

# Chapter 20

## Sequencing Method for the Mixed-Model Assembly Line to Realize Smoothed Production



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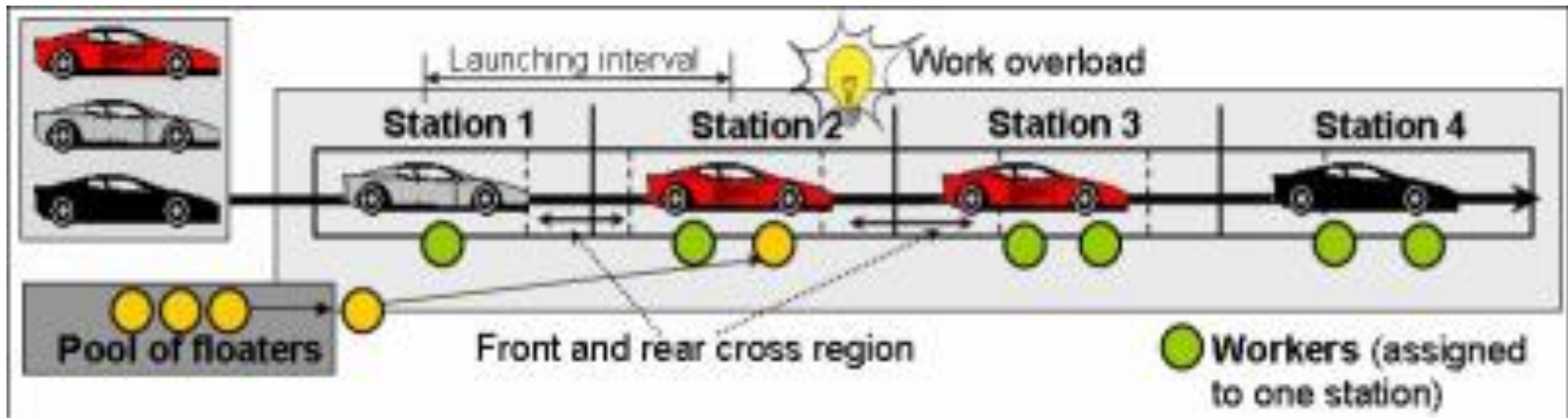
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Simultaneous achievement of two simplifying goals

# 20.1 GOALS OF CONTROLLING THE ASSEMBLY LINE

- ❖ A mixed-model includes several processes.
- ❖ It can operate different car models.
- ❖ It involves many types of parts.



# 20.1 GOALS OF CONTROLLING THE ASSEMBLY LINE

- ❖ Designing a mixed-model assembly line includes six steps.
  - Determine the cycle time.
  - Determine the minimum number of processes.
  - Prepare a diagram of integrated precedence relationship among elemental jobs.
  - Line balancing
  - Determine the sequence schedule for introducing various products to the line.
  - Determine the length of the operations range of each process.

# 20.1 GOALS OF CONTROLLING THE ASSEMBLY LINE

- ❖ There are two goals of controlling the mixed-model assembly line.
  - Leveling the total assembly time on each process within the line
  - Keeping a constant speed in consuming each part on the line
- ❖ Different goals require different sequences of introducing models to the line.

# 20.1 GOALS OF CONTROLLING THE ASSEMBLY LINE

- Indexes:

$i$  type of products (models)

$l$  process

- Parameters:

$\alpha$  number of types of products

$P$  number of processes

$Q_i$  planned production quantity of product  $i$  per day

$T_{il}$  unit operation time of product  $i$  on process  $l$

$C$  cycle time of the mixed-model assembly line

# 20.1 GOALS OF CONTROLLING THE ASSEMBLY LINE

- ❖ Goal one: **work load streamlining**
- ❖ A mixed-model assembly line has predetermined cycle time.
- ❖ Some products may have longer operation time than the cycle time.

