

Numerical Analysis ~~in Mechanical Eng.~~ for Eng. Appls.

$$Ax = b$$

- Linear diff'l eqs.
- Nonlinear " "
- Steady diff'l eqs
- Unsteady " "

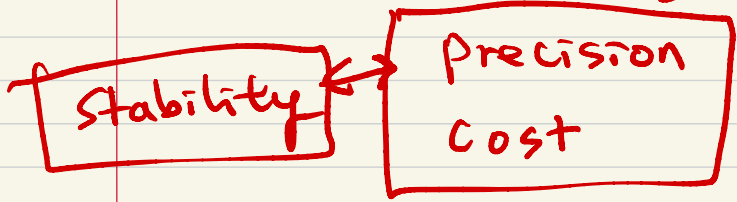
Numerical
methods

System of algebraic eqs.
(linear or nonlinear)

computer / numerical methods

Answers

$$x = A^{-1}b ?$$



Ch. 0 Linear Algebra

- Matrix-matrix multiplication

$$\begin{matrix} A & B & \longrightarrow & C = AB \\ m \times n & n \times l & & m \times l \end{matrix}$$

matrix elements $c_{ij} = a_{ik} b_{kj}$

$$i = 1, 2, \dots, m$$

$$j = 1, 2, \dots, l$$

$$k = 1, 2, \dots, n$$

Let the columns of B be denoted $\underline{b}^1, \underline{b}^2, \underline{b}^3, \dots, \underline{b}^l$

$$AB = [A\underline{b}^1, A\underline{b}^2, \dots, A\underline{b}^l]_{n \times l}$$

Operation counts

$$A\underline{b}^1$$

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & \dots & a_{2n} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ a_{m1} & a_{m2} & \dots & \dots & a_{mn} \end{bmatrix} \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{bmatrix}$$

multiplications:

$$m \times n$$

additions:

$$(n-1) \times m$$

if $m = n$, $\mathcal{O}(n^2)$ operations are required.

$A \ B \xrightarrow{\hspace{2cm}}$ $m \times n \times l$ multiplications
 $m \times n \times l$ additions

if $m = n = l$, $\mathcal{O}(n^3)$ operations.

ex) $n = 500 \rightarrow n^3 = 1.25 \times 10^8$

Intel Itanium 2 \rightarrow 300 MFlops
 \Downarrow floating-point operations per second
CRAY C90 3×10^8 operations/sec

$$\text{For } AB, \quad t = \frac{1.25 \times 10^8}{3 \times 10^8} \approx 0.3 \text{ sec}$$

in mid 1980's, IBM PC \rightarrow 8000 Fops $\Rightarrow t \doteq 4$ hrs.

NEC SX5 \rightarrow 3 GFops $\rightarrow t \approx 0.03$ sec

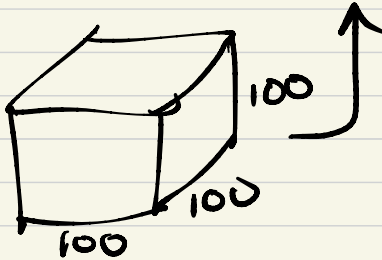
GAIA 1536 cores \rightarrow 23,370 TFops $\rightarrow t = 0.5 \times 10^{-5}$ sec

GPU Nvidia A100 \rightarrow 20 TFlops

CPU AMD Ryzen Threadripper 3990X

\rightarrow 3 TFlops $\rightarrow t = 4 \times 10^5$ sec.

How about $n = 10^6 \rightarrow n^3 = 10^{18}$



Intel Itanium 2 $\rightarrow t = \frac{10^{18}}{3 \times 10^8} \approx 10^6$ hrs

NEC Sx5 $\rightarrow t = 10^5$ hrs $\rightarrow 10^9$ KW

(1M Korean Won for 100 CPU hrs)

GAI A $t = \frac{10^{18}}{23.37 \times 10^{12}} \approx 42$ sec.

