



# Transport Layer

## - High Performance Transport -

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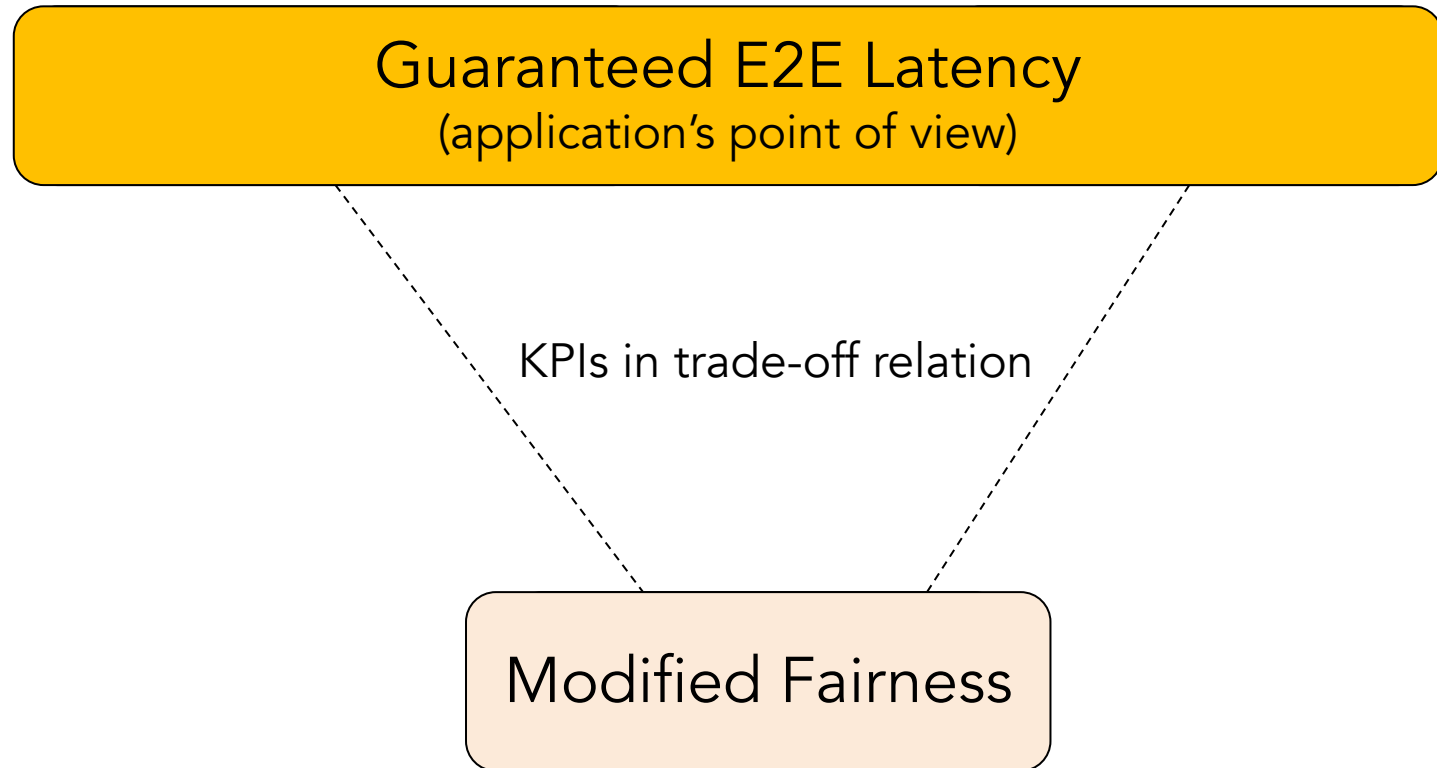
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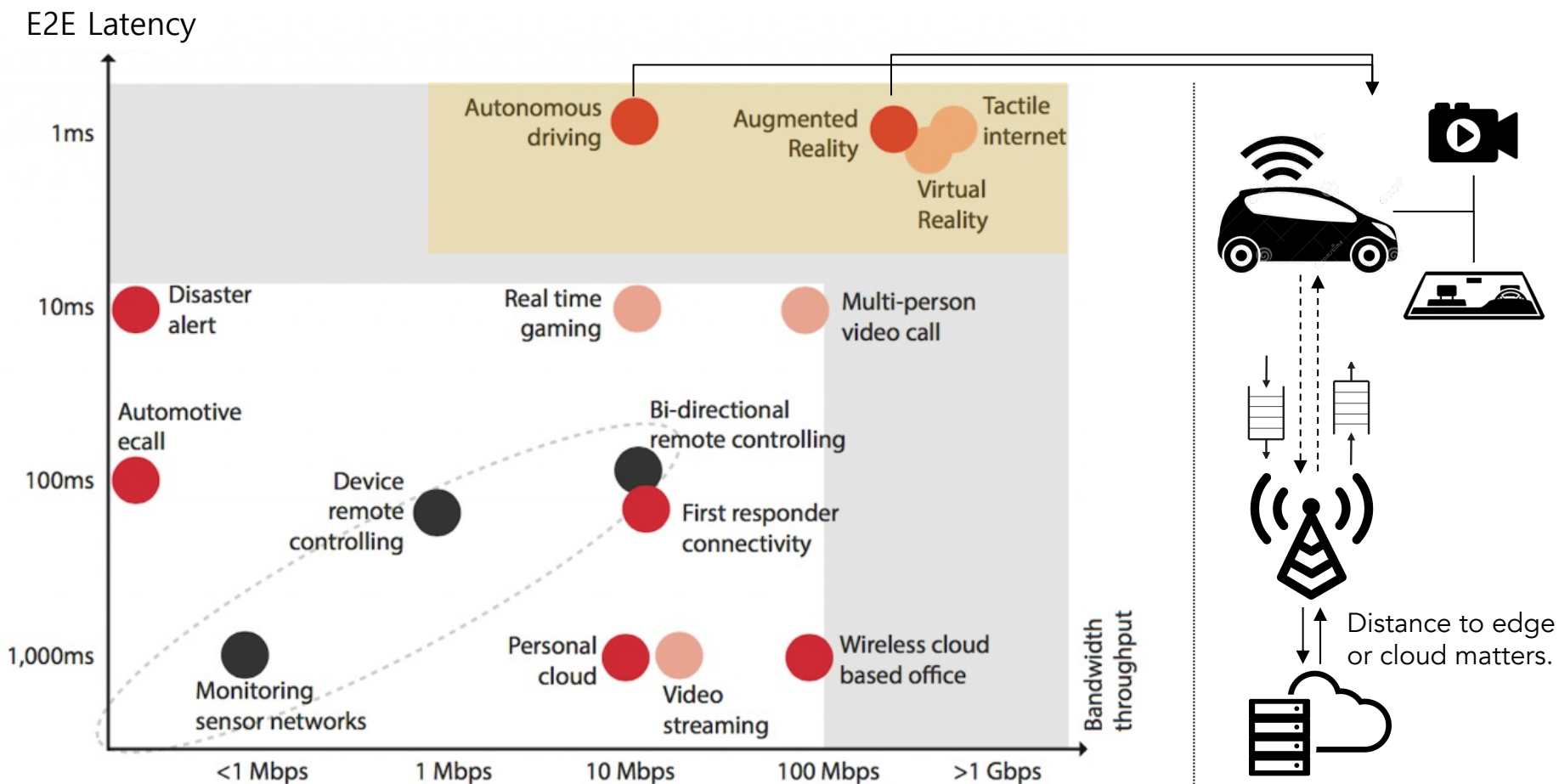


# KPIs for Future Transport Layer Protocols



# E2E Latency is Everything

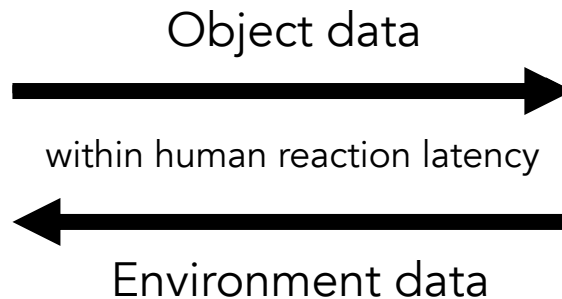
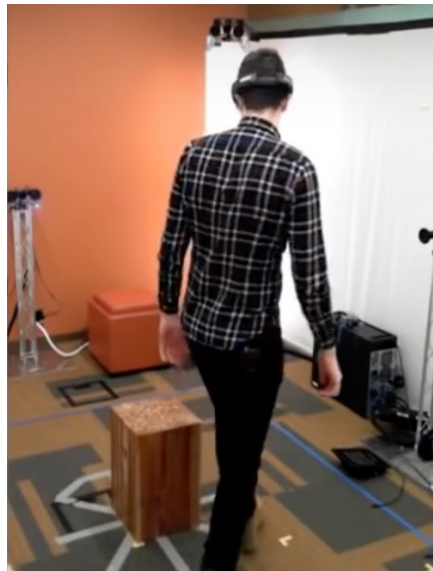
$$\text{E2E Latency} = \text{First-byte-Delay} + \text{DataVolume/Throughput}$$



Mark Glimour, "5G – Latency: New use cases and the need for network slicing," InformaTech reports, Feb. 2017.

