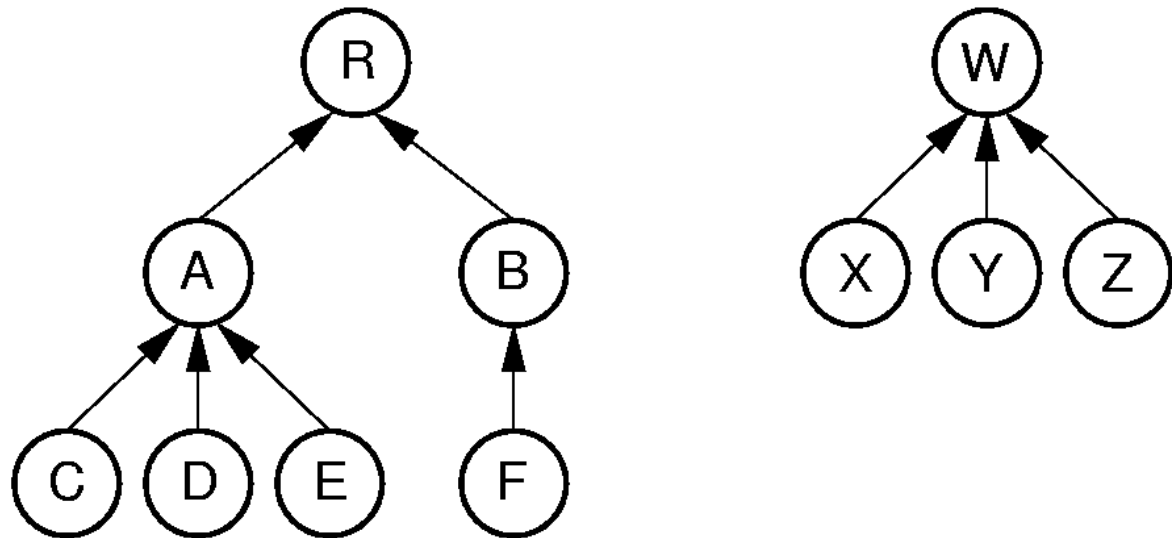


Well defined?





# Parent Pointer Implementation



Parent's Index		0	0	1	1	1	2		7	7	7
Label	R	A	B	C	D	E	F	W	X	Y	Z
Node Index	0	1	2	3	4	5	6	7	8	9	10

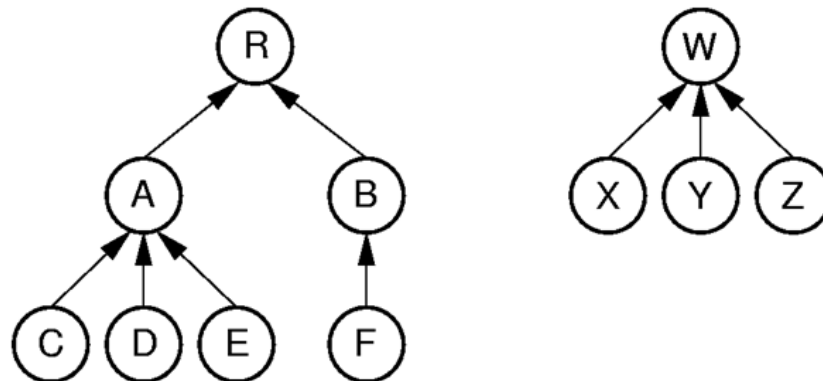


# Equivalence Class Problem

- The parent pointer representation is good for answering:
  - Are two elements in the same tree?

```
/** Determine if nodes in different trees */  
public boolean differ(int a, int b) {  
    Integer root1 = FIND(a);  
    Integer root2 = FIND(b);  
    return root1 != root2;  
}
```

**FIND function  
returns the root**

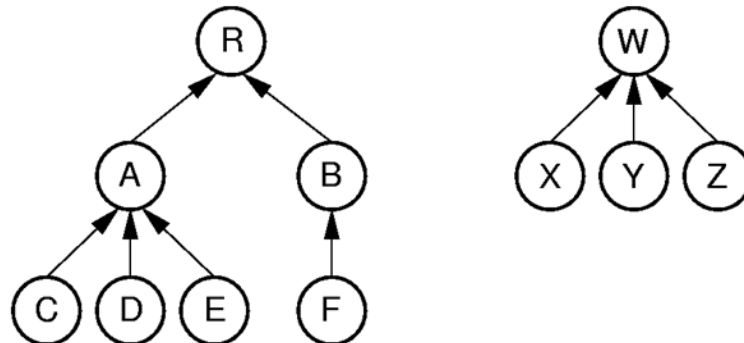




# Find

```
/** return the root of the curr's tree */  
public Integer FIND(Integer curr) {  
    if (array[curr] == null) return curr;  
    while (array[curr] != null)  
        curr = array[curr];  
    return curr;  
}
```

```
/** Determine if nodes in different trees */  
public boolean differ(int a, int b) {  
    Integer root1 = FIND(a);  
    Integer root2 = FIND(b);  
    return root1 != root2;  
}
```