\Rightarrow Matrix-matrix multiplication is
a very expensive job!

$$Ax = b$$
$$x = A^{-1}b$$

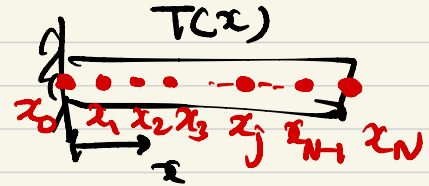
AB

• Banded matrix

diff'l eq $\xrightarrow{\text{numerical method}}$ system of algebraic eqs.

○ ODE (ordinary diff'l eq)

$$\boxed{\frac{d^2 T}{dx^2} - a^2 T = Q(x)}$$



Introduce a discrete set of points $x_j \quad j=0, 1, 2, \dots, N$
grid pts.

Find $T(x_j)$ or T_j

finite difference method (FDM)

Taylor series expansion $h_j = x_{j+1} - x_j$: grid spacing

$h_1 = h_2 = \dots = h$: uniform
grid spacing

$$T(x_{j+1}) = T(x_j) + h \frac{dT}{dx}|_j + \frac{1}{2} h^2 \frac{d^2 T}{dx^2}|_j + \frac{1}{6} h^3 \frac{d^3 T}{dx^3}|_j + \frac{1}{24} h^4 \frac{d^4 T}{dx^4}|_j + \dots$$

$$- \left[T(x_{j-1}) = \dots - \dots + \dots - \dots + \dots - \dots \right]$$

$$T(x_{j+1}) - T(x_{j-1}) = 2h \frac{dT}{dx}|_j + \frac{1}{3} h^3 \frac{d^3 T}{dx^3}|_j + \dots$$

$$\Rightarrow \boxed{\frac{dT}{dx}|_j = \frac{T(x_{j+1}) - T(x_{j-1}))}{2h}} - \frac{1}{6} h^2 \frac{d^3 T}{dx^3}|_j + \dots$$

↑
FDM

leading error term

$$\left(\frac{dT}{dx} = \lim_{\Delta x \rightarrow 0} \frac{T(x_{j+1}) - T(x_{j-1}))}{\Delta x} \right)$$

$\mathcal{O}(h^2)$

Second-order FDM

$$\boxed{\frac{dT}{dx}|_j = \frac{T(x_{j+1}) - T(x_j)}{h}} - \frac{1}{2} h \frac{d^2 T}{dx^2}|_j + \dots$$

first-order FDM

$$\frac{d^2 T}{dx^2} - dT = Q$$

+

$$T(x_{j+1}) + T(x_{j-1}) = 2T(x_j) + h^2 \frac{d^2 T}{dx^2} \Big|_j + \frac{1}{12} h^4 \frac{d^4 T}{dx^4} \Big|_j + \dots$$

$$\Rightarrow \frac{d^2 T}{dx^2} \Big|_j = \frac{T(x_{j+1}) - 2T(x_j) + T(x_{j-1}))}{h^2} - \frac{1}{12} h^2 \frac{d^4 T}{dx^4} \Big|_j + \dots$$

second-order FDM

leading error term $O(h^2)$

→ substitute this to diff'l eq.

$$\textcircled{a} \bar{j}, \begin{cases} \frac{T_{\bar{j}+1} - 2T_{\bar{j}} + T_{\bar{j}-1}}{h^2} - dT_{\bar{j}} = Q_{\bar{j}} \\ T_0 = 0 \\ T_N = S \end{cases}$$

$$\bar{j} = 1, 2, \dots, N-1$$



boundary conditions

$$\frac{d^2 T}{dx^2} - \alpha^2 T = Q(x)$$

$$\left\{ \begin{array}{l} \frac{T_{j+1} - 2T_j + T_{j-1}}{h^2} - \alpha^2 T_j = Q_j, \quad j=1, 2, \dots, N-1 \\ T_0 = 0 \\ T_N = S \end{array} \right.$$

System of algebraic eqs.