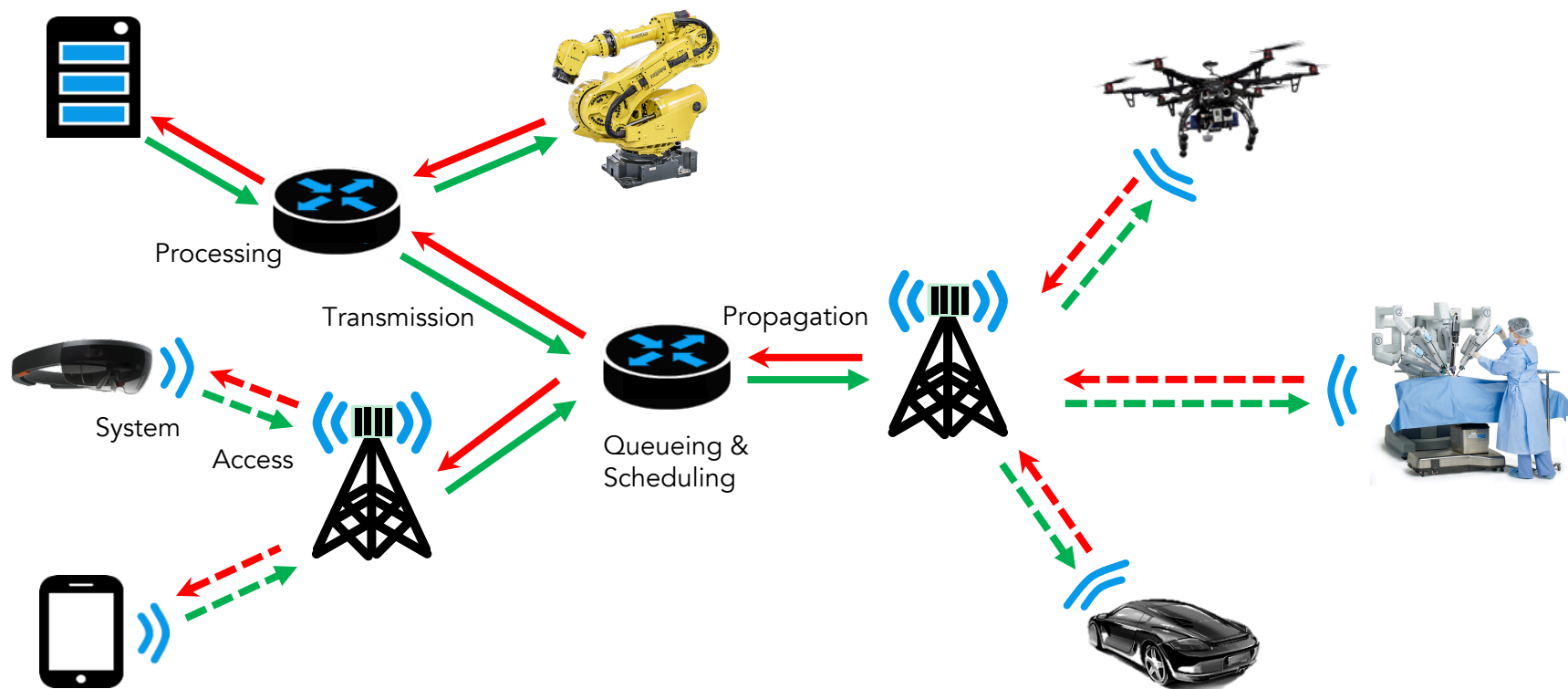
# Ultimate Service: Low-latency Telepresence

- **Low-latency telepresence** is the idea of providing people the experience as if they are present at a place other than their physical location **by making remote interactions indistinguishable from local interactions.**





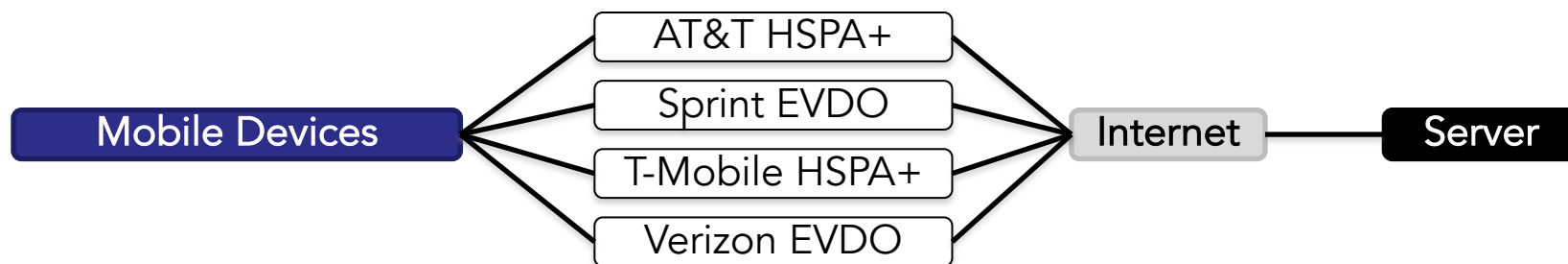
# Latency Decomposition



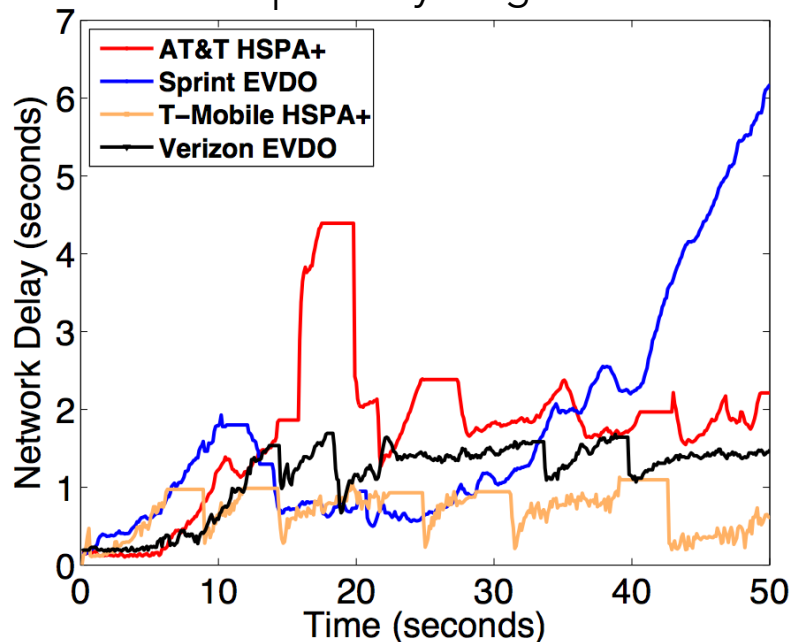
| App Delay                          | Transmission Delay<br>(1 us – 100ms) | Processing Delay<br>(< 1 us) | Scheduling Delay<br>(< 1 ms) | Queueing Delay<br>(< 500 ms) | Propagation Delay<br>(1 ms @ 200km) |
|------------------------------------|--------------------------------------|------------------------------|------------------------------|------------------------------|-------------------------------------|
| Socket control,<br>Encoder control | Redundancy<br>elimination            | FPGA,<br>ASIC                | Short TTI,<br>Priority sched | Congestion<br>control        | Untouchable                         |

# Latency Reality in Cellular Networks

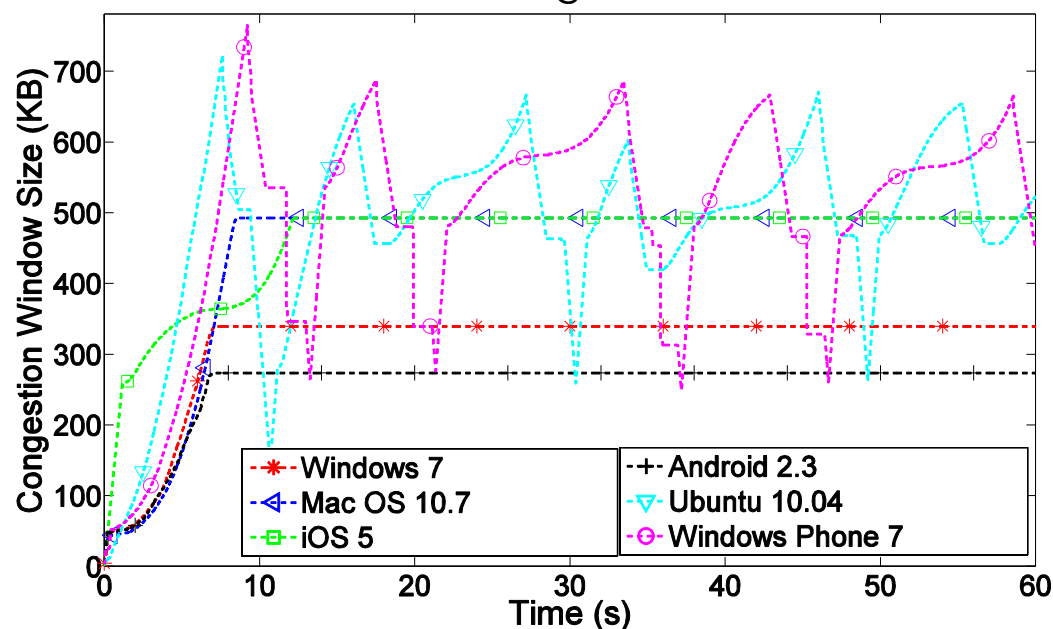
Tackling bufferbloat in 3G/4G networks, ACM IMC 2012



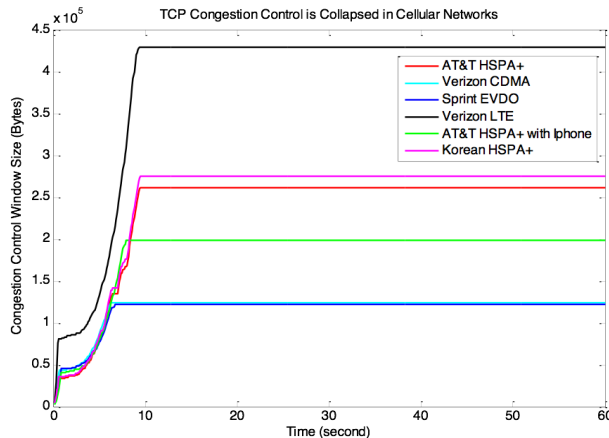
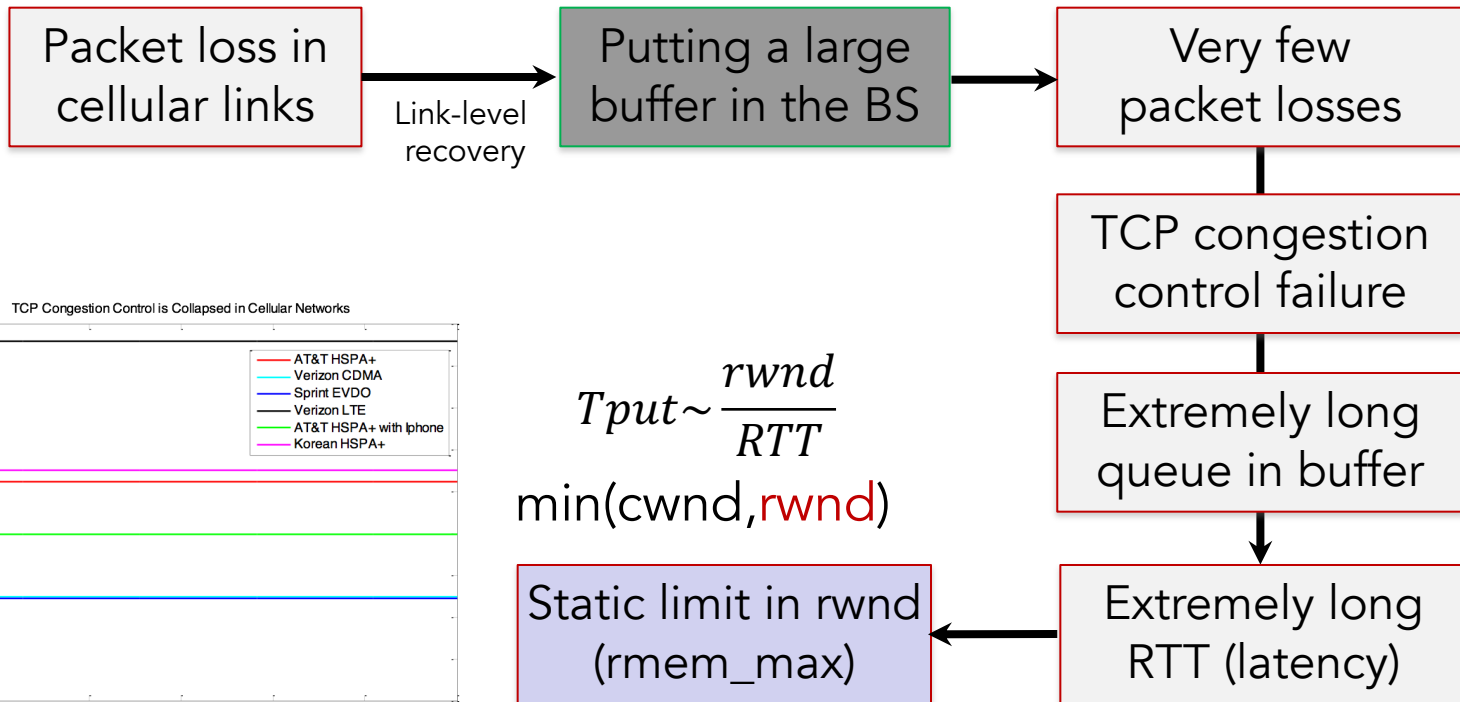
Exceptionally long RTT



