Intro to DB



# Chapter 3: SQL

- Overview
- Data Definition
- Basic Structure
- Additional Basic Operations
- Set Operations
- Null Values
- Aggregate Functions
- Nested Subqueries
- Modification of the Database

# **History**

- IBM Sequel language developed as part of System R project at the IBM San Jose Research Laboratory
- Renamed Structured Query Language (SQL)
- ANSI and ISO standard SQL:
  - SQL-86, SQL-89, SQL-92
  - SQL:1999 (Y2K!), SQL:2003, SQL:2008
- Commercial systems offer most, if not all, SQL-92 features, plus varying feature sets from later standards and special proprietary features.
  - Not all examples here may work on your particular system.

### Create Table Construct

• An SQL relation is defined using the **create table** command:

```
create table r (A_1 D_1, A_2 D_2, ..., A_n D_n,
(integrity-constraint<sub>1</sub>),
...,
(integrity-constraint<sub>k</sub>))
```

- r is the name of the relation
- each  $A_i$  is an attribute name in the schema of relation r
- $D_i$  is the data type of values in the domain of attribute  $A_i$
- Example:

insert into instructor values ('10211', 'Smith', 'Biology', 66000);

# **Domain Types in SQL**

- **char(n):** Fixed length character string, with user-specified length *n*.
- varchar(n): Variable length character strings, with user-specified maximum length *n*.
- int: Integer (a finite subset of the integers that is machine-dependent).
- smallint: Small integer (a machine-dependent subset of the integer domain type).
- numeric(p, d): Fixed point number, with user-specified precision of p digits, with n digits to the right of decimal point.
- **real, double precision:** Floating point and double-precision floating point numbers, with machine-dependent precision.
- **float(n):** Floating point number, with user-specified precision of at least *n* digits.

# **Integrity Constraints in Create Table**

- not null
- **primary key**  $(A_1, ..., A_n)$
- foreign key  $(A_m, ..., A_n)$  references r

Declare *dept\_name* as to be a reference to *department*.

primary key declaration on an attribute automatically ensures not null

# And a Few More Relation Definitions

• create table student (

IDvarchar(5),namevarchar(20) not null,dept\_namevarchar(20),tot\_crednumeric(3,0),primary key (ID),foreign key (dept\_name) references department) );

• create table takes (

varchar(5),
varchar(8),
varchar(8),
varchar(6),
<b>numeric</b> (4,0),
varchar(2),
(ID, course_id, sec_id, semester, year),
ID) references student,
course_id, sec_id, semester, year) references section );

Note: sec\_id can be dropped from primary key above, to ensure a student cannot be registered for two sections of the same course in the same semester

#### And more still

#### • create table course (

course_id	varchar(8) primary key,
title	<b>varchar(</b> 50),
dept_name	varchar(20),
credits	<b>numeric</b> (2,0),
foreign key	(dept_name) references department) );

Primary key declaration can be combined with attribute declaration as shown above

#### **Drop and Alter Table Constructs**

- **drop table:** deletes all information about the dropped relation from the database.
- alter table: used to add or drop attributes to an existing relation

#### alter table r add A D

where A is the name of the attribute to be added to relation r and D is the domain of A.

• *null* is assigned to the new attribute for each tuple

#### alter table *r* drop *A*

where *A* is the name of an attribute of relation *r* 

(dropping of attributes is not supported by many databases)

# **Basic Structure of SQL Queries**

- SQL is based on set and relational operations with certain modifications and enhancements
- A typical SQL query has the form:

```
select A_1, A_2, ..., A_n
from r_1, r_2, ..., r_m
where P
```

- *A<sub>i</sub>s* represent attributes
- *r<sub>i</sub>s* represent relations
- *P* is a predicate.
- The result of an SQL query is a relation.

#### The select Clause

Find the names of all instructors:

An asterisk in the select clause denotes "all attributes"

select \*
from instructor

- NOTE:
  - SQL does not permit the '-' character in names (use '\_' in a real implementation).
  - SQL names are case insensitive.