$$\rightarrow \left\{ \begin{array}{l} T_{j+1} - (2 + h^2 \alpha^2) T_j + T_{j-1} = h^2 Q_j \\ T_0 = 0 \\ T_N = S \end{array} \right.$$

$$j=1: T_2 - (2 + h^2 \alpha^2) T_1 + T_0 = h^2 Q_1$$

$$j=2: T_3 - \quad \quad T_2 + T_1 = h^2 Q_2$$

\vdots

$$j=N-1: T_N - \quad \quad T_{N-1} + T_{N-2} = h^2 Q_{N-1}$$

$$\begin{bmatrix}
 -(2+h^2\alpha^2) & 1 & 0 & \dots & 0 \\
 1 & -(2+h^2\alpha^2) & 1 & 0 & \dots & 0 \\
 0 & \dots & 0 & 1 & -(2+h^2\alpha^2) & 1 & 0 & \dots & 0 \\
 & & & & & & & & \\
 & & & & & & & & 1 & -(2+h^2\alpha^2)
 \end{bmatrix}
 \begin{bmatrix}
 T_1 \\
 T_2 \\
 \vdots \\
 T_j \\
 \vdots \\
 T_N
 \end{bmatrix}
 =
 \begin{bmatrix}
 h^2 Q_1 \\
 h^2 Q_2 \\
 \vdots \\
 h^2 Q_j \\
 \vdots \\
 h^2 Q_{N-1} \sim f
 \end{bmatrix}$$

tri-diagonal matrix $\underline{Ax} = b$

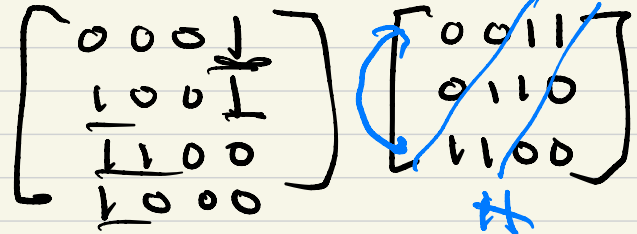
banded matrix : non-zero elements only around the main diagonal

↳ arises from FDM of diff'l eqs.

tri-diagonal matrix : $B [a_i, b_i, c_i]$

$$a_i = 1, \quad b_i = -(2+h^2\alpha^2), \quad c_i = 1$$

banded matrix \longleftrightarrow sparse matrix



② PDE (partial diff'l eq.)

$$\frac{\partial^2 T}{\partial x^2} + \frac{\partial^2 T}{\partial y^2} = Q(x, y)$$

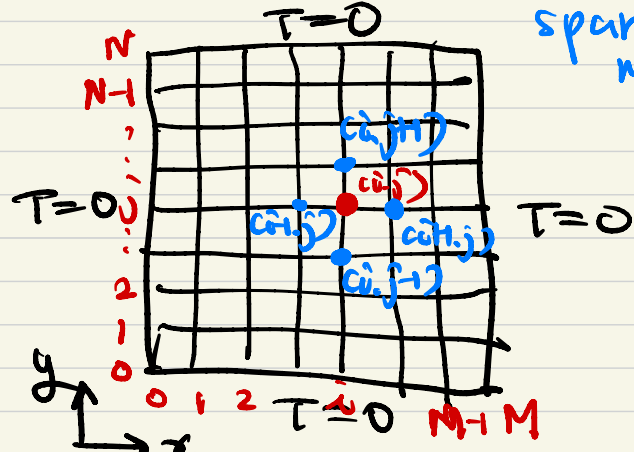
$\Delta x_i = x_{i+1} - x_i = h$
 $\Delta y_j = y_{j+1} - y_j = h$) uniform grid spacings

2nd-order FDM

$$\frac{T_{i+1,j} - 2T_{i,j} + T_{i-1,j}}{h^2} + \frac{T_{i,j+1} - 2T_{i,j} + T_{i,j-1}}{h^2} = Q_{i,j}$$