Flat TCP congestion window



# Understanding the Problem and the Trick



$$T_{put} \sim \frac{rwnd}{RTT} \min(cwnd, rwnd)$$

Static limit in rwnd (rmem\_max)

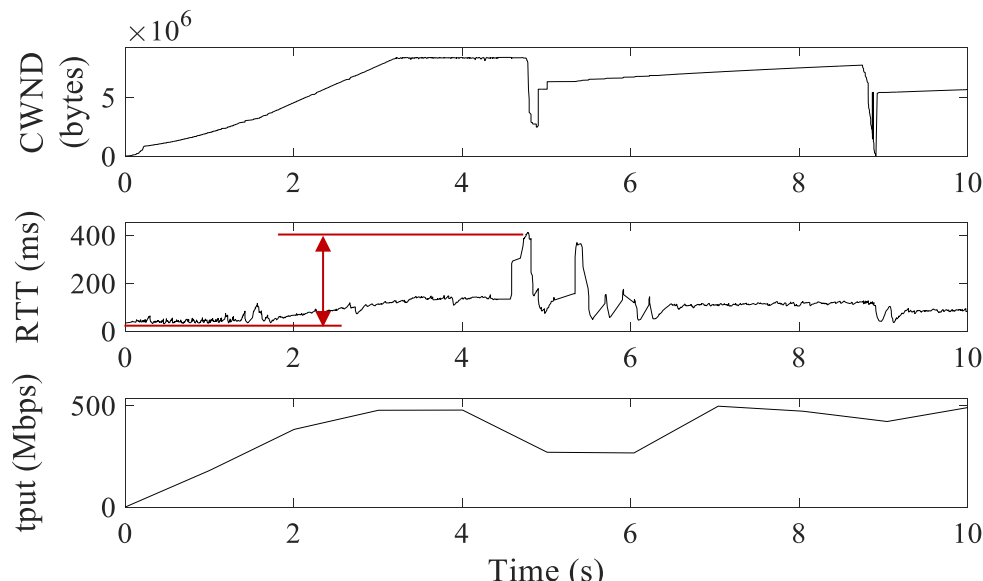
A potential solution (delay-based control)

DRWA (Dynamic Receive Window Adjustment)

$$T_{put} \sim \frac{rwnd(\Delta RTT)}{RTT}$$

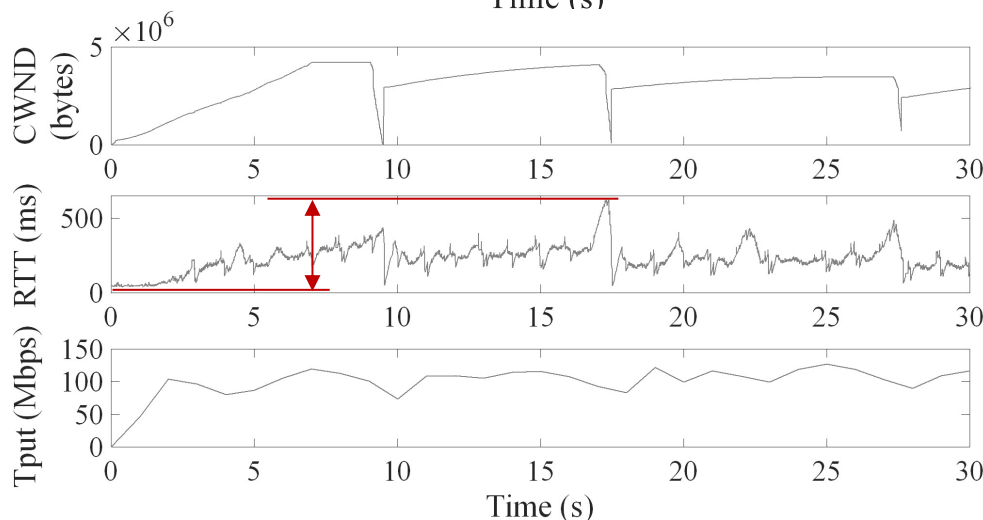
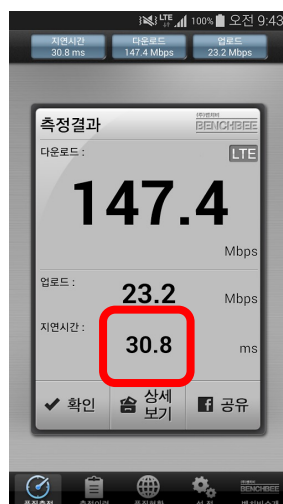
# Latency Reality in 5G/LTE Networks

5G  
(NSA)



Mean, Min RTT  
118.15, 35.28 ms  
(~3x inflation)

LTE

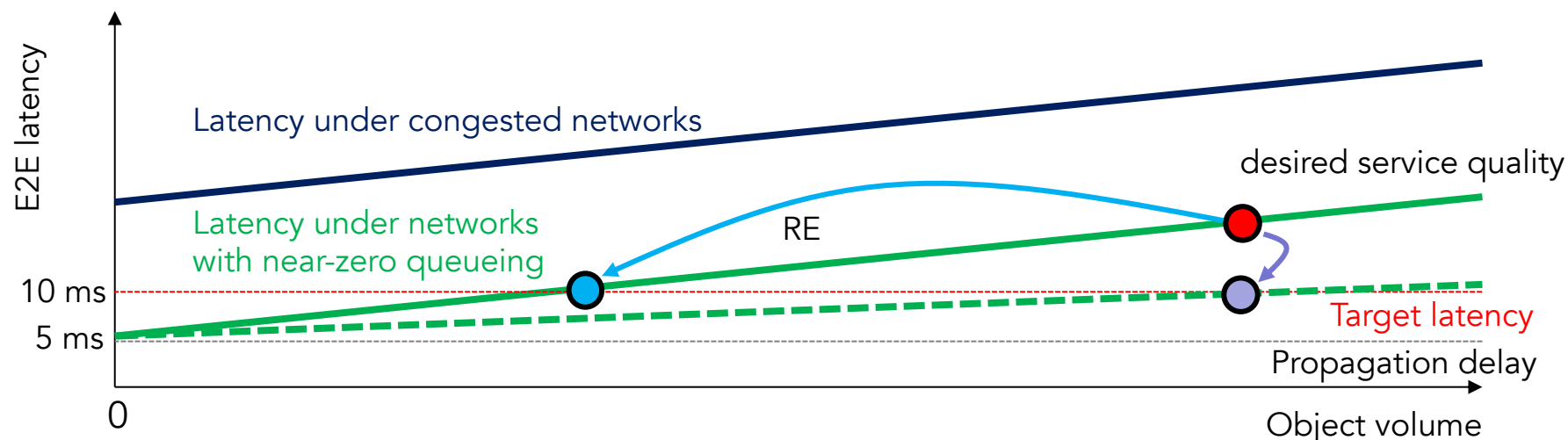


Mean, Min RTT  
226.10, 37.05 ms  
(~6x inflation)

# Core Challenges for Guaranteed E2E Latency

E2E latency: time taken to relocate an information object over a network  
(mostly Propagation delay + Queueing delay + Transmission delay + Retransmission delay)

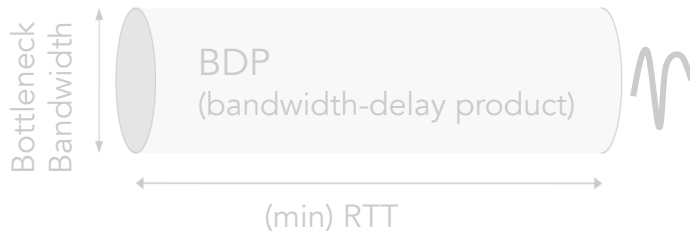
- $$= \frac{\text{volume} \downarrow}{\text{data rate} \uparrow}$$
- 1. Low-latency transport protocol (congestion control)
  - 2. Redundancy elimination (socket)
  - 3. Fast Convergence (+ PHY evolution)
  - 4. Zero retransmission (socket)



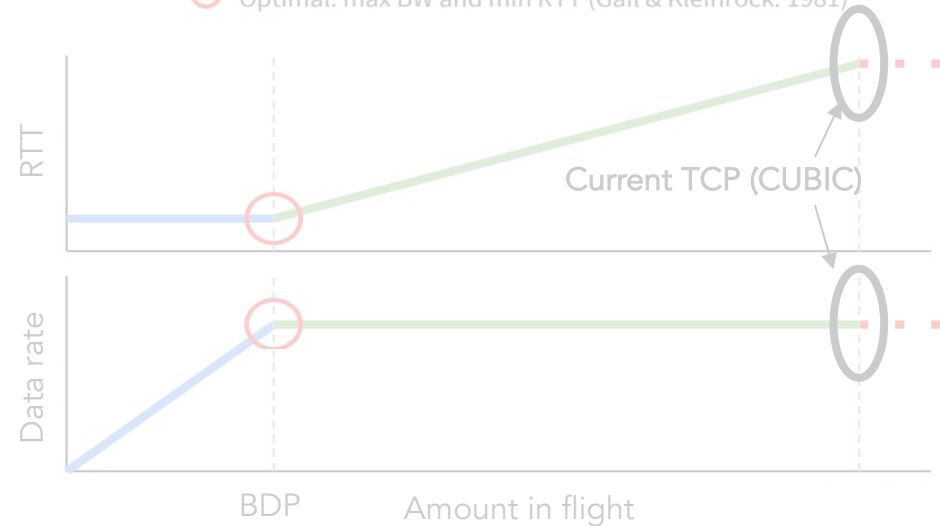
# Challenge 1+3: Low-Latency CC (+ Fast Convergence)