# 20.1 GOALS OF CONTROLLING THE ASSEMBLY LINE

- ❖ In terms of averages, the total operation time for all models, each weighted by its production ratio, should satisfy

$$\max_{l=1, \dots, P} \left\{ \frac{\sum_{i=1}^{\alpha} Q_i T_{il}}{\sum_{i=1}^{\alpha} Q_i} \right\} \leq C$$

$$C = \text{cycle time} = \frac{\text{total operation time per day}}{\sum_{i=1}^{\alpha} Q_i}$$



# 20.1 GOALS OF CONTROLLING THE ASSEMBLY LINE

- ❖ Suppose a line that produces three models with two processes.

$$\alpha=3, P=2$$

$$Q_1=4 \text{ units}, \quad Q_2=3 \text{ units}, \quad Q_3=2 \text{ units}$$

$$T_{1,1}=2 \text{ hours/unit}, \quad T_{2,1}=5 \text{ hours/unit}, \quad T_{3,1}=2 \text{ hours/unit}$$

$$T_{1,2}=3 \text{ hours/unit}, \quad T_{2,2}=1 \text{ hour/unit}, \quad T_{3,2}=4 \text{ hours/unit}$$

# 20.1 GOALS OF CONTROLLING THE ASSEMBLY LINE

❖ For Process 1, the average operation time is

$$\frac{Q_1T_{1,1}+Q_2T_{2,1}+Q_3T_{3,1}}{Q_1+Q_2+Q_3} = \frac{4 \times 2 + 3 \times 5 + 2 \times 2}{4+3+2} = 3 \text{ hours/unit}$$

❖ For Process 2, the average operation time is

$$\frac{Q_1T_{1,2}+Q_2T_{2,2}+Q_3T_{3,2}}{Q_1+Q_2+Q_3} = \frac{4 \times 3 + 3 \times 1 + 2 \times 4}{4+3+2} = \frac{23}{9} \text{ hours/unit}$$

$$\max \left\{ 3, \frac{23}{9} \right\} = 3 \leq C$$

# 20.1 GOALS OF CONTROLLING THE ASSEMBLY LINE

- ❖ If products with relatively longer operation time are successively introduced into the line, there would be a delay in completing the product.

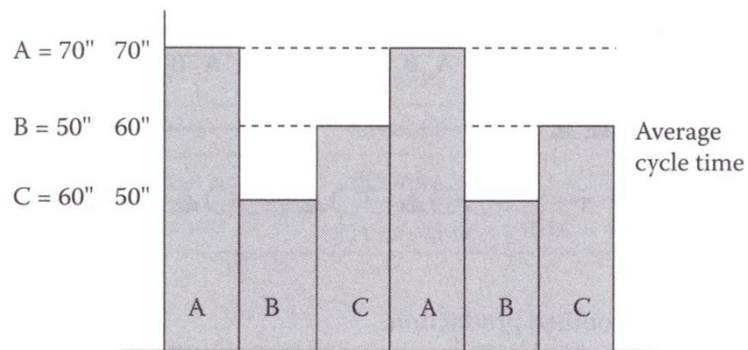


FIGURE 5.5  
Sequence schedule that enables assembly within the average cycle time.

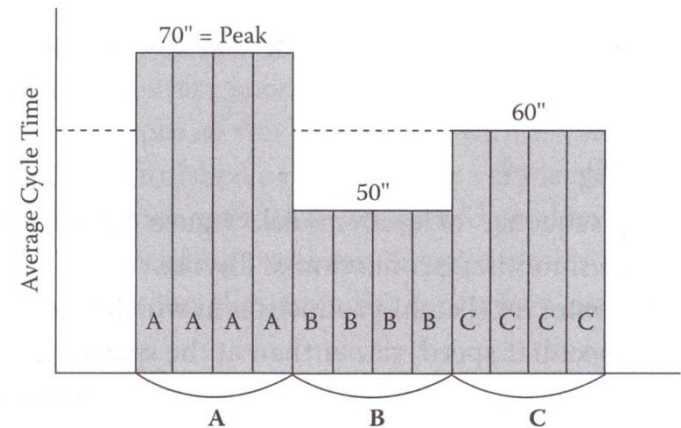


FIGURE 5.6  
Sequence schedule that causes line stoppage.