# Union/Find

- Two important operations for equivalence class problem
  - UNION: merge two sets together
  - FIND: determine if two objects are in the same sets
  
- Applications
  - Graph connectivity
  - Clustering



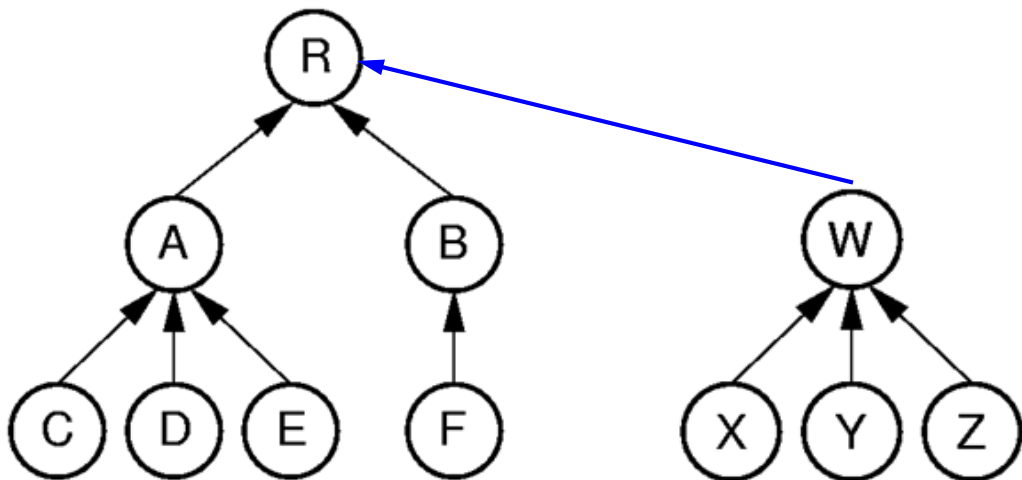
# Implementations for Union/Find

- UNION: merge two sets together
- FIND: determine if two objects are in the same sets; i.e., we want to get  $\text{id}(i)$  for an object  $i$
- Naïve algorithm  $n$ : # of items in  $A \cup B$ 
  - UNION of  $A$  and  $B$ : update  $\text{id}(i)$  for all  $i$  in  $A \cup B$ . Takes  $O(n)$
  - FIND of  $i$ :  $O(1)$  using  $\text{id}(i)$
- Using parent pointer implementation for tree
  - UNION of  $A$  and  $B$ :  $O(\text{FIND}(i)) = O(\log n)$
  - FIND of  $i$ :  $O(\log n)$ . If we use path compression, it can be  $O(1)$  for most  $i$



# Union

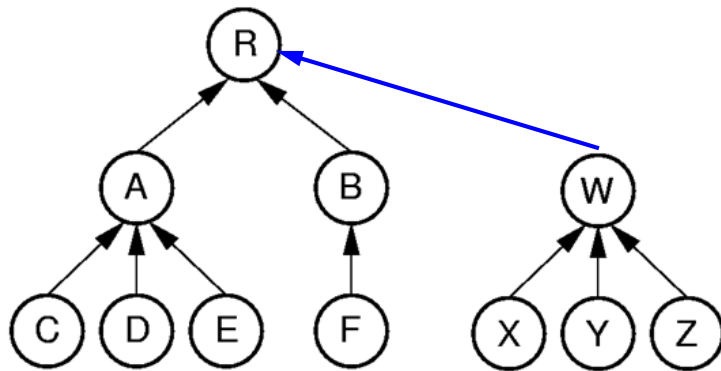
```
/** Merge two subtrees */  
public void UNION(int a, int b) {  
    Integer root1 = FIND(a); // Find a's root  
    Integer root2 = FIND(b); // Find b's root  
    if (root1 != root2) array[root2] = root1;  
}
```



