

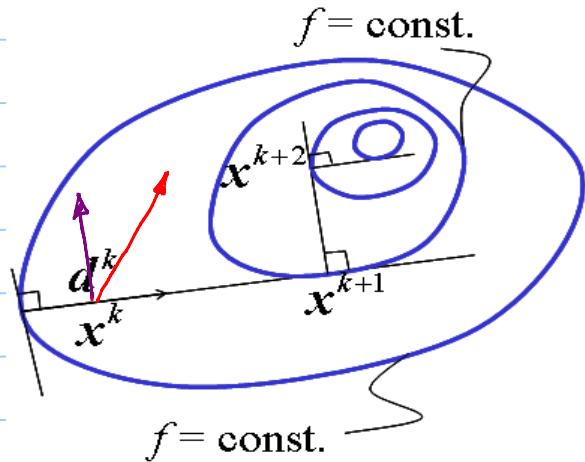
3-3 : Conjugate Gradient Method

노트 제목

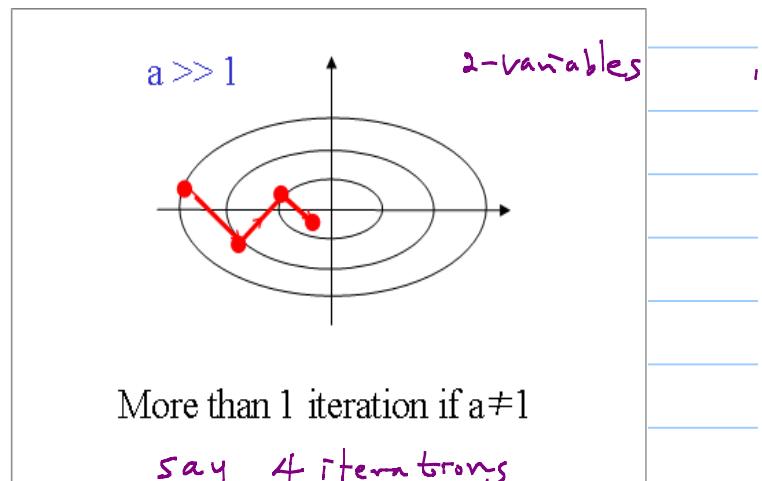
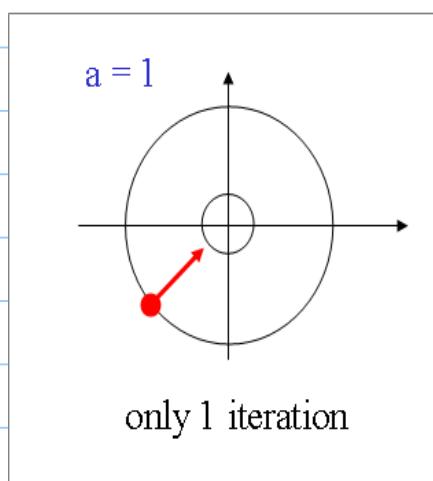
①

< Motivation >

- i) Recall
the convergence
pattern of the
steepest descent
method



- ii) Consider $f(x_1, x_2) = x_1^2 + \alpha x_2^2$ ← quadratic ftn
of 2 variables



⇒ We want a method yielding a min of $f(x)$
within "n" iterations if $f(x)$ is a
quadratic function of "n" variables

(2)

⇒ "Conjugate Gradient Method"

(a very powerful method that
also works for general non-quadratic
functions)

Will study

1) What are conjugate directions?

i.e., definition of conjugate direction

2) Proof to show n-iteration convergence
of $Q(x_1, \dots, x_n)$ with conjugate
Search direction.

3) How to find conjugate direction?

(approximate for non-quadratic fns)

(1) < Definition of Conjugate Directions >

Given $A: n \times n$ positive-definite
symmetric matrix

\underline{d}^i : conjugate direction

if $(\underline{d}^i)^T A \underline{d}^j = 0$ for $i \neq j$

(3)

Examples of conjugate directions

Let $A: 3 \times 3$ sym matrix

e_i : eigenvector of A ($e_i^T A e_j = \delta_{ij}$)

Then

$$\left\{ \begin{array}{l} (e_1, e_2, e_3) : \text{conjugate directions} \\ (e_1 + e_2, e_1 - e_2, e_3) : " \\ \vdots \end{array} \right.$$

• X A special case of conjugate directions
= eigenvectors

(2) < proof of n-iteration convergence >

Min of a quadratic ftn of n

variables is found in n iterations or less.