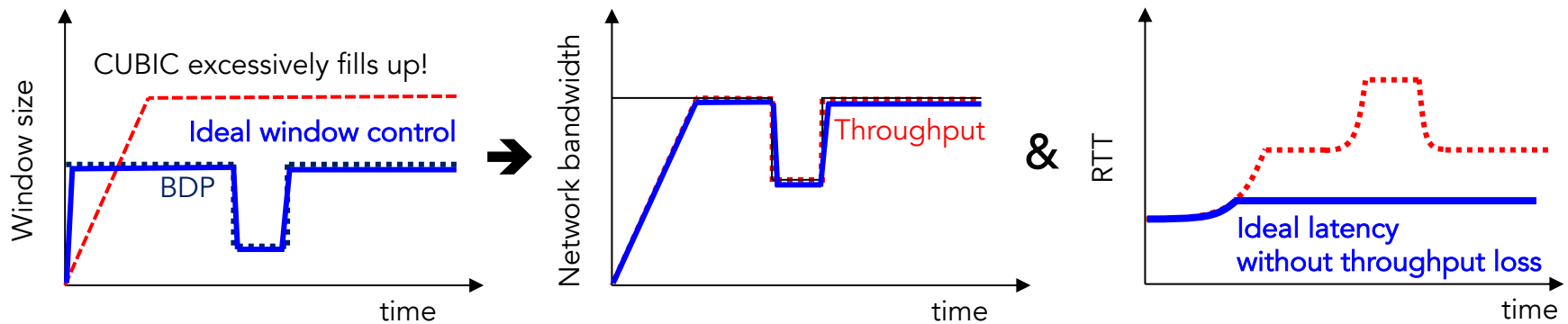
Transport protocol controls "windows (W)"  
Data rate ( $W/RTT$ ) is the result of sending W over RTT.



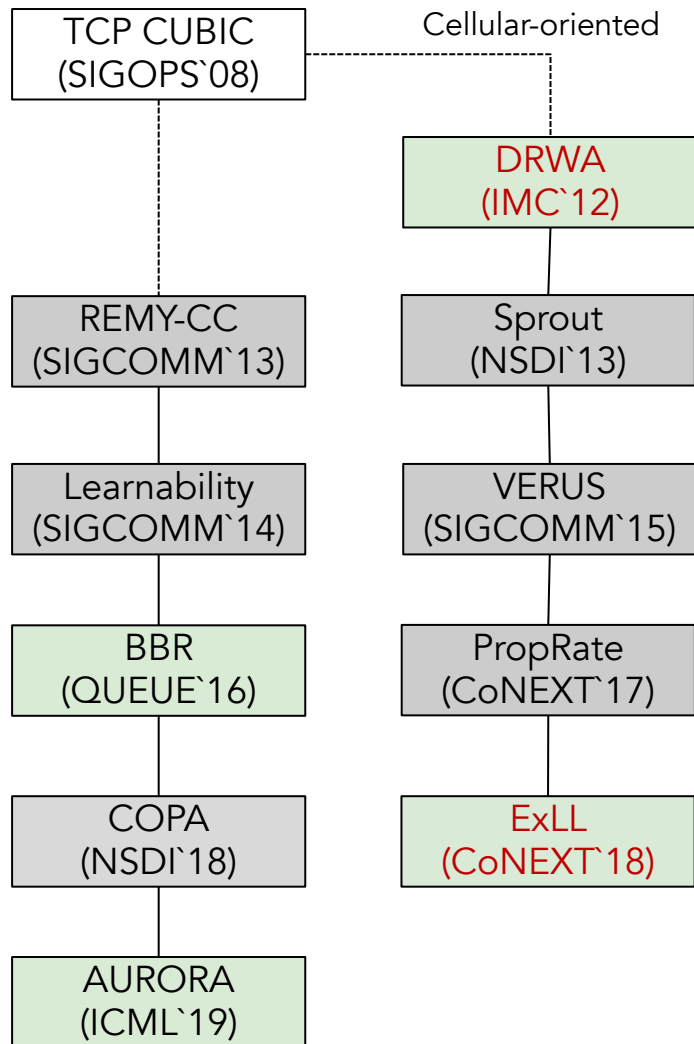
○ Optimal: max BW and min RTT (Gail & Kleinrock, 1981)



RTT in a cellular network with fluctuating bandwidth

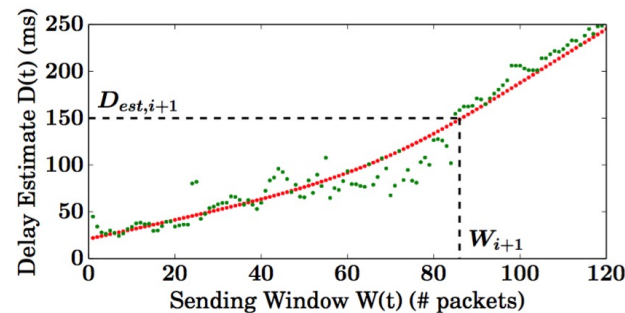


# Low-Latency Transport Protocols

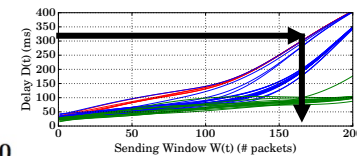


## VERUS

builds delay profiles from its window controls



Rebuild profile in every 1s

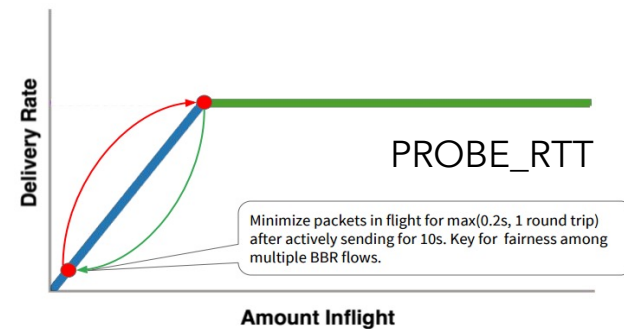
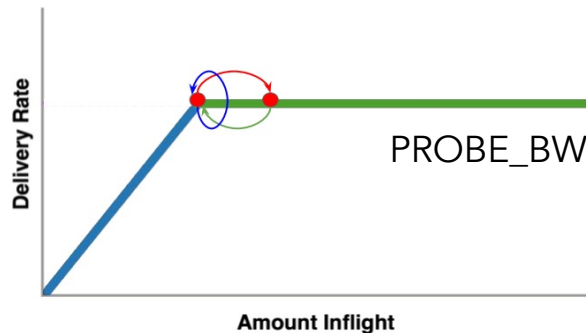
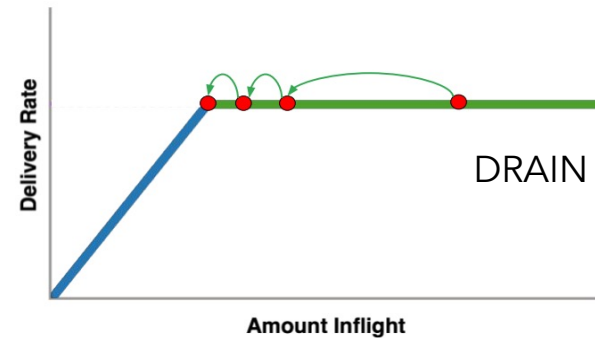
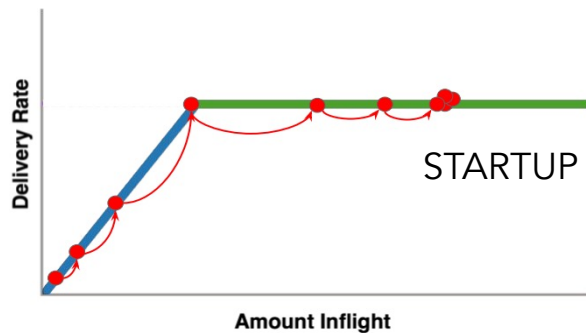


Reduce window if max delay out of a bound

$$D_{est,i+1} = \begin{cases} D_{est,i} - \delta_2 & \text{if } \frac{D_{max,i}}{D_{min}} > R \\ \text{Given more window budget if max delay within the bound} \\ D_{est,i} + \delta_2 & \text{otherwise} \end{cases}$$