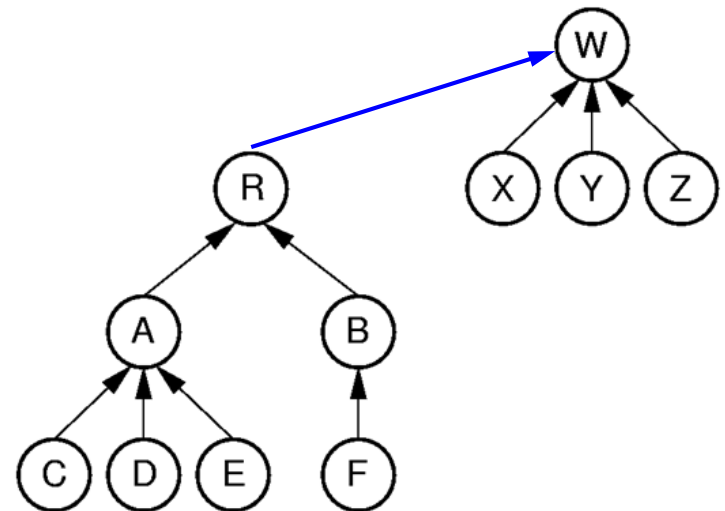


# Union

- In Union, we want to keep the depth of the tree small (Why?)
- Between the following two union methods, which is better? Why?



Method 1

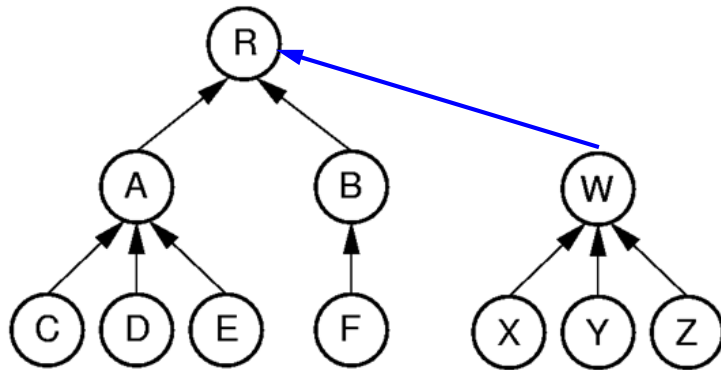


Method 2

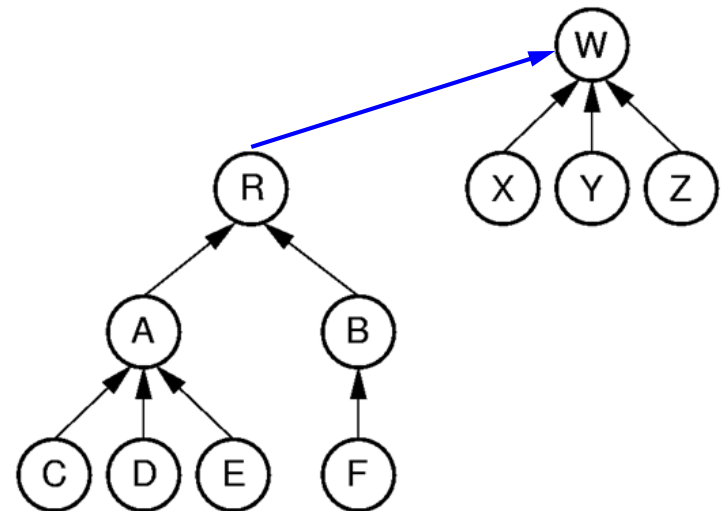


# Union

- Weighted union rule: join the tree with fewer nodes to the tree with more nodes.
  - Results in a tree with a small height



Method 1



Method 2





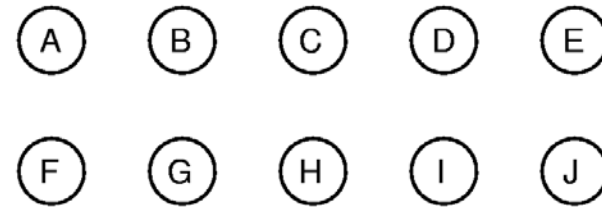
# Union

- (Theorem) Assume  $n$  nodes in  $n$  independent equivalent classes. Performing Union operations in any order results in a tree with depth at most  $\log n$ .
- (Proof)
  - Consider a node  $v$  with the maximum depth in the final tree. Initially,  $v$ 's depth was 0.  $v$ 's depth increased only when the subtree  $A$  containing  $v$  is merged with another tree  $B$ , and  $|A| < |B|$ . Then,  
 $v$ 's final depth = (# of times  $v$ 's depth increased)  $\leq \log n$ .