■ Consider $Q(\tilde{x}) = \frac{1}{2} \tilde{x}^T A \tilde{x} + B^T \tilde{x} + C$

a quadratic ftn

↑
positive-definite
symmetric matrix

$$\tilde{x} = \{x_1, \dots, x_n\}^T$$

(4)

ProofLet $\tilde{x}^* = \min$ pointthen $\nabla Q(\tilde{x}^*) = \underline{A} \tilde{x}^* + \underline{B} = 0 \quad (1)$

$$\begin{aligned}
\frac{\partial Q}{\partial x_k} &= \frac{\partial}{\partial x_k} \left\{ \frac{1}{2} \sum_{i=1}^n \sum_{j=1}^n x_i A_{ij} x_j + \sum_{j=1}^n B_j x_j + C \right\} \\
&= \frac{1}{2} \left[\sum_{i=1}^n \sum_{j=1}^n (A_{ik} A_{ij} x_j + x_i A_{ij} A_{jk}) \right. \\
&\quad \left. + \sum_{j=1}^n B_j A_{jk} \right] \\
&= \frac{1}{2} \left[\sum_{j=1}^n A_{kj} x_j + \sum_{i=1}^n x_i A_{ik} \right] + B_k \\
&= \frac{1}{2} \left[\sum_{i=1}^n A_{ki} x_i + \sum_{i=1}^n \underbrace{A_{ik} x_i}_{A_{kk}} \right] + B_k \\
&= \sum_{i=1}^n A_{ki} x_i + B_k
\end{aligned}$$

Expand \tilde{x}^* in terms of orthogonal $\underline{d}^i \{i=0 \dots n-1\}$
wrt \underline{A}

Such that

$$\tilde{x}^* = \tilde{x}^0 + \sum_{i=0}^{n-1} \beta^i \underline{d}^i \quad (2)$$

↑ used to represent
an arbitrary starting point

(5)

To determine β^i , (2) \rightarrow (1)

$$\underline{\beta} + \underline{A} \underline{x}^0 + \sum_{i=0}^{n-1} \beta^i \underline{A} \underline{d}^i = 0 \quad (3)$$

$(\underline{d}^i)^T (3)$:

$$\underline{d}^i T (\underline{\beta} + \underline{A} \underline{x}^0) + \sum_{i=0}^{n-1} \beta^i \underline{d}^i \underline{A} \underline{d}^i = 0 \quad (4)$$

$\underbrace{= 0 \text{ if } i \neq j}$

\downarrow

$$\beta^i \underline{d}^i T \underline{A} \underline{d}^i$$

Thus

$$\beta^i = - \frac{(\underline{\beta} + \underline{A} \underline{x}^0)^T \underline{d}^i}{(\underline{d}^i)^T \underline{A} \underline{d}^i} \quad (5)$$

$i = 0, 1, 2, \dots, n-1$

\Rightarrow Message: • For "any" starting point \underline{x}^0 ,
 the min \underline{x}^* of $\underline{Q}(\underline{x})$ is
 expressed in terms of
 (\underline{d}^i) ; expansion coefficients β_i are
 computed by (5).

②

Let us now prove that min of quadratic function $Q(\underline{x})$ can be found with n iterations if we update \underline{x}^i using \underline{d}^i as

$$\begin{array}{l} \underline{x}^{i+1} = \underline{x}^i + \lambda^{i*} \underline{d}^i \\ \quad \quad \quad (i=0, 1, \dots, n-1) \end{array} \quad \text{conjugate direction}$$

where λ^{i*} is determined by 1-D search by

$$\min Q(\underline{x}^i + \lambda^{i*} \underline{d}^i)$$

$$\Leftrightarrow \boxed{0 = \frac{dQ}{d\lambda^i} \Big|_{\lambda^i = \lambda^{i*}}} \quad (7)$$

* If we can prove that

$$\boxed{\lambda^{i*} = \beta^i}, \quad (i=0, \dots, n-1)$$

then

$$\underline{x}^1 = \underline{x}^0 + \lambda^{0*} \underline{d}^0 = \underline{x}^0 + \beta^0 \underline{d}^0$$

$$\underline{x}^2 = \underline{x}^1 + \lambda^{1*} \underline{d}^1 = \underline{x}^0 + \beta^0 \underline{d}^0 + \beta^1 \underline{d}^1$$

:

(7)

$$\underline{x}^n = \underline{x}^0 + \sum_{i=0}^{n-1} \beta^i \underline{d}^i$$

i.e., updating by (6) and (7) yields solutions within n iterations.