#### The select Clause (cont.)

- SQL allows duplicates in relations as well as in query results.
- To force the elimination of duplicates, insert the keyword **distinct** after **select**.
- Find the names of all departments with instructor, and remove duplicates

The keyword all specifies that duplicates not be removed

select all dept\_name
from instructor

#### The select Clause (cont.)

- The **select** clause can contain arithmetic expressions
  - +, -, \*, /
  - on constants or attributes of tuples.
- The query:

would return a relation that is the same as the *instructor* relation, except that the value of the attribute *salary* is divided by 12.

#### The where Clause

- Corresponds to the selection predicate of the relational algebra.
- Predicate involving attributes of the relations that appear in the from clause.
- Find all instructors in Comp. Sci. dept with salary > 80000

- Comparison conditions: >, <, =, <=, >=, !=
- Logical connectives: and, or, not
- Comparisons can be applied to results of arithmetic expressions
- SQL includes a **between** comparison operator

where salary between 90000 and 100000

### The from Clause

- Lists the relations to be scanned in the evaluation of the expression.
- Corresponds to the Cartesian product operation of the relational algebra.
- Instructor x teaches

select \*
from instructor, teaches

#### instructor

#### teaches

Π	D	nı	ame	dep	ot_name	sa	lary	[	ID	course_id	sec_id	l semest	ter	year
101	.01	Sri	nivasan	Co	omp. Sci.	65	5000	Ī	10101	CS-101	1	Fall		2009
121	.21	Wı	1	Fi	nance	90	0000		10101	CS-315	1	Spring	g	2010
151	.51	Mo	ozart	M	usic	40	0000		10101	CS-347	1	Fall		2009
222	222	Eir	nstein	Ph	iysics	95	5000		12121	FIN-201	1	Spring	3	2010
323	343	El	Said	Hi	story	60	0000		15151	MU-199	1	Spring	3	2010
~~~·			11 1	731	•		7000 I		22222	PHY-101	1	Fall		2009 I
	inst.	ID	name		dept_nan	ne	salary	teac	ches.ID	course_id	sec_id	semester	year	
[	1010	)1	Sriniva	san	Comp.	Sci.	65000	1	0101	CS-101	1	Fall	2009	
	1010	01	Sriniva	san	Comp. S	Sci.	65000	1	0101	CS-315	1	Spring	2010	
	1010	01	Sriniva	san	Comp. S	Sci.	65000	1	0101	CS-347	1	Fall	2009	
	1010	01	Sriniva	san	Comp. S	Sci.	65000	12	2121	FIN-201	1	Spring	2010	
	1010	01	Sriniva	san	Comp. S	Sci.	65000	1	5151	MU-199	1	Spring	2010	
	1010	01	Sriniva	san	Comp.	Sci.	65000	2	2222	PHY-101	1	Fall	2009	
					•••									
	1212	21	Wu		Finance		90000	1	0101	CS-101	1	Fall	2009	
	1212	21	Wu		Finance	2	90000	1	0101	CS-315	1	Spring	2010	
	1212	21	Wu		Finance		90000	1	0101	CS-347	1	Fall	2009	
	1212	21	Wu		Finance		90000	1	2121	FIN-201	1	Spring	2010	
	1212	21	Wu		Finance		90000	1	5151	MU-199	1	Spring	2010	
	1212	21	Wu		Finance		90000	2	2222	PHY-101	1	Fall	2009	
												••••		

select \* from instructor, teaches

#### Joins

 For all instructors who have taught some course, find their names and the course ID of the courses they taught.

• Find the course ID, semester, year and title of each course offered by the Comp. Sci. department

section	course
<u>course_i</u> d <u>sec_id</u> <u>semeste</u> r <u>year</u> building	<u>course_i</u> d title dept_name credits
room_no time_slot_id	

### **Schema Diagram for the University**



#### **SQL Examples**

• Find the IDs of students advised by an instructor named Einstein.