$i = 1, 2, \dots, M-1$

$j = 1, 2, \dots, N-1$



⚡ sparse matrix

$$\Rightarrow T_{i+1,j} - 4T_{i,j} + T_{i-1,j} + T_{i,j+1} + T_{i,j-1} = h^2 Q_{i,j}$$

$M=4$
 $N=4$

$$\bar{i}=1, \bar{j}=1 : T_{2,1} - 4T_{1,1} + T_{0,1} + T_{1,2} + T_{1,0} = h^2 Q_{1,1}$$

$$\bar{i}=2, \bar{j}=1 : T_{3,1} - 4T_{2,1} + T_{1,1} + T_{2,2} + T_{2,0} = h^2 Q_{2,1}$$

$$\vdots$$

$$\bar{i}=3, \bar{j}=3 : T_{4,3} - 4T_{3,3} + T_{2,3} + T_{3,4} + T_{3,2} = h^2 Q_{3,3}$$

9 unknowns

$T_{1,1} \dots T_{3,3}$

-4	1	0	1	0	0	0	0	0	0	$T_{1,1}$ $T_{2,1}$ $T_{3,1}$ $T_{1,2}$ $T_{2,2}$ $T_{3,2}$ $T_{3,1}$ $T_{3,2}$ $T_{3,3}$
1	-4	1	0	1	0	0	0	0	0	
0	1	-4	0	0	1	0	0	0	0	
1	0	0	-4	1	0	1	0	0	0	
0	1	0	1	-4	1	0	1	0	0	
0	0	1	0	1	-4	0	0	1	0	
0	0	0	1	0	0	-4	1	0	0	
0	0	0	0	1	0	1	-4	1	0	
0	0	0	0	0	1	0	1	-4	0	

$h^2 Q_{1,1}$ $h^2 Q_{2,1}$ \vdots \vdots \vdots \vdots \vdots $h^2 Q_{3,3}$	$)$
---	-----

Block-tridiagonal matrix

$$\begin{pmatrix} B_1 & C_1 & 0 & & 0 \\ A_2 & B_2 & C_2 & 0 & \\ 0 & A_3 & B_3 & C_3 & \\ 0 & & \dots & \dots & \dots \end{pmatrix}$$

⑤ Solution technique

• diagonal matrix

$$\begin{pmatrix} 1 & 0 & 0 \\ 0 & 2 & 0 \\ 0 & 0 & 3 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 1 \\ 2 \\ 3 \end{pmatrix}$$

• upper triangular matrix

$$\begin{pmatrix} 1 & 2 & 3 \\ 0 & 4 & 5 \\ 0 & 0 & 6 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 7 \\ 8 \\ 9 \end{pmatrix}$$

• lower triangular matrix

$$\begin{pmatrix} 1 & 0 & 0 \\ 2 & 3 & 0 \\ 4 & 5 & 6 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 7 \\ 8 \\ 9 \end{pmatrix}$$

* Gauss elimination (GE) : to make upper triangular matrix

$$\begin{cases} 4x_2 - x_3 = 5 \\ x_1 + x_2 + x_3 = 6 \\ 2x_1 - 2x_2 + x_3 = 1 \end{cases} \rightarrow \begin{pmatrix} 0 & 4 & -1 \\ 1 & 1 & 1 \\ 2 & -2 & 1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 5 \\ 6 \\ 1 \end{pmatrix}$$

Interchange the first two rows

$$\begin{pmatrix} 1 & 1 & 1 \\ 0 & 4 & -1 \\ 2 & -2 & 1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 6 \\ 5 \\ 1 \end{pmatrix}$$

Subtract twice the first row from last row:

$$\begin{pmatrix} 1 & 1 & 1 \\ 0 & 4 & -1 \\ 0 & -4 & -1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 6 \\ 5 \\ -11 \end{pmatrix} \rightarrow \begin{pmatrix} 1 & 1 & 1 \\ 0 & 4 & -1 \\ 0 & 0 & -2 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 6 \\ 5 \\ -6 \end{pmatrix}$$

upper triangular matrix

Stop
GE.