# 20.1 GOALS OF CONTROLLING THE ASSEMBLY LINE

Goal two: sequencing model for **parts usage streamlining**

- ❖ In the pull system, the variation in production quantities or conveyance times at **preceding processes** should be minimized.



- ❖ Therefore, the quantity used per hour for each part in the mixed-model line should be kept as constant as possible.

# 20.1 GOALS OF CONTROLLING THE ASSEMBLY LINE

- Parameters:

$Q$  total production quantity of all products =  $\sum_{i=1}^{\alpha} Q_i$

$\beta$  total number of types of parts

$N_j$  necessary quantity of part  $j$  for producing  $Q$

$X_{j,k}$  necessary quantity of part  $j$  for producing the products of determined sequence from first to  $k$ th

$N_j/Q$  average necessary quantity of part  $j$  for one product

$kN_j/Q$  average necessary quantity of part  $j$  for  $k$  products

# 20.1 GOALS OF CONTROLLING THE ASSEMBLY LINE

- ❖ To keep the consumption speed of part  $j$  constant, the value of  $X_{j,k}$  should be as close as possible to the value of  $kN_j/Q$ ,  $k=1,2,\dots,Q$ .

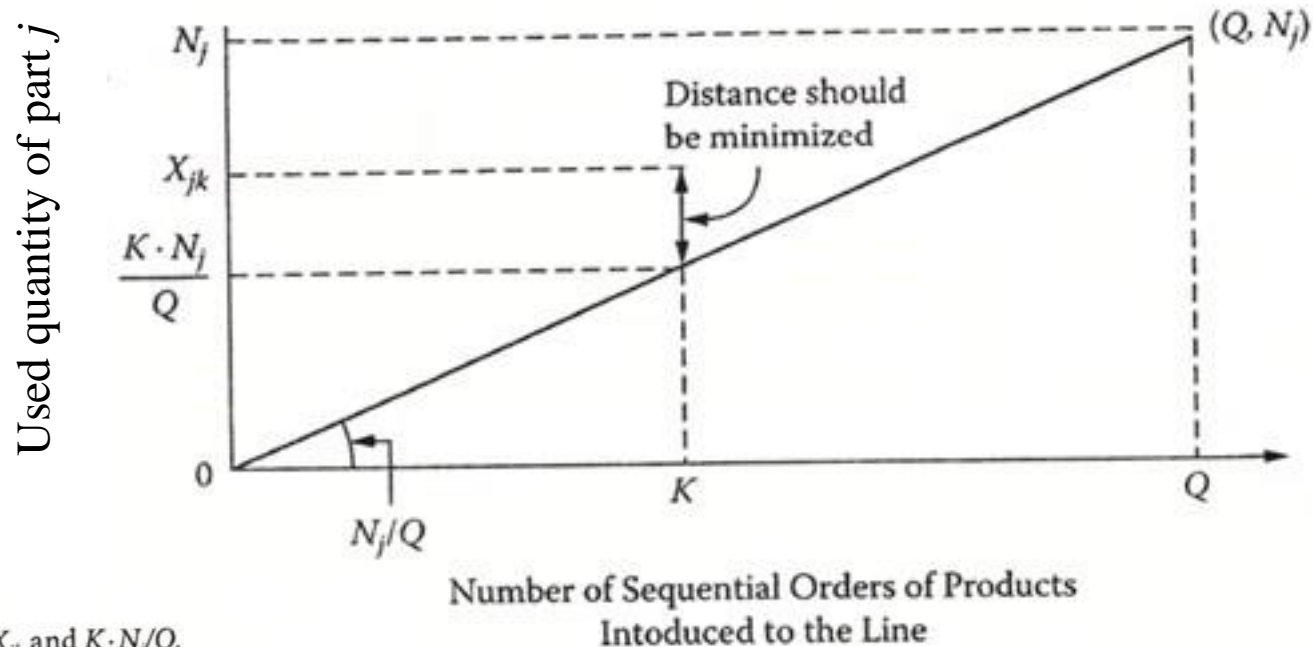


FIGURE 20.1  
Relationship between  $X_{jk}$  and  $K \cdot N_j/Q$ .