# Equiv Class Processing (1)

/	/	/	/	/	/	/	/	/	/
A	B	C	D	E	F	G	H	I	J
0	1	2	3	4	5	6	7	8	9

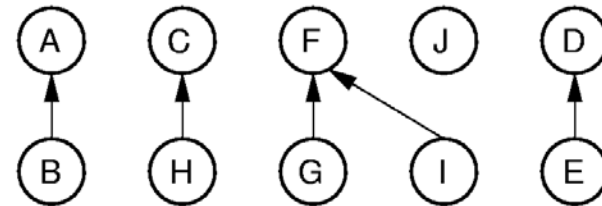


(Equivalence Relation)

(A,B) (C,H) (G,F) (D,E) (I,F)

(a)

/	0	/	/	3	/	5	2	5	/
A	B	C	D	E	F	G	H	I	J
0	1	2	3	4	5	6	7	8	9

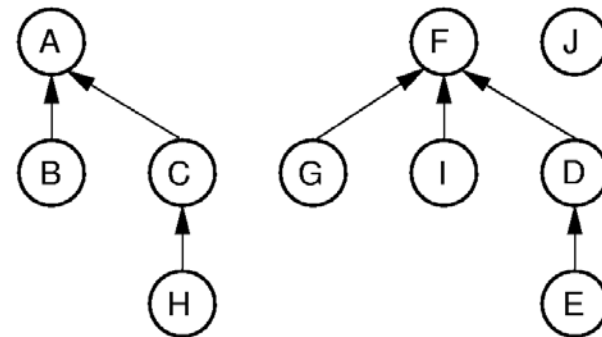


(Equivalence Relation)

(H,A) (E,G)

(b)

/	0	0	5	3	/	5	2	5	/
A	B	C	D	E	F	G	H	I	J
0	1	2	3	4	5	6	7	8	9

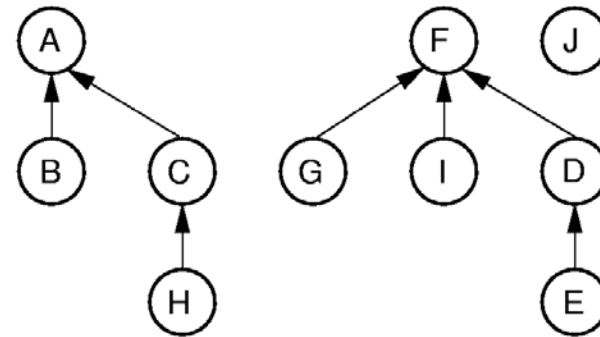


(c)



# Equiv Class Processing (2)

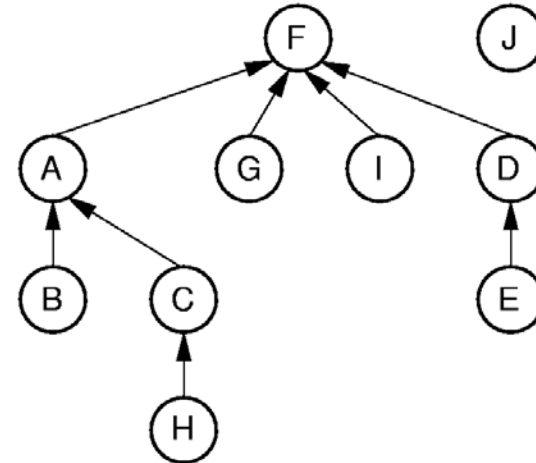
/	0	0	5	3	/	5	2	5	/
A	B	C	D	E	F	G	H	I	J
0	1	2	3	4	5	6	7	8	9



(c)

(Equivalence Relation)  
(H,E)