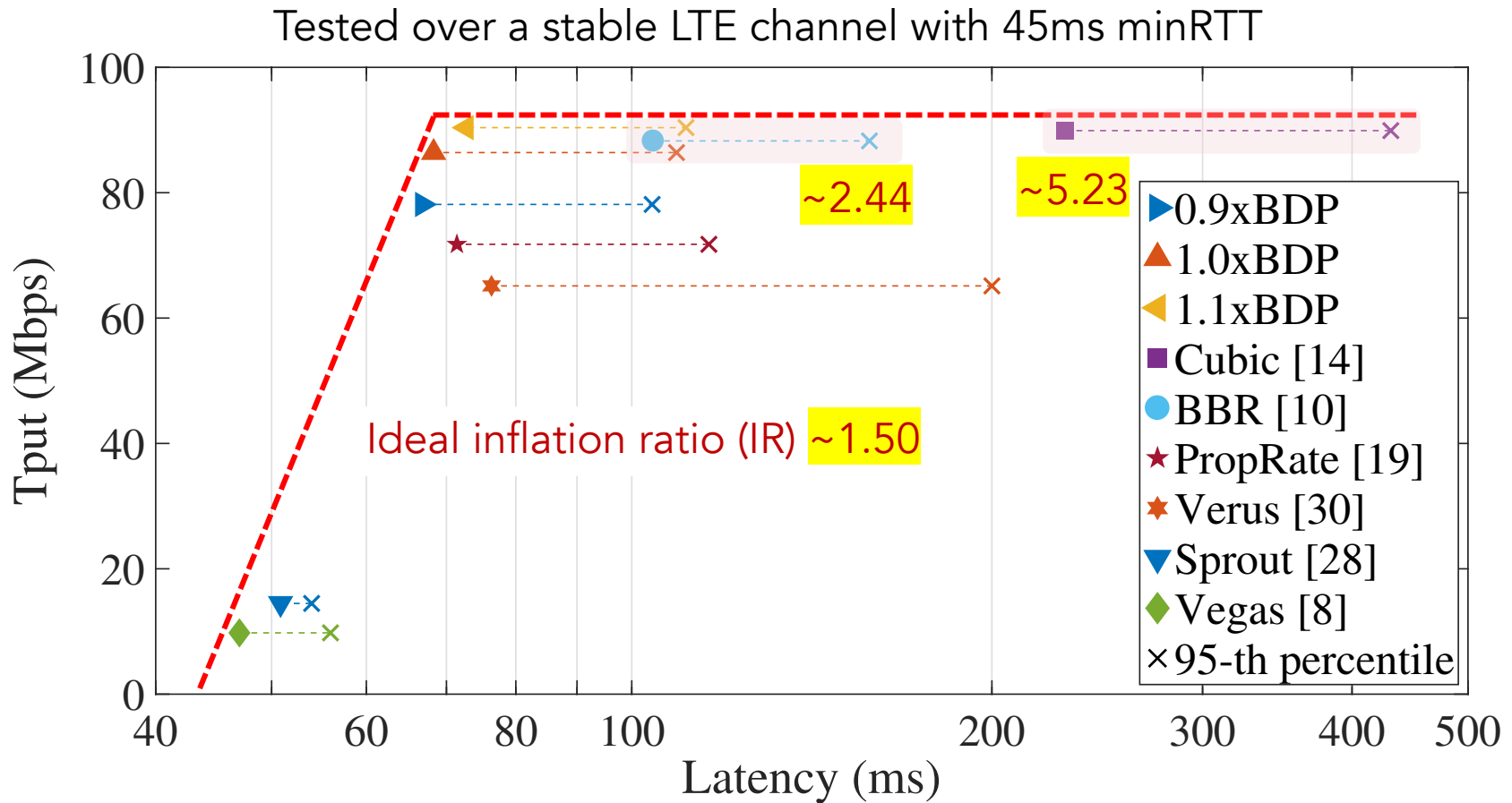
# BBR (Bottleneck Bandwidth and Round-trip propagation time)

- BBR cycles 4 operations (BBR is in the Linux kernel from 4.9)
  - STARTUP - exponential BW search
  - DRAIN - drain the queue created during startup
  - PROBE\_BW - explore max BW, drain queue, cruise
  - PROBE\_RTT - guarantee fairness among multiple BBR flows



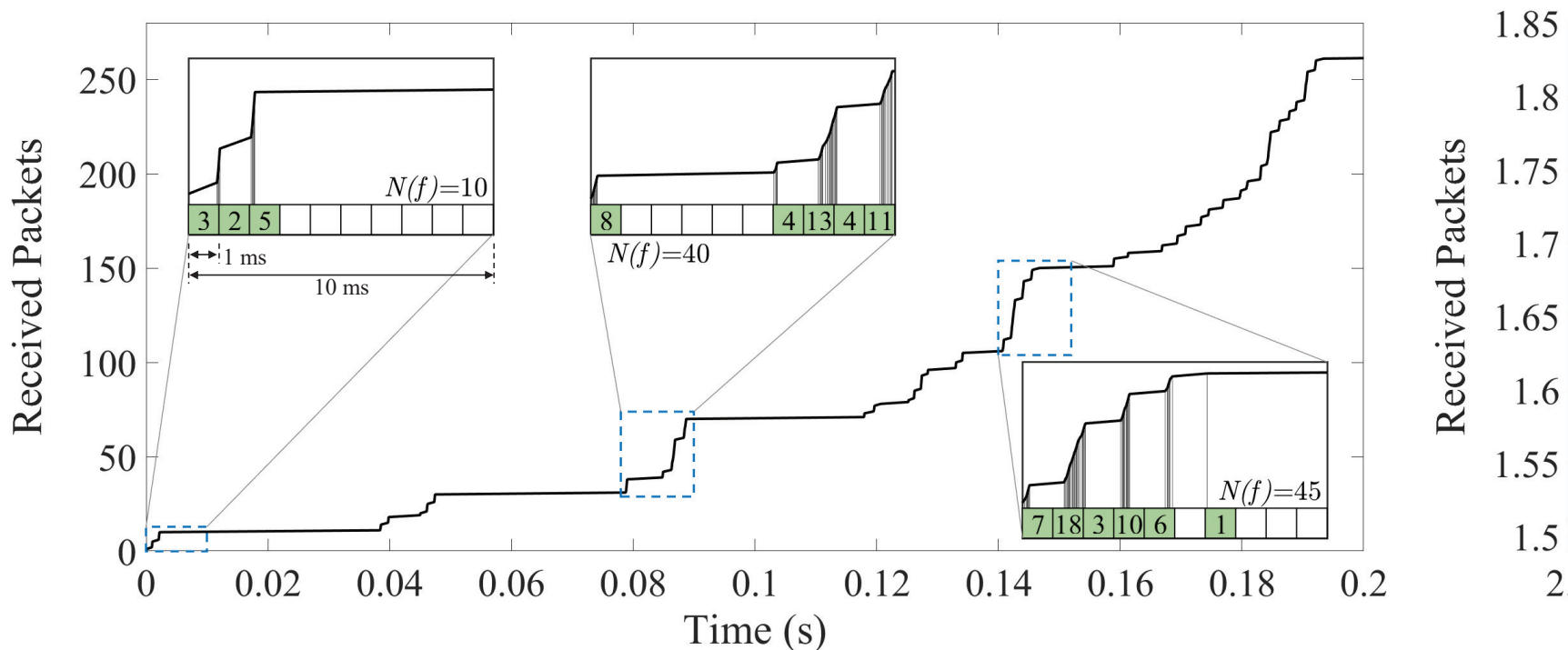


# Latency Limit without Throughput Loss



# Eliminating Probing Inefficiency

- Letting the cellular receiver infer the available throughput as well as minimum achievable RTT for tracking BDP
  - Packet reception pattern in the receiver has abundant information
    - Consecutive packet interval  $\rightarrow$  available PHY rate (at the chosen MCS)
    - Per-radio frame reception  $\rightarrow$  available PHY rate + user scheduling situation



# Window Control Implementation

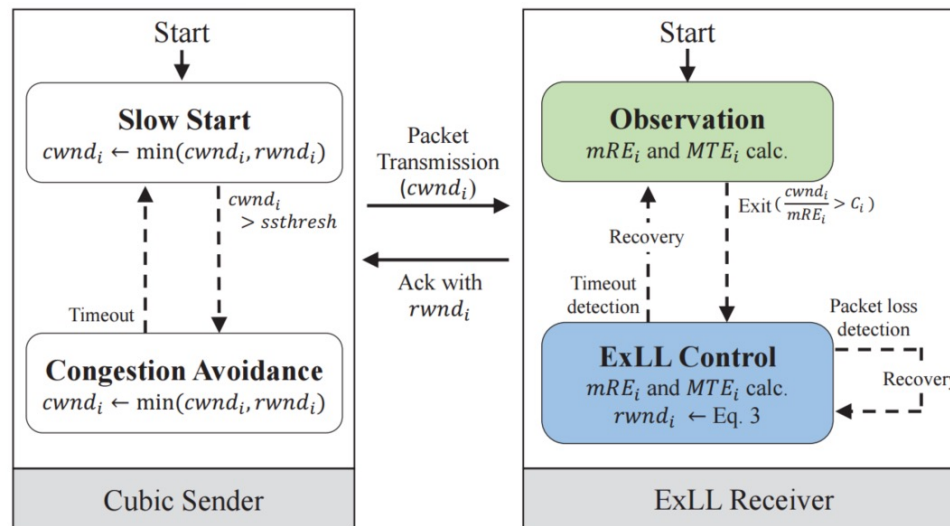
## □ ExLL: window control with **MTE** and **mRE**

- Converges to the solution of a NUM (network utilization maximization):  $\max_{x \geq 0} \sum_i \alpha_i \log x_i$

$$w_{i+1} = (1 - \gamma) \cdot w_i + \gamma \left( \frac{mRE_i}{R_i} \cdot w_i + \alpha \left( 1 - \frac{T_i}{MTE_i} \right) \right)$$

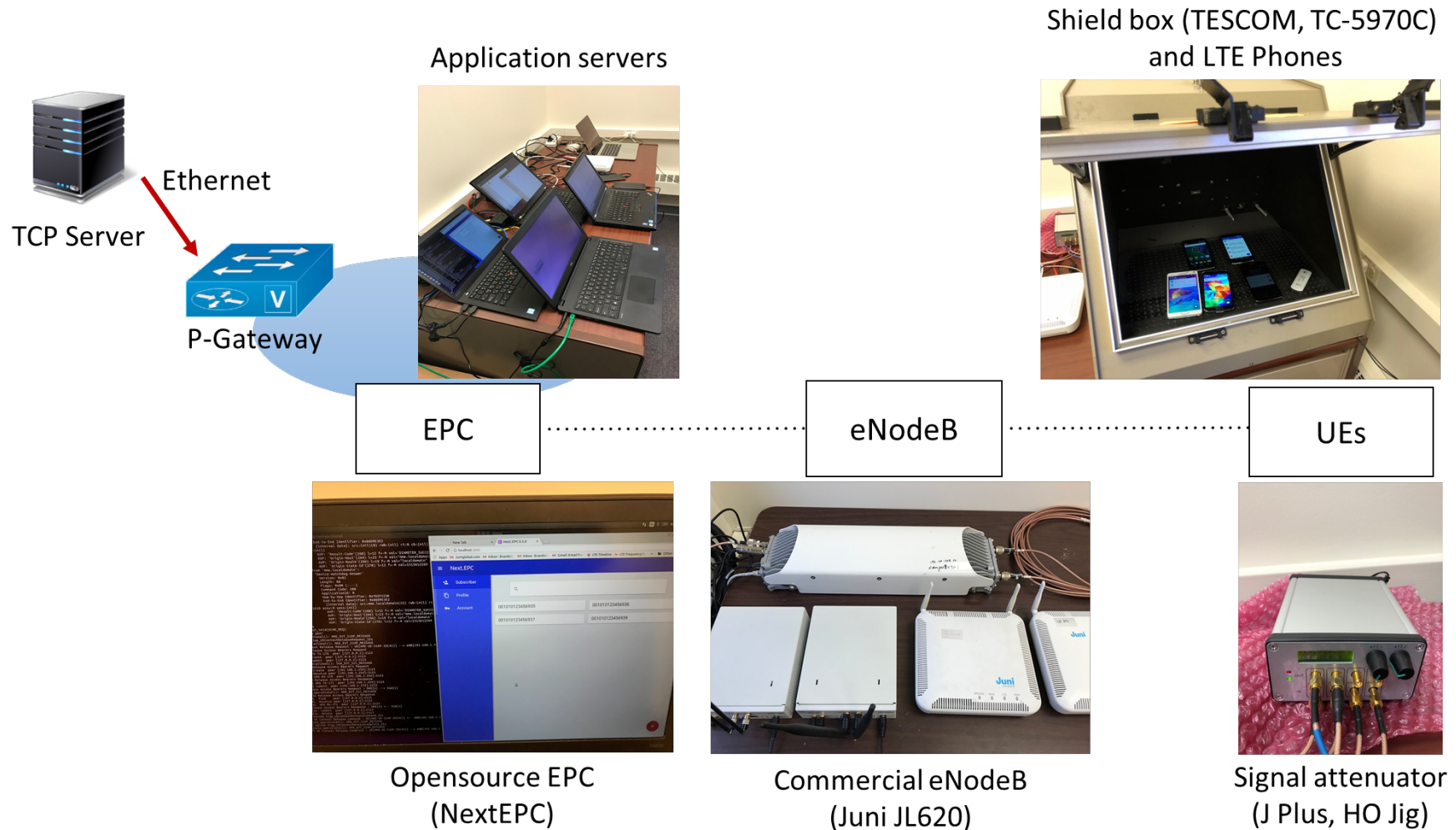
Probes minRTT

Probes max throughput, goes to 0 when MTE is achieved



ExLL receiver makes any CUBIC sender to be a low-latency sender.