# 20.1 GOALS OF CONTROLLING THE ASSEMBLY LINE

❖ It can be further defined that

Independent of the sequence

$$\text{A point } G_k = \left( kN_1 / Q, kN_2 / Q, \dots, kN_\beta / Q \right)$$

$$\text{A point } P_k = \left( X_{1,k}, X_{2,k}, \dots, X_{\beta,k} \right)$$

Dependent on the sequence

# 20.1 GOALS OF CONTROLLING THE ASSEMBLY LINE

- ❖ If a sequence schedule assures the constant speed of consuming each part, point  $P_k$  must be as close as possible to point  $G_k$ .
- ❖ Let the distance be  $D_k$ , we have

$$\min D_k = \|G_k - P_k\| = \sqrt{\sum_{j=1}^{\beta} \left( \frac{kN_j}{Q} - X_{j,k} \right)^2}$$



## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

$b_{ij}$  necessary quantity of part  $j$  for product  $i$

$$i=1,2,\dots,\alpha; j=1,2,\dots,\beta$$

### ❖ Algorithm

Step 1: set  $k=1$ ,  $X_{j,0}=0$ , and  $S_0=\{1,2,\dots,\alpha\}$

## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

### ❖ Algorithm

Step 2: Set product  $i^*$  as the  $k$ th product in the sequence schedule that minimizes  $D_k$ .

$$D_{ki^*} = \min_{i \in S_{k-1}} \{ D_{ki} \}$$

$$D_{ki} = \sqrt{\sum_{j=1}^{\beta} \left( \frac{kN_j}{Q} - X_{j,k} \right)^2} = \sqrt{\sum_{j=1}^{\beta} \left( \frac{kN_j}{Q} - X_{j,k-1} - b_{ij} \right)^2}$$

## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

### ❖ Algorithm

Step 3: If all of product  $i^*$  have been included in the sequence, set  $S_k = S_{k-1} - \{i^*\}$ ;  
otherwise, set  $S_k = S_{k-1}$ .