or, augmented matrix

$$\left(\begin{array}{ccc|c} 0 & 4 & 7 & 5 \\ 1 & 1 & 1 & 6 \\ 2 & -2 & 1 & 1 \end{array} \right) \xrightarrow{GE} \left(\begin{array}{ccc|c} 1 & 1 & 1 & 6 \\ 0 & 4 & -1 & 5 \\ 0 & 0 & -2 & -6 \end{array} \right)$$

$$\rightarrow \left(\begin{array}{ccc|c} 1 & 1 & 1 & 6 \\ 0 & 4 & 0 & 8 \\ 0 & 0 & -2 & -6 \end{array} \right) \rightarrow \left(\begin{array}{ccc|c} 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 2 \\ 0 & 0 & 1 & 3 \end{array} \right)$$

$$\rightarrow x_1 = 1, x_2 = 2, x_3 = 3$$

Gauss-Jordan elimination

• General matrix system

$$Ax = b$$

$n \times n$

(GE)

$$\left(\begin{array}{cccc} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \dots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} \end{array} \right) \times \frac{a_{22}}{a_{11}}$$

0

To eliminate a_{21} , $l_2 = a_{21}/a_{11}$.

Multiply the first row by l_2 , and subtract from 2nd row.

Continue until all elts below a_{11} are eliminated.

$$\begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ 0 & a'_{22} & \dots & a'_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & a''_{22} & \dots & a''_{2n} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} = \begin{pmatrix} b_1 \\ b'_2 \\ \vdots \\ b_n \end{pmatrix}$$

By same way (we drop (') for convenience)

$$\begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ 0 & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & & & a_{nn} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{pmatrix}$$

Stop G.E.

backward sweep

$$\begin{cases} x_n = b_n / a_{nn} \\ x_{n-1} = (b_{n-1} - a_{n-1,n} x_n) / a_{n-1,n-1} \\ x_j = (b_j - \sum_{k=j+1}^n a_{jk} x_k) / a_{jj} \end{cases} \quad j = n-1, n-2, \dots, 1$$

* Operation counts for obtaining sol. of $Ax=b$
 $Ax=b$ A : full matrix

$$b = \begin{pmatrix} a_{21} \\ a_{11} \end{pmatrix} \times \begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & & a_{2n} \\ \vdots & \vdots & & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{pmatrix}$$

GE To eliminate a_{21} : $1D, nM, nA$
 To " the first column: $(n-1)D, n(n-1)M, n(n-1)A$
 To " " 2nd " : $(n-2)D, (n-2)(n-1)M, (n-2)(n-1)A$
 \vdots
 To " " $(n-1)^{th}$ " : $1D, 1-2M, 1-2A$

Total divisions: $\sum_{k=2}^{n-1} k = \frac{1}{2}n(n-1)$

" multiplications: $\sum_{k=1}^{n-1} k(k+1) = \frac{1}{3}(n^3 - n)$

additions

$$\frac{1}{2}n^2 = \frac{1}{6}n(n+1)(2n+1)$$

$\Rightarrow \frac{1}{2}n(n-1)D, \frac{1}{3}(n^3-n)M, \frac{1}{3}(n^3-n)A$ $\mathcal{O}(n^3)$
Gauss elimination requires $\mathcal{O}(n^3/3)$!

Backward sweep

$$z_n = b_n / a_{nn}$$

$$x_j = (b_j - \sum_{k=j+1}^n a_{jk} z_k) / a_{jj}, \quad j = n-1, n-2, \dots, 1$$

for each j , $1D, (n-j)M, (n-j)A$

total $\sum_{j=1}^{n-1} (n-j) = \frac{1}{2}n(n-1)A, M$ $\mathcal{O}(n^2)$

nD

negligible as compared to GE requiring $\mathcal{O}(n^3)$ for $n \gg 1$.

* Gauss-Jordan elimination requires same operation counts.