# Testbed with Reproducible Environment

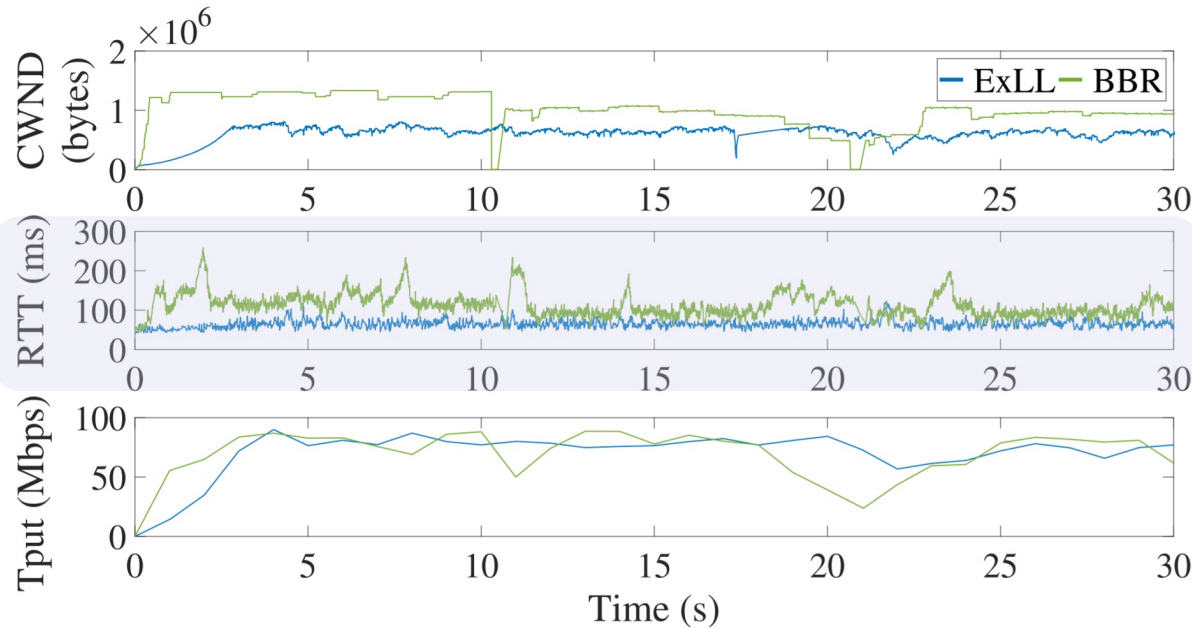


Open-source EPC  
(NextEPC)

Commercial eNodeB  
(Juni JL620)

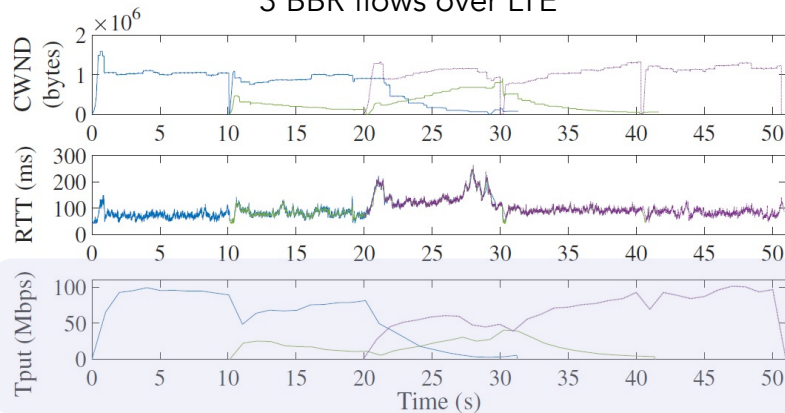
Signal attenuator  
(J Plus, HO Jig)