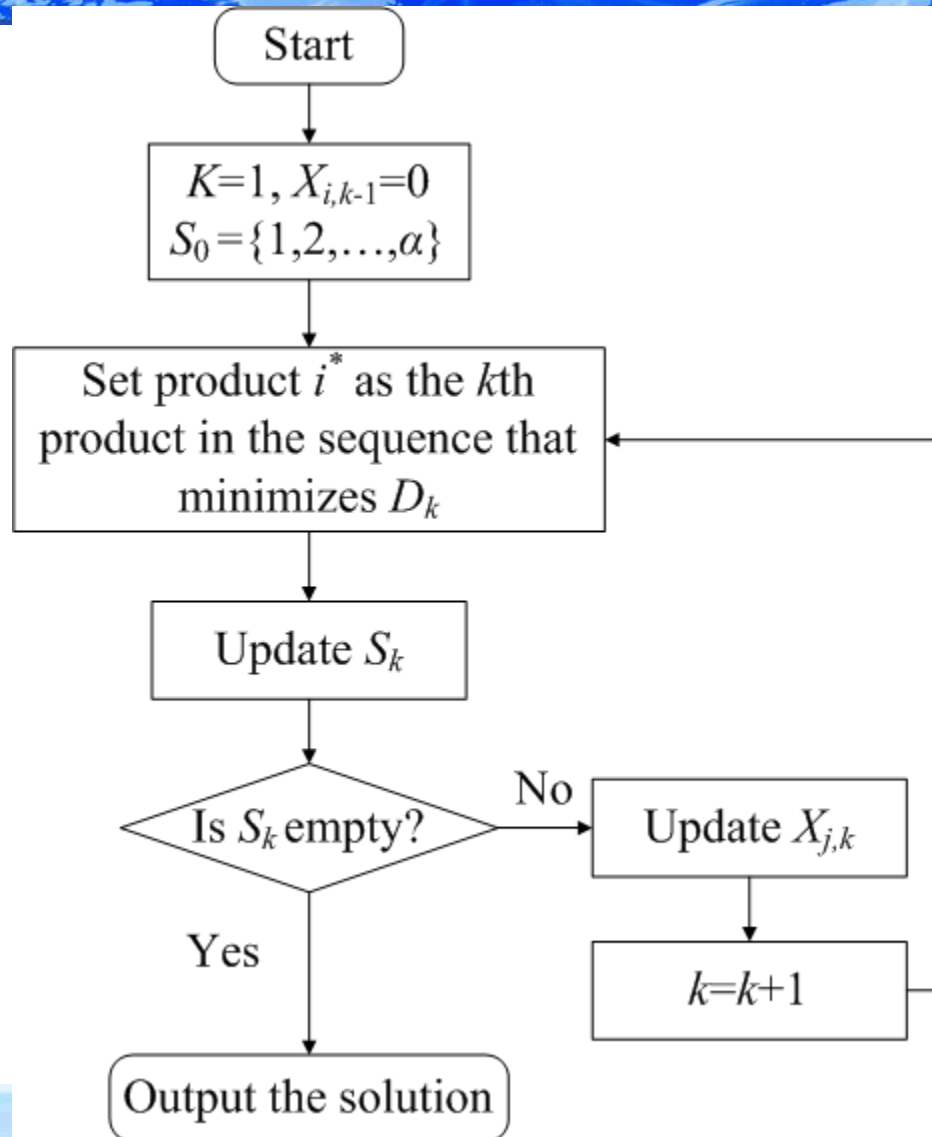
## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

### ❖ Algorithm

Step 4: If  $S_k = \emptyset$ , stop.

If  $S_k \neq \emptyset$ , then set  $X_{i,k} = X_{i,k-1} + b_{i*j}, j=1,2,\dots,\beta$ .  
set  $k = k+1$ ;  
go to Step 2.

# 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE



# 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

Denote:

$b_{ij}$  = Necessary quantity of the part  $a_i$  ( $j = 1, \dots, \beta$ ) for producing one unit of the product  $A_i$  ( $i = 1, \dots, \alpha$ ).

Other notations are already defined.

Then,

*Step 1* Set  $K = 1$ ,  $X_{j,k-1} = 0$ , ( $j = 1, \dots, \beta$ ),  $S_{k-1} = (1, 2, \dots, \alpha)$ .

*Step 2* Set as  $K$ th order in the sequence schedule the product  $A_{i^*}$  which minimizes the distance  $D_k$ . The minimum distance will be found by the following formula:

$$D_{ki^*} = \min\{D_{ki}\}, i \in S_{k-1} \text{ (i)} \rightarrow \text{should be removed}$$

where  $D_{ki} = \sqrt{\sum_{j=1}^{\beta} \left( \frac{K \cdot N_j}{Q} - X_{j,k-1} - b_{ij} \right)^2}$ .

should be "all"

*Step 3* If **some** units of product  $A_{i^*}$  were ordered and included in the sequence schedule, then

$$\text{Set } S_k = S_{k-1} - \{i^*\}.$$

If some units of a product  $A_{i^*}$  are still remaining as being not ordered, then set  $S_k = S_{k-1}$ .

*Step 4* If  $S_k = \emptyset$  (empty set), the algorithm will end.

If  $S_k \neq \emptyset$ , then compute  $X_{jk} = X_{j,k-1} + b_{i^*j}$  ( $j = 1, \dots, \beta$ ) and go back to Step 2 by setting  $K = K + 1$ .