Q Let us prove $\lambda^{i*} = \beta^i$

Use Eq. (7)

$$0 = \frac{d Q(\underline{x}^i + \lambda^i \underline{d}^i)}{d x^i} \Big|_{\lambda^i = \lambda^{i*}}$$

$$= \frac{\partial}{\partial \lambda^i} \left[\frac{1}{2} (\underline{x}^i + \lambda^i \underline{d}^i)^T \Delta (\underline{x}^i + \lambda^i \underline{d}^i) + \underline{B}^T (\underline{x}^i + \lambda^i \underline{d}^i) + C \right] \Big|_{\lambda^i = \lambda^{i*}}$$

$$= \underline{d}^i T \Delta (\underline{x}^i + \lambda^{i*} \underline{d}^i) + \underline{B}^T \underline{d}^i$$

$$= \lambda^{i*} (\underline{d}^{i* T} \Delta \underline{d}^i) + \underbrace{\underline{d}^{i* T} \Delta \underline{x}^i}_{\text{Scalar}} + \underline{B}^T \underline{d}^i$$

$$= \underline{x}^{i* T} \Delta \underline{d}^i$$

$$= \underline{x}^{i* T} \underbrace{\Delta}_{\text{Scalar}} \underline{d}^i$$

(8)

$$\therefore \lambda^{i*} = - \frac{(\underline{B}^T + \underline{x}^{i*} \Delta) \underline{d}^i}{\underline{d}^{iT} \Delta \underline{d}^i} \quad (8)$$

To simplify (8), note that

$$\begin{aligned} \textcircled{1} \quad \underline{x}^i &= \underline{x}^{i-1} + \lambda^{(i-1)*} \underline{d}^{i-1} \\ &= \underline{x}^{i-2} + \lambda^{(i-2)*} \underline{d}^{i-2} + \lambda^{(i-1)*} \underline{d}^{i-1} \\ &= \underline{x}^0 + \sum_{k=0}^{i-1} \lambda^{k*} \underline{d}^k \end{aligned}$$

$$\begin{aligned} \textcircled{2} \quad \underline{x}^{iT} \Delta \underline{d}^i &= (\underline{x}^0 + \sum_{k=0}^{i-1} \lambda^{k*} \underline{d}^k)^T \Delta \underline{d}^i \\ &= \underline{x}^{0T} \Delta \underline{d}^i + \sum_{k=0}^{i-1} \lambda^{k*} \underline{d}^{kT} \underbrace{\Delta}_{=0} \underline{d}^i \\ &= \underline{x}^{0T} \Delta \underline{d}^i \\ &= \underline{x}^{0T} \Delta^T \underline{d}^i \\ &= (A \underline{x}^0)^T \underline{d}^i \end{aligned}$$

because $\begin{cases} \underline{d}^k, \text{conjugate} \\ i \neq j \end{cases}$

$$\therefore \lambda^{i*} = - \frac{(\underline{B} + A \underline{x}^0)^T \underline{d}^i}{\underline{d}^{iT} \Delta \underline{d}^i} \equiv \beta^i$$


(9)

(3) < How to find \underline{d} ? >

- Result:

① for quadratic function $Q(\underline{x}) = \frac{1}{2} \underline{x}^T \underline{A} \underline{x} + \underline{B}^T \underline{x} + C$

$$\cdot \underline{d}^0 = -\underline{g}^0 \quad (\underline{g} = \nabla Q)$$

$$\cdot \underline{d}^{k+1} = -\underline{g}^{k+1} + \beta^k \underline{d}^k \quad (k \geq 1)$$

$$\text{where } \beta^k = \frac{(\underline{g}^{k+1})^T (\underline{A} \underline{d}^k)}{\underline{d}^k T (\underline{A} \underline{d}^k)}$$

② for non-quadratic function $f(\underline{x})$

replace β^k as

$$\beta^k = -\frac{\underline{g}^{k+1}^T \underline{g}^{k+1}}{\underline{g}^k T \underline{g}^k}$$

(Key point: no second-derivative information such as Hessian ($\approx \underline{A}$) is needed
 \rightarrow due to Fletcher-Reeve)

Overall approach

i) derive \underline{d} for $Q(\underline{x})$

ii) extend the result for non-quadratic functions

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< Derivation >

(A) For $Q(\underline{x}) = \frac{1}{2} \underline{x}^T A \underline{x} + \underline{B}^T \underline{x} + C$ first