# ExLL Provides Near-Ideal Performance

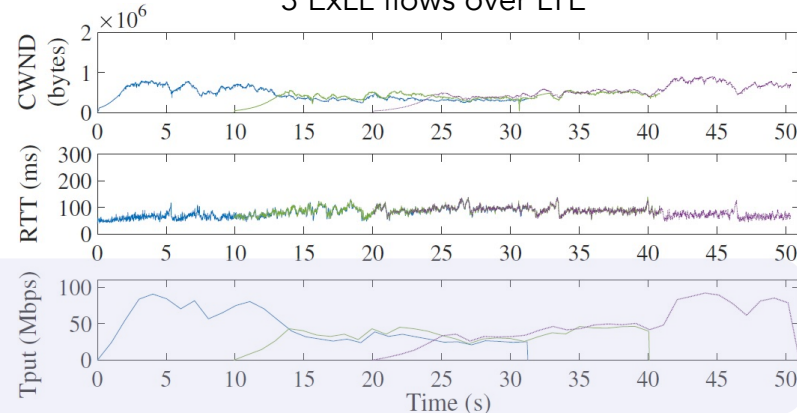


Extremely  
suppressed  
latency

3 BBR flows over LTE

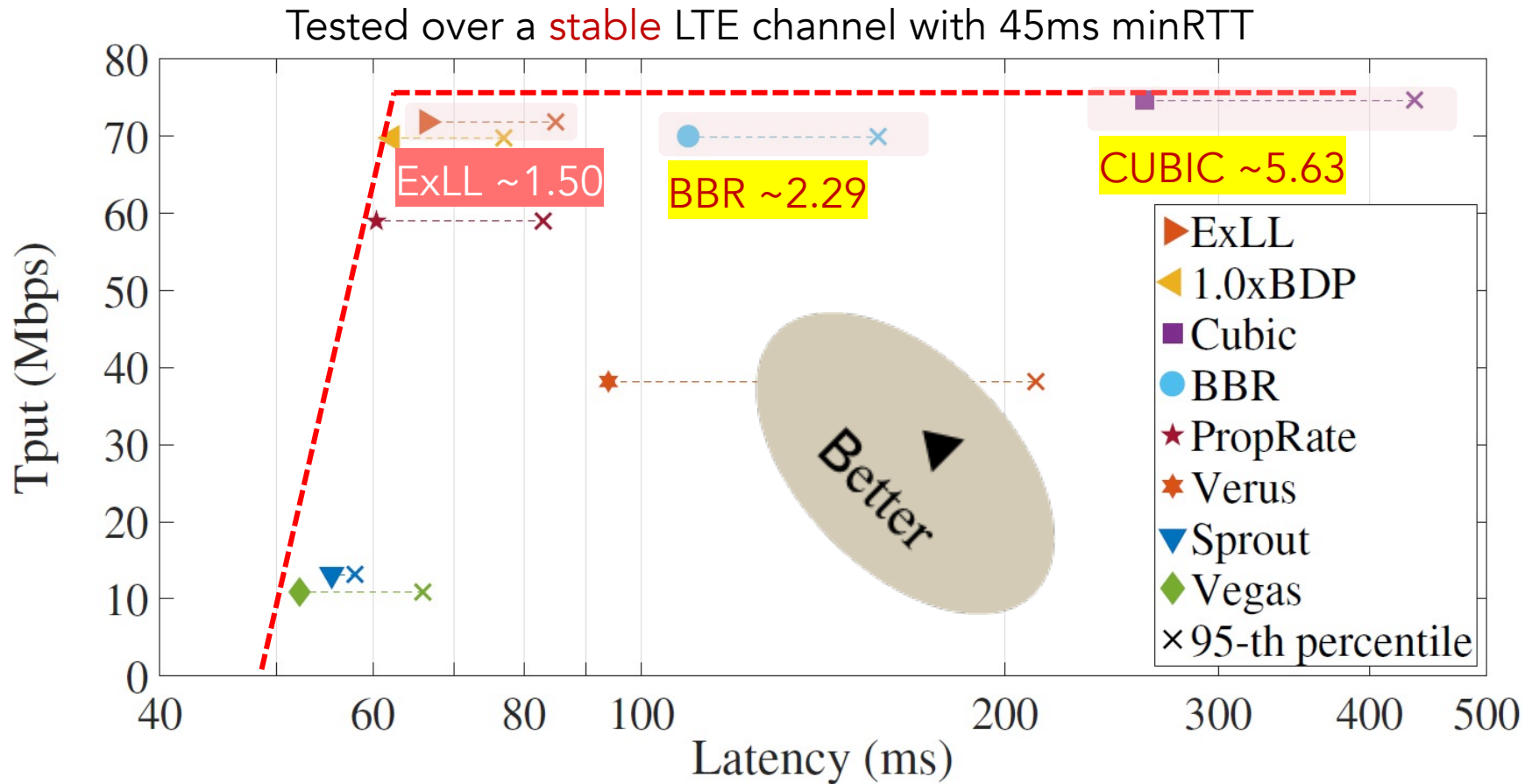


3 ExLL flows over LTE

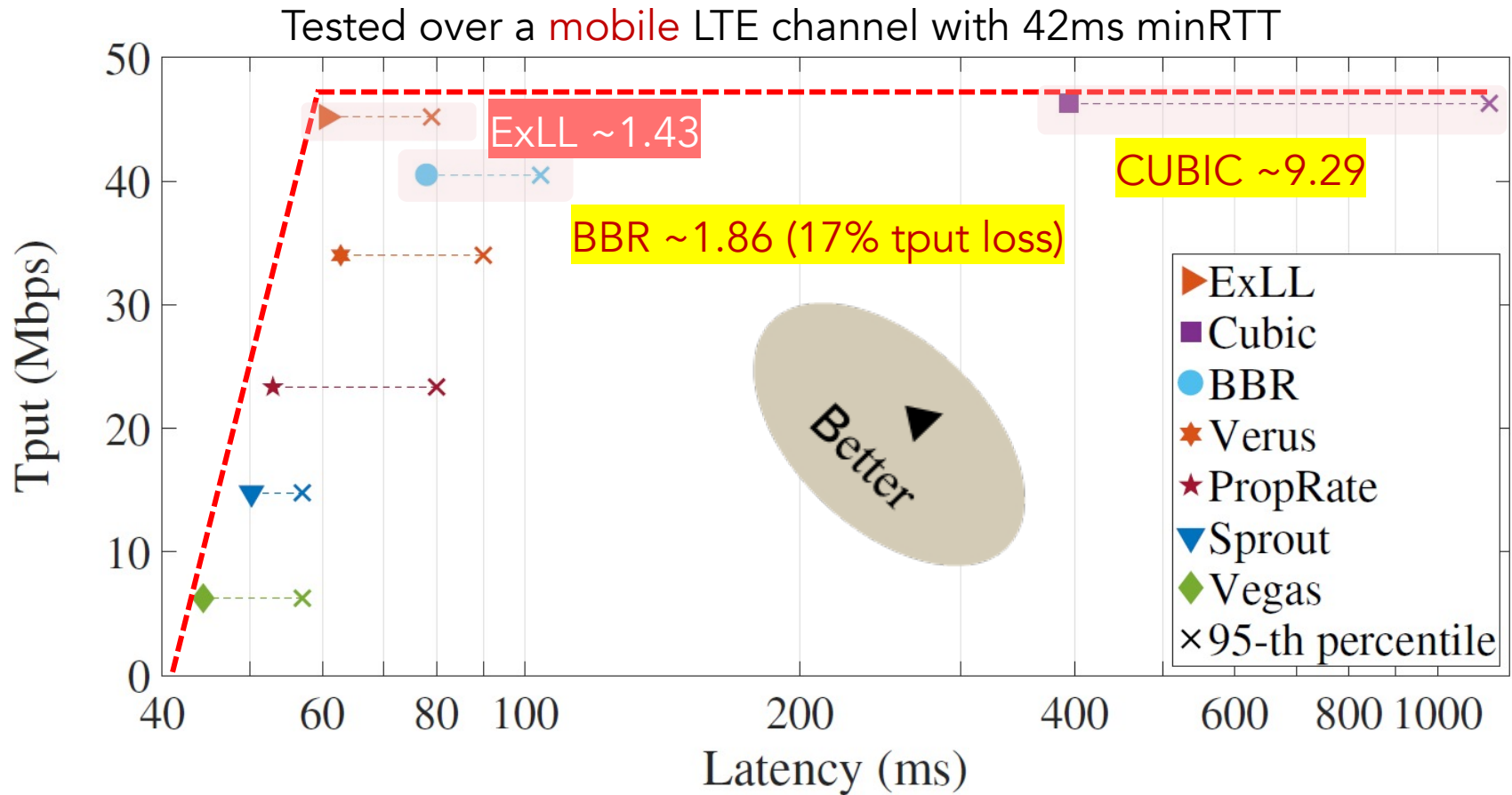


Guaranteed  
fairness

# ExLL Provides Near-Ideal Performance



# ExLL Provides Near-Ideal Performance



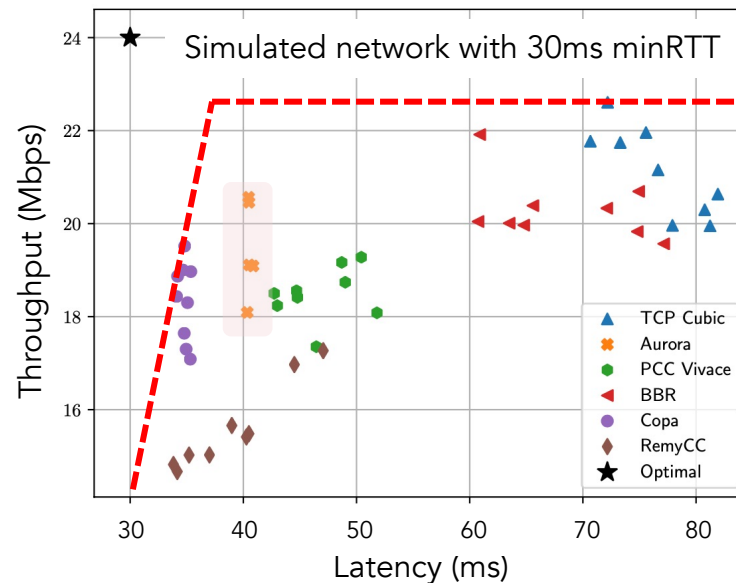


# Low-Latency Transport with ML

- People started to test RL (e.g., Aurora [ICML'19]) for congestion control

$$R_t = \mathbb{E}[\sum_t \gamma^t \cdot r_t]$$

$$r_t = 10 * \text{throughput} - 1000 * \text{latency} - 2000 * \text{loss}$$



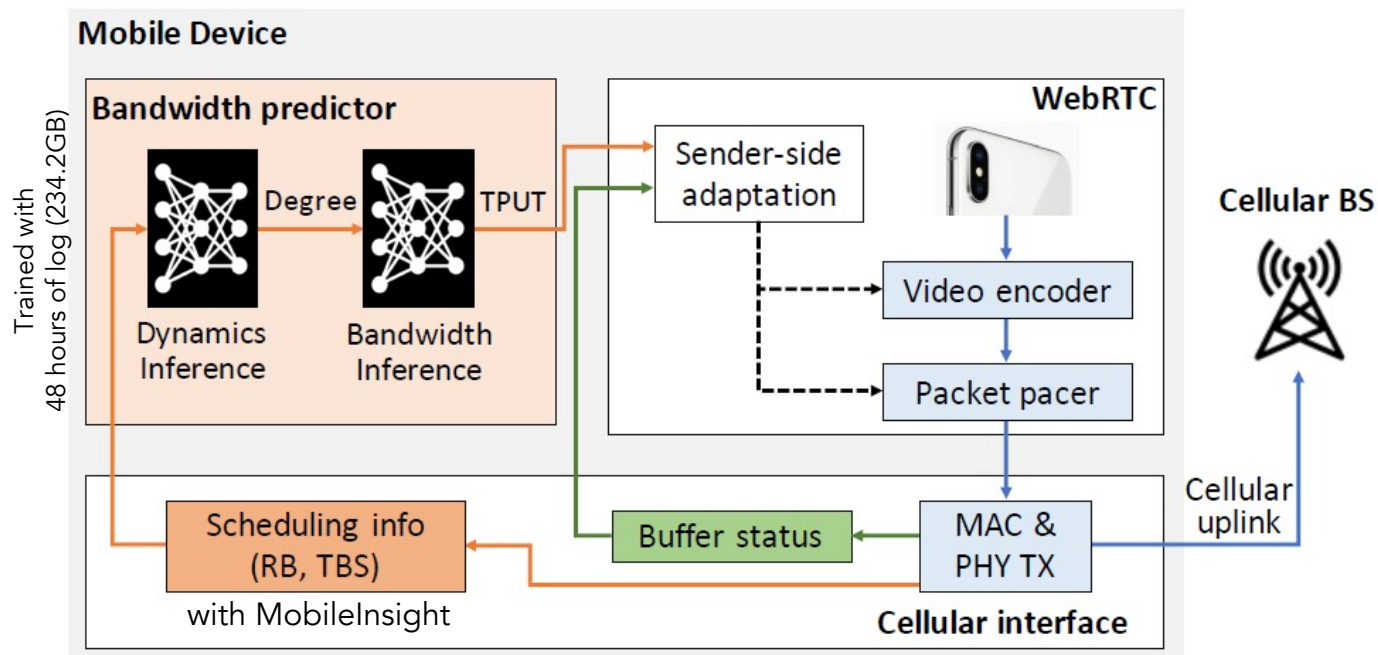
- Training for all possible cases is impossible. (Exploration inefficiency exists.)
- Fairness between flows has not been addressed.



# Low-Latency Transport with ML

- Throughput prediction for near-zero queueing (under mobility/handoff/user-contention)

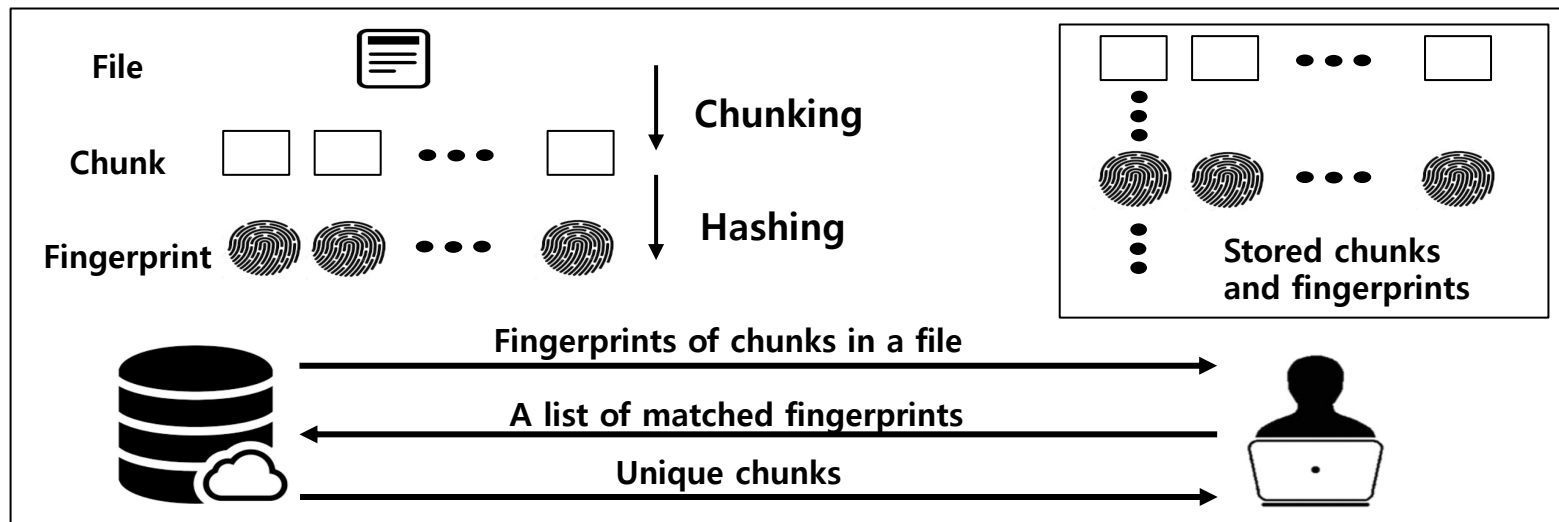
PERCEIVE [MobiSys'20] 2-stage LSTM on Pixel 3 for cellular uplinks gives  
 6.25% tput error @ 224ms  
 8.73% tput error @ 22.3ms (with 10x model compression)