FIGURE 20.2  
Goal-chasing method I.

# 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

## ❖ An example

**TABLE 20.1**

Production Quantities  $Q_i$  and Parts Condition  $b_{ij}$ \*

	Product $A_i$		
	$A_1$	$A_2$	$A_3$
Planned Production Quantity $Q_j$	2	3	5

Products $A_i$	Parts $a_j$			
	$a_1$	$a_2$	$a_3$	$a_4$
$A_1$	1	0	1	1
$A_2$	1	1	0	1
$A_3$	0	1	1	0

## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

❖ The total production quantity

$$\sum_{i=1}^3 Q_i = 2 + 3 + 5 = 10$$

$$[N_j] = [Q_i] [b_{ij}]$$

$$= [2, 3, 5] \begin{bmatrix} 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix} = [5, 8, 7, 5]$$



## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

$$\left[ N_j / Q \right] = \left[ \frac{5}{10}, \frac{8}{10}, \frac{7}{10}, \frac{5}{10} \right] \quad j = 1, 2, 3, 4$$

- ❖ Next, apply the values of  $[N_j / Q]$  and  $[b_{ij}]$  to the formula in Step 2.

## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

❖ When  $k = 1$ ,  $D_{ki}$  can be calculated as:

$$D_{1,1} = \sqrt{\left(\frac{1 \times 5}{10} - 0 - 1\right)^2 + \left(\frac{1 \times 8}{10} - 0 - 0\right)^2 + \left(\frac{1 \times 7}{10} - 0 - 1\right)^2 + \left(\frac{1 \times 5}{10} - 0 - 1\right)^2}$$
$$= 1.11$$

$$D_{1,2} = 1.01 \quad D_{1,3} = 0.79$$

$$D_{1,i^*} = \min \{1.11, 1.01, 0.79\} = 0.79$$

$$i^* = 3$$

## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

- ❖ Therefore, the first product introduced into the line is Product 3.

$$X_{j,k} = X_{j,k-1} + b_{3j}$$

$$X_{1,1} = 0 + 0 = 0 \quad X_{2,1} = 0 + 1 = 1$$

$$X_{3,1} = 0 + 1 = 1 \quad X_{4,1} = 0 + 0 = 0$$

## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

❖ When  $k = 2$ ,  $D_{ki}$  can be calculated as:

$$D_{2,1} = \sqrt{\left(\frac{2 \times 5}{10} - 0 - 1\right)^2 + \left(\frac{2 \times 8}{10} - 1 - 0\right)^2 + \left(\frac{2 \times 7}{10} - 1 - 1\right)^2 + \left(\frac{2 \times 5}{10} - 0 - 1\right)^2}$$
$$= 0.85$$

$$D_{2,2} = 0.57 \quad D_{2,3} = 1.59$$

$$D_{2,i^*} = \min \{0.85, 0.57, 1.59\} = 0.57$$

$$i^* = 2$$

## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

❖ Therefore, the second product introduced into the line is Product 2.

$$X_{j,k} = X_{j,k-1} + b_{2j}$$

$$X_{1,2} = 0 + 1 = 1$$

$$X_{2,2} = 1 + 1 = 2$$

$$X_{3,2} = 1 + 0 = 1$$

$$X_{4,2} = 0 + 1 = 1$$

# 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

❖ By repeating the above steps, we can obtain the solution.