 Find the IDs of all students who were taught by an instructor named Einstein; make sure there are no duplicates in the result.

### The Rename Operation

 The SQL allows renaming relations and attributes using the as clause: old-name as new-name

**select** *ID*, *name*, *salary/12* **as** *monthly\_salary* **from** *instructor* 

Find the names of all instructors who have a higher salary than some instructor in 'Comp. Sci'.

- *S*, *T* are called *tuple variables*
- Keyword **as** is optional and may be omitted *instructor* **as**  $T \equiv instructor T$

# **String Operations**

- For comparisons on character strings
- Patterns are described using special characters:
  - percent (%): matches any substring.
  - underscore (\_). matches any character.
- Find all courses whose title includes the substring "data".

- Escape character to specify % and \ within string like '100\%'
- A variety of string operations such as
  - concatenation (using "||")
  - case conversion, string length, substrings, etc.

# **Ordering the Display of Tuples**

List in alphabetic order all instructors

- desc for descending order or asc for ascending order (default)
  - order by name desc
- Can sort on multiple attributes
  - **order by** *dept\_name, name*
  - order by *dept-name* desc, *name* asc

#### **Set Operations**

- The set operations: union, intersect, except correspond to the relational algebra operations
   ∪, ∩, −.
- To retain all duplicates use multiset versions:

union all, intersect all and except all.

Suppose a tuple occurs *m* times in *r* and *n* times in *s*, then, it occurs:

![](_page_22_Figure_6.jpeg)

### **Set Operations (cont.)**

Find all customers who have a loan, an account, or both:

(select customer-name from depositor) union (select customer-name from borrower)

Find all customers who have both a loan and an account.

(select customer-name from depositor) intersect (select customer-name from borrower)

Find all customers who have an account but no loan.

(select customer-name from depositor)except(select customer-name from borrower)

#### **Null Values**

- It is possible for tuples to have a null value, denoted by *null*, for some of their attributes
- The result of any arithmetic expression involving null is null
  - Example: 5 + *null* returns null
- The predicate is null can be used to check for null values.
   Find all instructors whose salary is null.

select name from instructor where salary is null

• Why not the following?

select name
from instructor
where salary = null

### **Null Values and Three Valued Logic**

- Any comparison with null returns unknown
  - Example: 5 < null or null <> null or null = null
- Three-valued logic using the truth value *unknown*:
  - OR: (unknown or true) = true, (unknown or false) = unknown (unknown or unknown) = unknown
  - AND: (true and unknown) = unknown, (false and unknown) = false, (unknown and unknown) = unknown
  - NOT: (not unknown) = unknown
  - "*P* is unknown" evaluates to true if predicate *P* evaluates to unknown

# **Aggregate Functions**

• Operate on the multiset of values of a column of a relation, and return a value

avg: average valuemin: minimum valuemax: maximum valuesum: sum of valuescount: number of values

# **Aggregate Functions (Cont.)**

• Find the average salary of instructors in the Computer Science department