5	0	0	5	3	/	5	2	5	/
A	B	C	D	E	F	G	H	I	J
0	1	2	3	4	5	6	7	8	9

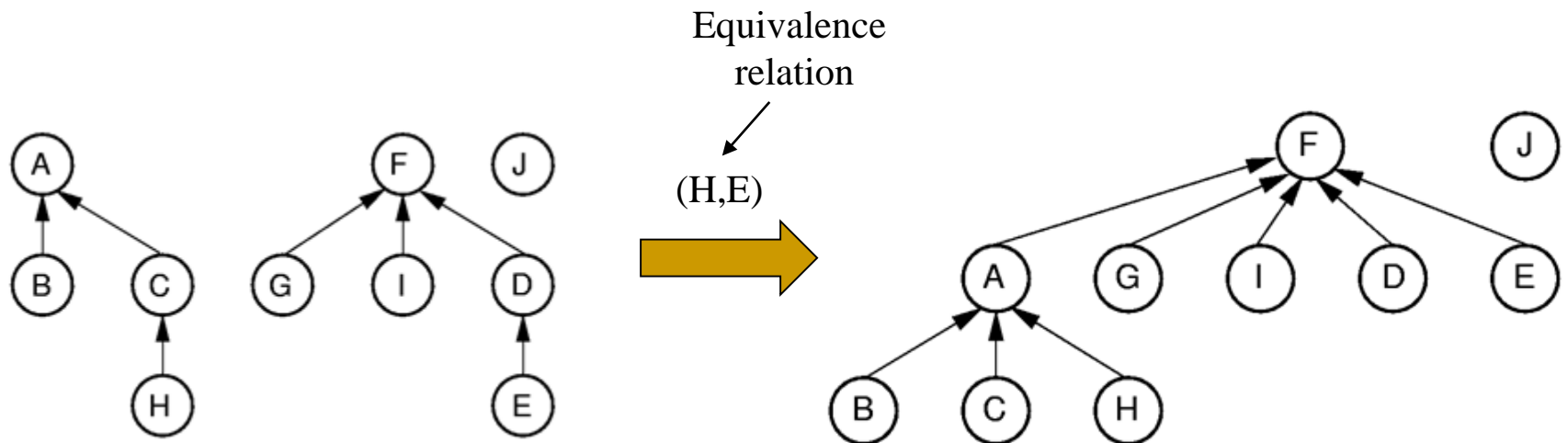


(d)



# Path Compression

- How can we further decrease the max depth of the tree?
  - Path Compression: while doing FIND, make each node on the path to point to root

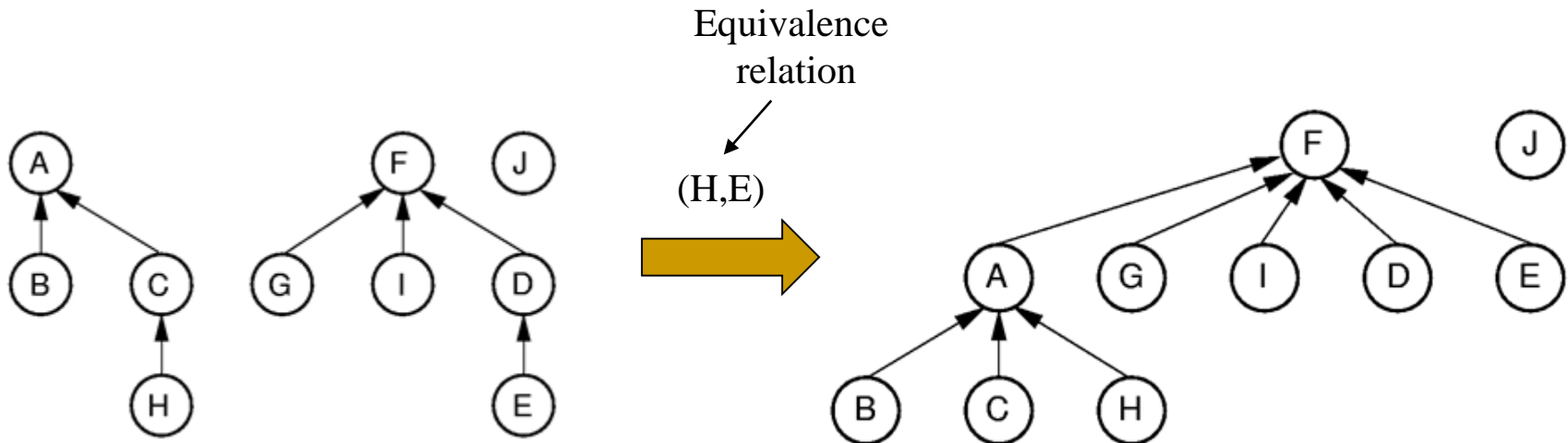




# Path Compression

```
/** FIND for path compression */  
public Integer FIND(Integer curr) {  
    if (array[curr] == null) return curr;  
    array[curr] = FIND(array[curr]);  
    return array[curr];  
}
```

← Compare this with  
FIND in slide 11





# What you need to know

- The ADT and operations of general tree
- Goal and implementation of the Union/Find operation
- Goal and implementation of Path compression for Union/Find



# Questions?