$K$	$D_{k1}$	$D_{k2}$	$D_{k3}$	Sequence Schedule	$X_{1k}$	$X_{2k}$	$X_{3k}$	$X_{4k}$
1	1.11	1.01	0.79	$A_3$	0	1	1	0
2	0.85	0.57*	1.59	$A_3 A_2$	1	2	1	1
3	0.82*	1.44	0.93	$A_3 A_2 A_1$	2	2	2	2
4	1.87	1.64	0.28*	$A_3 A_2 A_1 A_3$	2	3	3	2
5	1.32	0.87*	0.87	$A_3 A_2 A_1 A_3 A_2$	3	4	3	3
6	1.64	1.87	0.28*	$A_3 A_2 A_1 A_3 A_2 A_3$	3	5	4	3
7	0.93	1.21	0.82*	$A_3 A_2 A_1 A_3 A_2 A_3 A_3$	3	6	5	3
8	0.57*	0.85	1.59	$A_3 A_2 A_1 A_3 A_2 A_3 A_3 A_1$	4	6	6	4
9	1.56	0.77*	1.01	$A_3 A_2 A_1 A_3 A_2 A_3 A_3 A_1 A_2$	5	7	6	5
10	—	—	0*	<u><math>A_3 A_2 A_1 A_3 A_2 A_3 A_3 A_1 A_2 A_3</math></u>	5	8	7	5

FIGURE 20.3

Sequence schedule. (Note: \* indicates smallest distance  $D_{kj}$ .)

## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

- ❖ Suppose that  $[[kN_j/Q]]$  denotes the integer which is the closest to  $kN_j/Q$ .
- ❖ If  $X_{j,k} = [[kN_j/Q]]$ , the optimality is achieved in part  $j$  at the  $k$ th position of the sequence.

# 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

❖ The values of  $kN_j/Q$  and  $X_{j,k}$  for part  $j$  are depicted as follows.