```
select avg (salary)
from instructor
where dept_name= 'Comp. Sci.';
```

• Find the total number of instructors who teach a course in the Spring 2010 semester

Find the number of tuples in the *course* relation

select count (\*)
from course;

# **Aggregate Functions – Group By**

 Find the average salary of instructors in each department select dept\_name, avg (salary) from instructor group by dept\_name;

ID	name	dept_name	salary
76766	Crick	Biology	72000
45565	Katz	Comp. Sci.	75000
10101	Srinivasan	Comp. Sci.	65000
83821	Brandt	Comp. Sci.	92000
98345	Kim	Elec. Eng.	80000
12121	Wu	Finance	90000
76543	Singh	Finance	80000
32343	El Said	History	60000
58583	Califieri	History	62000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
22222	Einstein	Physics	95000

	Note: departments	with no	instructor w	vill not appear	in result
--	-------------------	---------	--------------	-----------------	-----------

dept_name	avg_salary
Biology	72000
Comp. Sci.	77333
Elec. Eng.	80000
Finance	85000
History	61000
Music	40000
Physics	91000

# **Aggregation (Cont.)**

- Attributes in **select** clause outside of aggregate functions must appear in **group by** list
  - /\* erroneous query \*/
     select dept\_name, ID, avg (salary)
     from instructor
     group by dept\_name;

# **Null Values and Aggregates**

Total all salaries

**select sum** (salary) **from** instructor

- Above statement ignores null amounts
- result is null if there is no non-null amount

# **Nested Subqueries**

A subquery is

• Common use of subqueries:

perform tests for set membership, set comparisons, and set cardinality.

# **Nested Query – Examples**

Find courses offered in Fall 2009 and in Spring 2010

select distinct course\_id
from section

Find courses offered in Fall 2009 but not in Spring 2010

select distinct course\_id
from section

# **Nested Query – Examples**

 Find the total number of (distinct) students who have taken course sections taught by the instructor with *ID* 10101

select count (distinct <i>ID</i> )	
from takes	

 Note: Above query can be written in a much simpler manner. The formulation above is simply to illustrate SQL features

## **Set Comparison**

 Find names of instructors with salary greater than that of some (at least one) instructor in the Biology department.

> **select distinct** *T.name* **from** *instructor* **as** *T*, *instructor* **as** *S*

Same query using > some clause

select name from instructor

# Set Comparison (cont.)

 Find the names of all instructors whose salary is greater than the salary of all instructors in the Biology department.

![](_page_35_Figure_2.jpeg)

# **Test for Empty Relations**

- The **exists** construct returns the value **true** if the argument subquery is nonempty.
- exists  $r \Leftrightarrow r \neq \emptyset$
- not exists  $r \Leftrightarrow r = \emptyset$

### **Correlation Variables**

 Yet another way of specifying the query "Find all courses taught in both the Fall 2009 semester and in the Spring 2010 semester"

```
select course_id
from section as S
where semester = 'Fall' and year= 2009 and
    exists (select *
        from section as T
        where semester = 'Spring' and year= 2010
        and S.course_id= T.course_id);
```

- Correlated subquery
- S: Correlation name or correlation variable

#### **Not Exists**

Find all students who have taken all courses offered in the Biology department.

```
select distinct S.ID, S.name
from student as S
where not exists ( (select course_id
            from course
            where dept_name = 'Biology')
```

- Note that  $X Y = \emptyset \iff X \subseteq Y$
- Note: Cannot write this query using = all and its variants

### **Modification of the Database – Deletion**

Delete all courses of Appl. Math department

delete from course
where dept\_name = 'Appl. Math'

 Delete all tuples in the *instructor* relation for those instructors associated with a department located in the Watson building.

### **Modification of the Database – Deletion**

Delete all instructors whose salary is less than the average salary of instructors.

delete from instructor

where salary< (select avg (salary) from instructor);</pre>

Problem:

- Solution used in SQL:
  - 1. First, compute avg salary and find all tuples to delete
  - 2. Next, delete all tuples found above

(without recomputing **avg** or retesting the tuples)

### **Modification of the Database – Insertion**

• Add a new tuple to *course* 

```
insert into course
values ('CS-437', 'Database Systems', 'Comp. Sci.', 4);
```

or equivalently

Add a new tuple to student with tot\_creds set to null

insert into student
values ('3003', 'Green', 'Finance', null);

# **Insertion (Cont.)**

Add all instructors to the student relation with tot\_creds set to 0

insert into student
 select ID, name, dept\_name, 0
 from instructor

- The select from where statement is evaluated fully before any of its results are inserted into the relation
  - otherwise queries like the following would cause problems (if *table1* did not have any primary key defined)

insert into table1 select \* from table1

### **Modification of the Database – Updates**

- Increase salaries of instructors whose salary is over \$100,000 by 3%, and all others receive a 5% raise
  - Write two update statements:

![](_page_43_Picture_3.jpeg)

# END OF CHAPTER 3