- \underline{d}_0 : can be any non-zero vector to be a conjugate vector at the beginning

- thus, choose $d_0 = -\frac{g}{d}$.

where

$$\begin{aligned} \nabla Q(x^k) &= A x^k + B \\ &\approx \quad (a) \end{aligned}$$

- ## Search Algorithm

$$\underline{x}^{k+1} = \underline{x}^k + \alpha^k \underline{d}^k \quad (b)$$

found by 1-D search

- Conjugate direction ok

$$\tilde{d}^{k+1} = -\underline{g}^{k+1} + \beta^k \underline{d}^k \quad (c2)$$

(11)

- To determine β^k , use the conjugate condition.

$$\underline{d}^{k+1} \perp \underline{d}^{k+1} = 0$$

$$[-\underline{g}^k + \beta^{k+1} \underline{d}^{k+1}]^T \perp \underline{d}^{k+1} = 0$$

$$\rightarrow \beta^{k+1} = \frac{(\underline{g}^k)^T \perp \underline{d}^{k+1}}{(\underline{d}^{k+1})^T \perp \underline{d}^{k+1}}$$

or

$$\boxed{\beta^k = \frac{(\underline{g}^k)^T (\perp \underline{d}^k)}{(\underline{d}^k)^T (\perp \underline{d}^k)}} \quad (d)$$

for non-quadratic ftn,
 \perp can be calculated logically
as $H(\underline{x}^k)$, but calculation
of $\perp H(\underline{x}^k)$ should be avoided
for computational efficiency

①②

Extension to Non-quadratic functions

(avoid explicit calculation of Δ)

Recall

$$\underline{A} \underline{x}^{k+1} = (\underline{A}(\underline{x}^k + \alpha^k \underline{d}^k)) \quad \text{appearing in (d)}$$

use (b)

$$(\underline{A} \underline{d}^k) = \frac{1}{\alpha^k} (\underline{A} \underline{x}^{k+1} - \underline{A} \underline{x}^k)$$

$$= \frac{1}{\alpha^k} [(\underline{A} \underline{x}^{k+1} + \beta) - (\underline{A} \underline{x}^k + \beta)]$$

$$= \frac{1}{\alpha^k} (\underline{g}^{k+1} - \underline{g}^k) \quad (e)$$

Expressed only in term of
gradients

(e) \rightarrow (d) :

$$\beta^k = \frac{(\underline{g}^{k+1})^T (\underline{g}^{k+1} - \underline{g}^k)}{(\underline{g}^k)^T (\underline{g}^{k+1} - \underline{g}^k)} \quad (f)$$

To simplify (f), use

$$\left\{ \begin{array}{l} ① (\underline{g}^{k+1})^T \underline{d}^k = 0 \\ ② (\underline{d}^k)^T \underline{g}^k = -(\underline{g}^k)^T \underline{d}^k \\ ③ (\underline{g}^{k+1})^T \underline{g}^k = 0 \end{array} \right.$$

(13)

Then $\beta^k = \frac{(\underline{g}^{k+1})^T \underline{g}^{k+1} - (\underline{g}^k)^T \underline{g}^k}{(\underline{d}^k)^T \underline{g}^{k+1} - (\underline{d}^k)^T \underline{g}^k}$

$\therefore \boxed{\beta^k = \frac{(\underline{g}^{k+1})^T \underline{g}^k}{(\underline{g}^k)^T \underline{g}^k}}$

Proofs

proof of ①: $(\underline{g}^{k+1})^T \underline{d}^k = 0$

Recall: $f(\underline{x}) = f(\underbrace{\underline{x}^k}_{\text{known}} + \alpha \underbrace{\underline{d}^k}_{\text{known}})$

\rightarrow becomes min at $\alpha = \alpha^k$ such that

$$\begin{aligned} 0 &= \left. \frac{df}{d\alpha} \right|_{\alpha=\alpha^k} = \left[(\nabla f)^T \underline{d}^k \right]_{\alpha=\alpha^k} \\ &= \left[\nabla f (\underbrace{\underline{x}^k + \alpha^k \underline{d}^k}_{\underline{x}^{k+1}}) \right]^T \underline{d}^k \\ &= (\underline{g}^{k+1})^T \underline{d}^k \end{aligned}$$