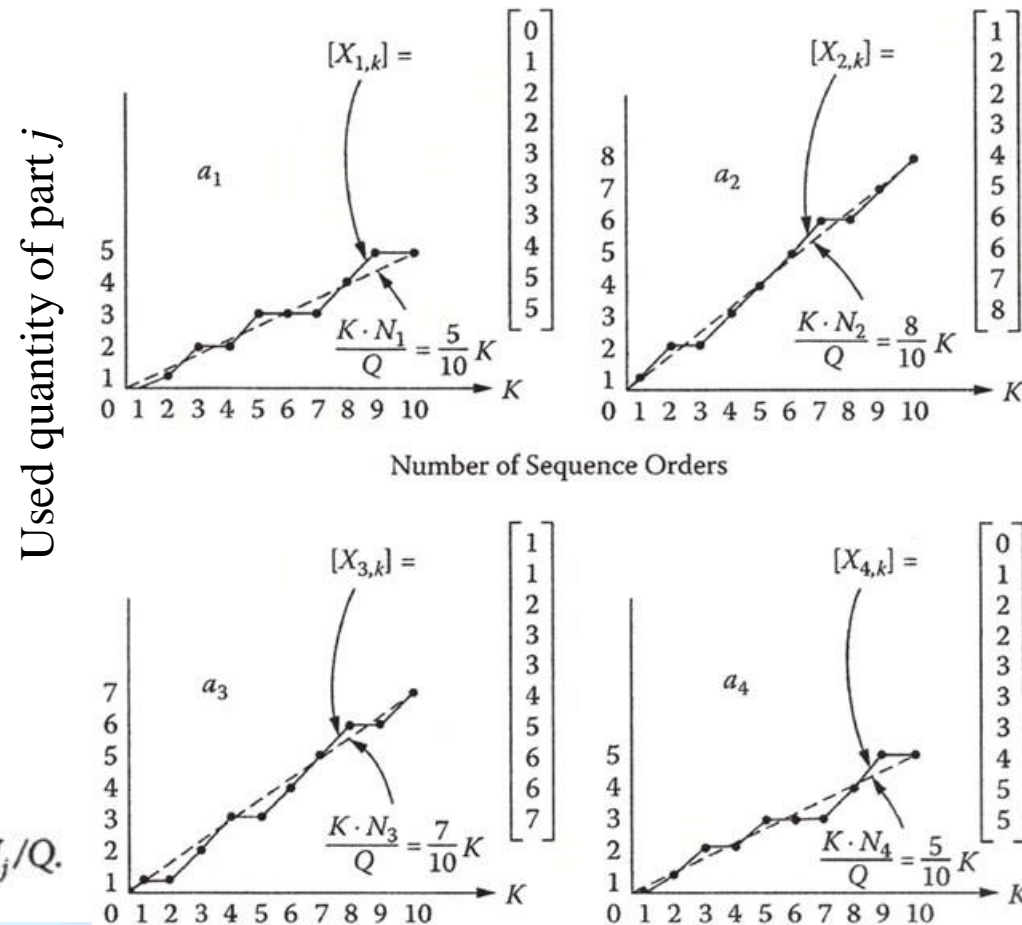


FIGURE 20.4

How  $X_{j,k}$  approached  $K \cdot N_j/Q$ .



## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

- ❖ To further evaluate this algorithm, we calculate the mean and standard deviation of the following values.

$$\frac{kN_j}{Q} - X_{j,k} \quad k = 1, 2, \dots, Q$$

## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

- ❖ When the number of varieties in parts and/or the number of varieties in models are **increased**, both the mean and the standard deviation are **increased**.
- ❖ When the production quantity is **increased**, both the mean and the standard deviation are **decreased**.

The more the tendency to produce multi-varieties in each small quantity is promoted, the less likely smoothing of production will be attained.

## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

- ❖ Consider the case with a large total production quantity (e.g.,  $Q=1000$ ).
- ❖ The sequence determined by the algorithm can be divided into 16 equal ranges, each of which corresponds to one hour.

## 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

- ❖ The quantity of each part in each range are computed.
- ❖ The standard deviation of these quantities are computed.

# 20.2 GOAL-CHASING METHOD: A NUMERICAL EXAMPLE

❖ The values of  $\sigma$  are fairly small. The coefficient of variation is also small.  $(\sigma / \bar{x})$

Kind of front axes \ Range	1				2				3				4				$\bar{x}$	$\sigma$
	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16		
a <sub>1</sub>	9	7	7	9	8	7	8	8	8	8	7	8	9	7	7	8	7.8	0.73
a <sub>2</sub>	6	5	7	6	5	6	7	5	7	6	5	7	6	6	5	6	5.9	0.75
a <sub>3</sub>	5	6	5	5	6	6	4	6	4	6	6	5	4	6	5	6	5.3	0.77
a <sub>4</sub>	3	3	3	2	3	3	3	3	3	2	3	3	3	2	3	3	2.8	0.33
a <sub>5</sub>	2	2	2	2	3	2	2	2	2	3	2	1	3	2	2	2	2.1	0.48
a <sub>6</sub>	1	1	1	1	1	2	1	1	2	0	2	1	1	1	1	1	1.1	0.48

**FIGURE 20.5**

Distribution of each kind of front axle used.

## 20.3 SIMULTANEOUS ACHIEVEMENT OF TWO SIMPLIFYING GOALS

- ❖ It is important to avoid successive processing of the products that have larger load of assembly time.
- ❖ Products have different loads at different processes.

## 20.3 SIMULTANEOUS ACHIEVEMENT OF TWO SIMPLIFYING GOALS



The car models with longer assembly time



Larger loads at every process

- ❖ Avoid introducing successively the same product requiring longer assembly time to the line.

## 20.3 SIMULTANEOUS ACHIEVEMENT OF TWO SIMPLIFYING GOALS

- ❖ All models are classified according to large ( $a_l$ ), medium ( $a_m$ ), or small ( $a_s$ ) total assembly time.
- ❖ Each model of  $a_j$  must be introduced to the line to keep the speed constant,  $j=l, m, s$ .



## 20.3 SIMULTANEOUS ACHIEVEMENT OF TWO SIMPLIFYING GOALS

- ❖ In practice, Toyota “weights” important subassemblies and provides constraints, such as facility capacities .
- ❖ The categories  $(a_l, a_m, a_s)$  are also given some weight to solve the conflict between the **line balancing goal** and the **part smoothing goal**.