# Challenge 2: Volume Reduction

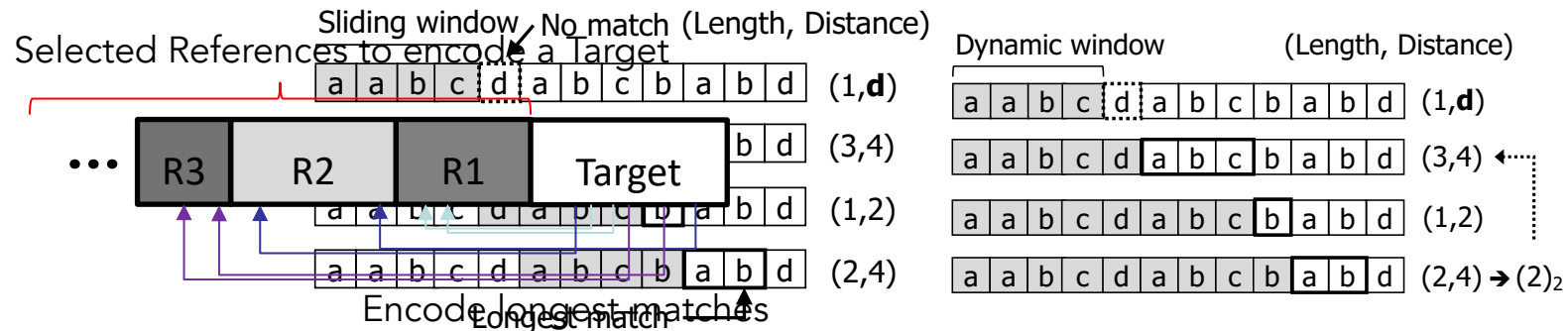
- RE is a network (middlebox) technique that **can remove duplicate data** from within arbitrary network flows for traffic volume reduction.

End-host RE ~ Transmission with Dedup



B. Aggarwal, A. Akella, A. Anand, A. Balachandran, P. Chitnis, C. Muthukrishnan, R. Ramjee and G. Varghese, "EndRE: An End-System Redundancy Elimination Service for Enterprises," NSDI 2010.

# SyncCoding: RE with Pre-Synced Data

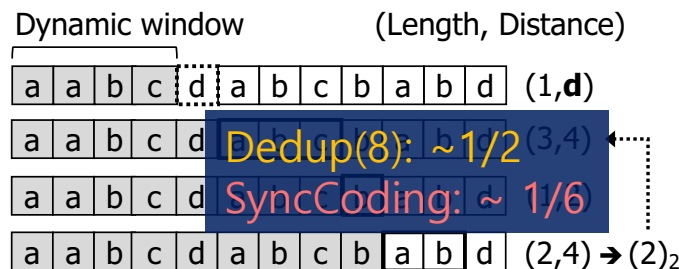


SyncCoding

vs.

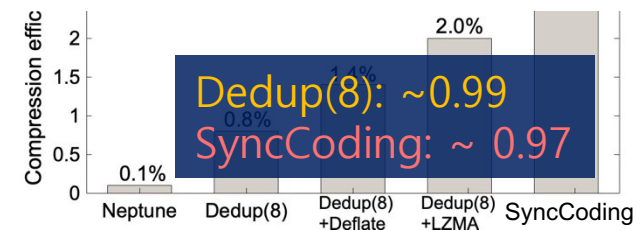
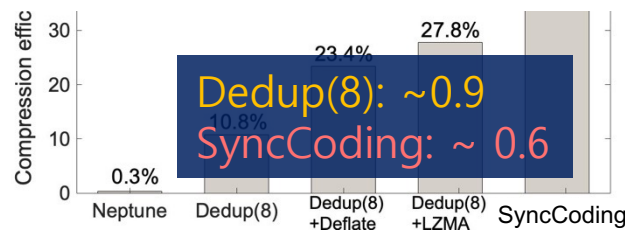
Dedup

(using n-byte

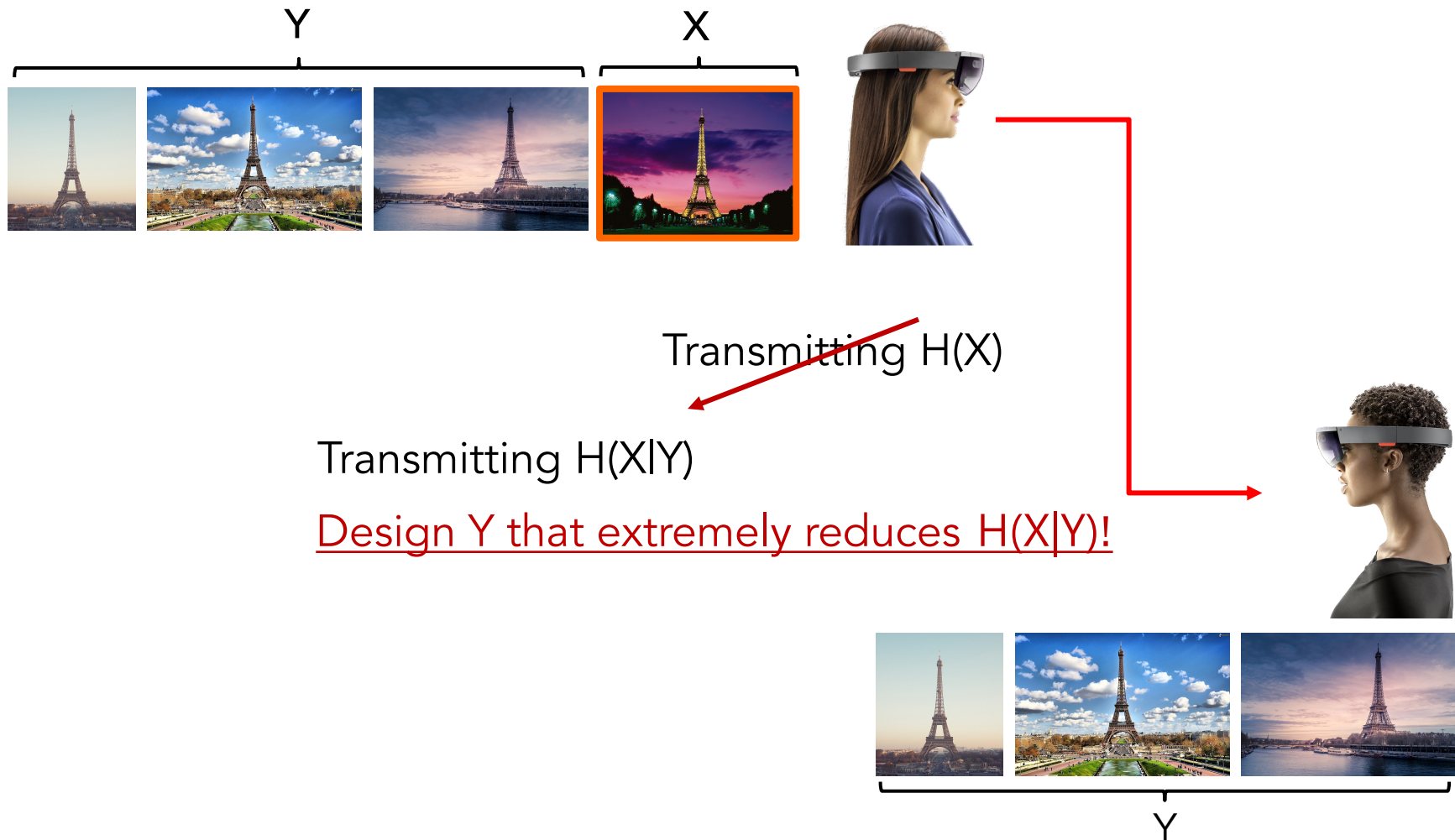


Dedup(8):  $\sim 1/2$   
SyncCoding:  $\sim 1/25$

Compression effi  
compressed amount  
original size



# Ultimate RE inside Transport Layer



# Challenge 3: Retransmission Delay

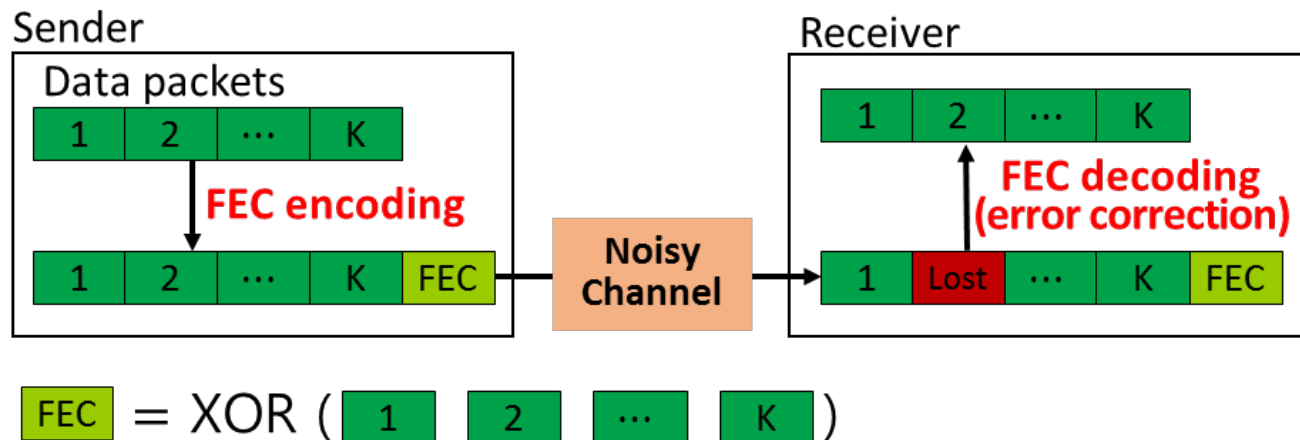
- HINT: QUIC (Quick UDP Internet Connections)
  - UDP based transport protocol with functions of TCP and TLS.
  - Standardized as HTTP/3 (from 2018.11)
  
- QUIC advantages
  - Connection establishment latency
  - Forward error correction
  - Improved congestion control
  - Multiplexing without head-of-line blocking
  - Connection migration



# QUIC Protocol: FEC

## □ Forward Error Correction

- To recover from lost packets without waiting for a retransmission
- QUIC uses a simple **XOR FEC (with session multiplexing)**



- A more efficient FEC can be employed for the same purpose.