(14)

$$\text{proof of } \textcircled{2} : (\underline{d}^k)^T \underline{g}^k = -(\underline{g}^k)^T \underline{g}^k$$

$$\begin{aligned} (\underline{d}^k)^T \underline{g}^k &= \underset{(c)}{=} \left[-\underline{g}^k + \beta^{k-1} \underline{d}^{k-1} \right]^T \underline{g}^k \\ &= -(\underline{g}^k)^T \underline{g}^k + \beta^{k-1} (\underline{d}^{k-1})^T \underline{g}^k \\ &= -(\underline{g}^k)^T \underline{g}^k \quad (\underline{g}^k)^T \underline{d}^{k-1} \stackrel{\text{by } \textcircled{1}}{=} 0 \end{aligned}$$

$$\text{proof of } \textcircled{3} : (\underline{g}^{k+1})^T \underline{g}^k = 0$$

$$\begin{aligned} (\underline{g}^{k+1})^T \underline{g}^k &= \underset{(c)}{=} (\underline{g}^{k+1})^T \left[-\underline{d}^k + \beta_{k-1} \underline{d}^{k-1} \right] \\ &= -\underbrace{(\underline{g}^{k+1})^T \underline{d}^k}_{\equiv 0 \text{ by } \textcircled{1}} + \beta_{k-1} (\underline{g}^{k+1})^T \underline{d}^{k-1} \\ &= \beta_{k-1} (\underline{g}^{k+1})^T \underline{d}^{k-1} \quad (\textcircled{2}) \end{aligned}$$

Now rewrite $\textcircled{2}$ as

$$\underline{g}^{k+1} = \underline{g}^k + \alpha^k \underline{A} \underline{d}^k \quad (\textcircled{3})$$

Then

$$\begin{aligned} \underline{g}^{k+1} &= \beta_{k-1} \left[\underbrace{\underline{g}^k + \alpha_{k-1} \underline{A} \underline{d}^k}_{(\textcircled{3})} \right]^T \underline{d}^{k-1} \\ &= \cancel{\beta_{k-1} (\underline{g}^k)^T \underline{d}^{k-1}}_{\text{by } \textcircled{1}} + \beta_{k-1} \alpha_{k-1} \cancel{(\underline{A} \underline{d}^k)^T \underline{d}^{k-1}}_{\text{by conjugate cond.}} = 0 \end{aligned}$$

(15)

< Fletcher-Reeve Conjugate Gradient Method >

- Start with any \underline{x}_0

- Search direction

$$\text{if } k = 0$$

$$\underline{d}^k = -\underline{g}^k = -\nabla f(\underline{x}^k)$$

else

$$\underline{d}^k = -\underline{g}^k + \beta^{k-1} \underline{d}^{k-1}$$

$$\beta^{k-1} = \frac{(\underline{g}^k)^T \underline{g}^k}{(\underline{g}^{k-1})^T \underline{g}^{k-1}} = \frac{\|\nabla f(\underline{x}^k)\|^2}{\|\nabla f(\underline{x}^{k-1})\|^2}$$

- Repeat until convergence

- Remark :
- ① this method applies to non-quadratic function minimization
 - ② n-iteration convergence is valid only for quadratic ftn
 - ③ may need to reset \underline{d}^k for every n step

Example

Consider $f = x_1^2 + 4x_2^2$, $\mathbf{x}_0 = (1,1)^T$

- the steepest descent iteration

$$\mathbf{d}_0 = -\nabla f(\mathbf{x}_0) = -(2, 8)^T$$

$$f(\alpha) = f(\mathbf{x}_0 + \alpha \mathbf{d}_0) \Rightarrow \alpha_0 = 0.1308$$

$$\therefore \mathbf{x}_1 = \mathbf{x}_0 + \alpha_0 \mathbf{d}_0 = (0.7385, -0.0462)^T$$

- the conjugate gradient iteration

$$\beta_0 = \frac{\|\nabla f(\mathbf{x}_1)\|^2}{\|\nabla f(\mathbf{x}_0)\|^2} = 0.0341$$

$$\mathbf{d}_1 = -\nabla f^T(\mathbf{x}_1) + \beta_0 \mathbf{d}_0 = \begin{pmatrix} -1.5451 \\ 0.0966 \end{pmatrix}$$

$$f(\alpha) = f(\mathbf{x}_1 + \alpha \mathbf{d}_1) \Rightarrow \alpha_1 = 0.4779$$

$$\therefore \mathbf{x}_2 = \mathbf{x}_1 + \alpha_1 \mathbf{d}_1 \cong (0.0, 0.0)^T$$

